

# Graphical Models

## Discrete Inference and Learning

### Lecture 2

MVA

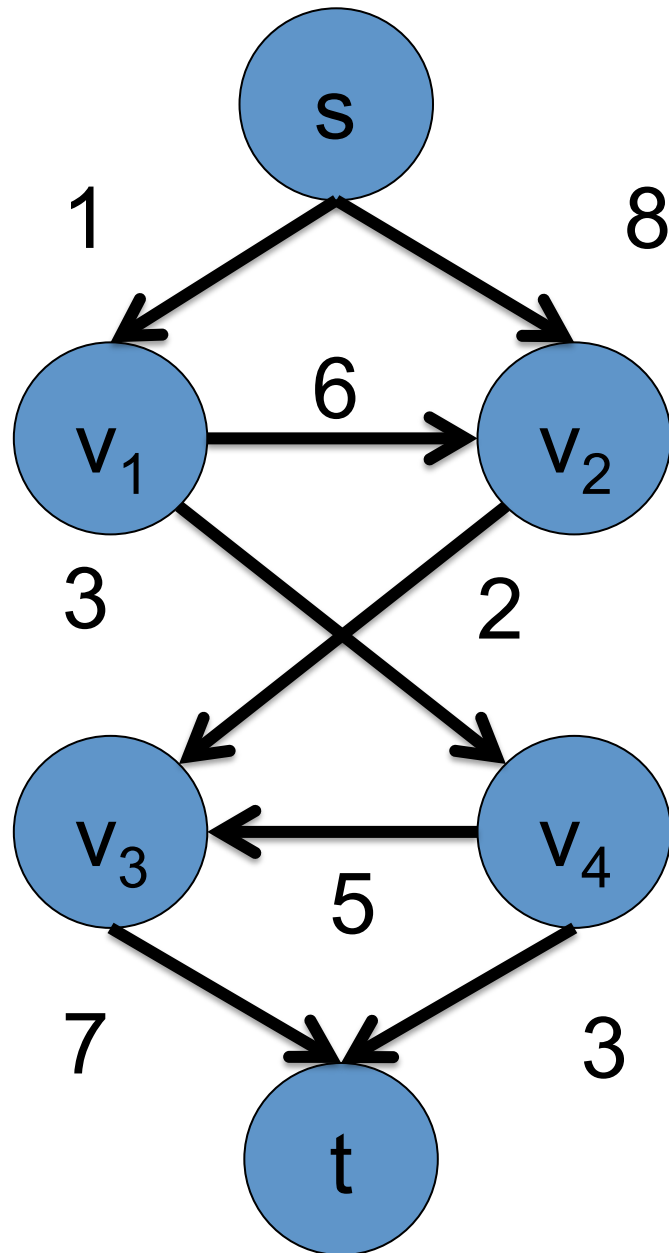
2019 – 2020

<http://thoth.inrialpes.fr/~alahari/disinflern>

# Outline

- Preliminaries
  - **s-t Flow**
  - s-t Cut
  - Flows vs. Cuts (details in next lecture)
- Maximum Flow
- Algorithms
- Energy minimization with max flow/min cut

# s-t Flow



Function flow:  $A \rightarrow R$

Flow of arc  $\leq$  arc capacity

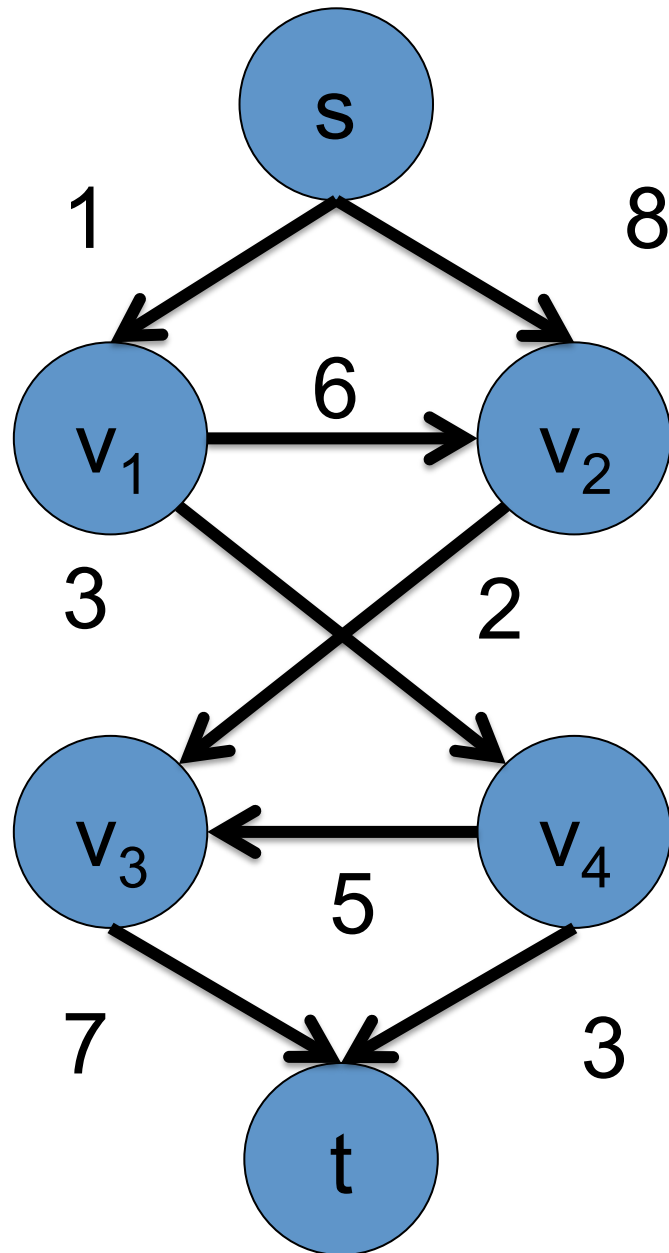
Flow is non-negative

For all vertex except  $s, t$

Incoming flow

= Outgoing flow

# s-t Flow



Function flow:  $A \rightarrow R$

$\text{flow}(a) \leq c(a)$

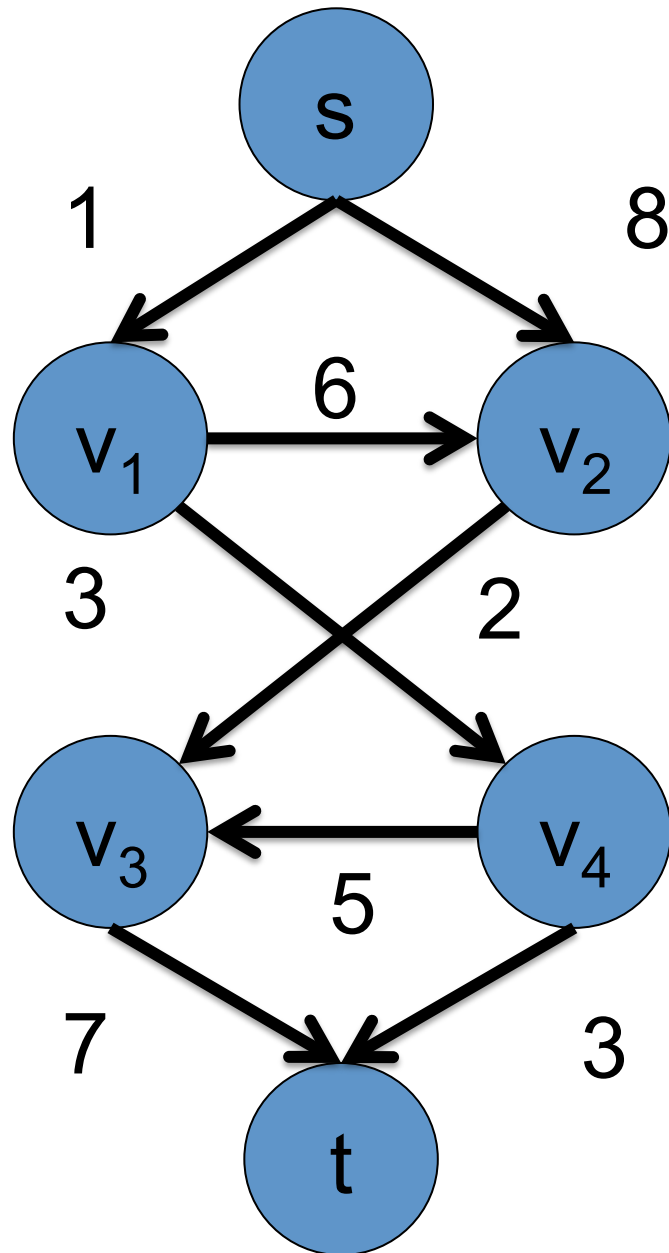
Flow is non-negative

For all vertex except  $s, t$

Incoming flow

= Outgoing flow

# s-t Flow



Function flow:  $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

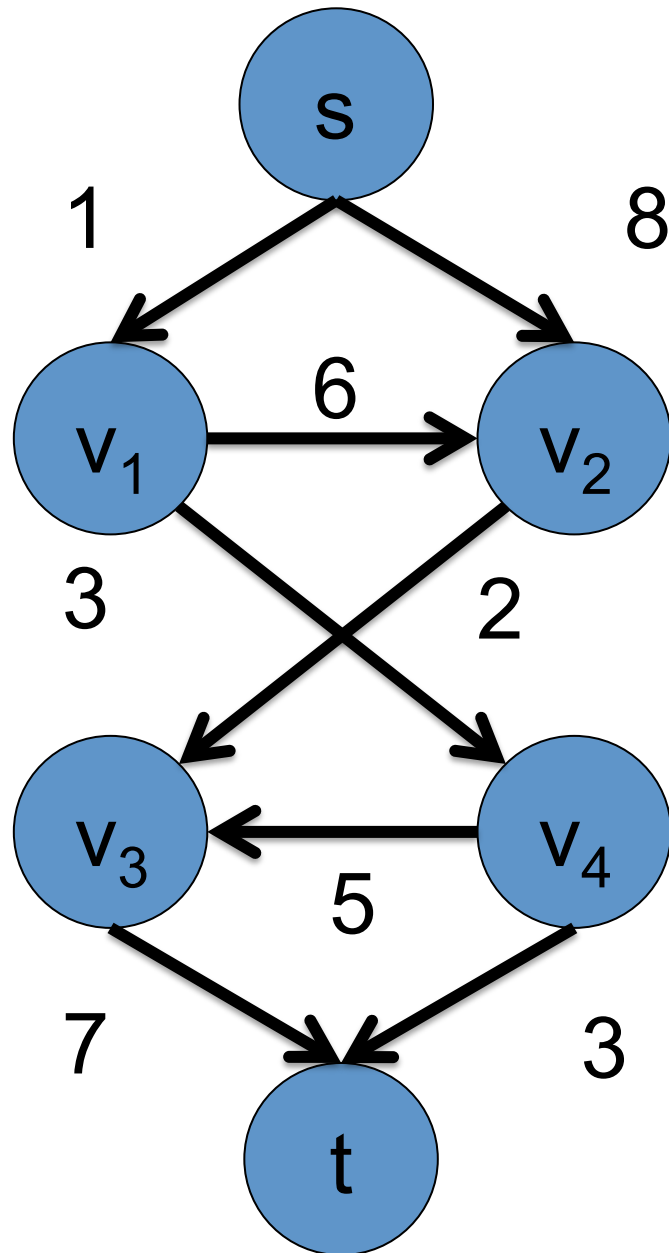
$$\text{flow}(a) \geq 0$$

For all vertex except  $s, t$

Incoming flow

= Outgoing flow

# s-t Flow



Function flow:  $A \rightarrow R$

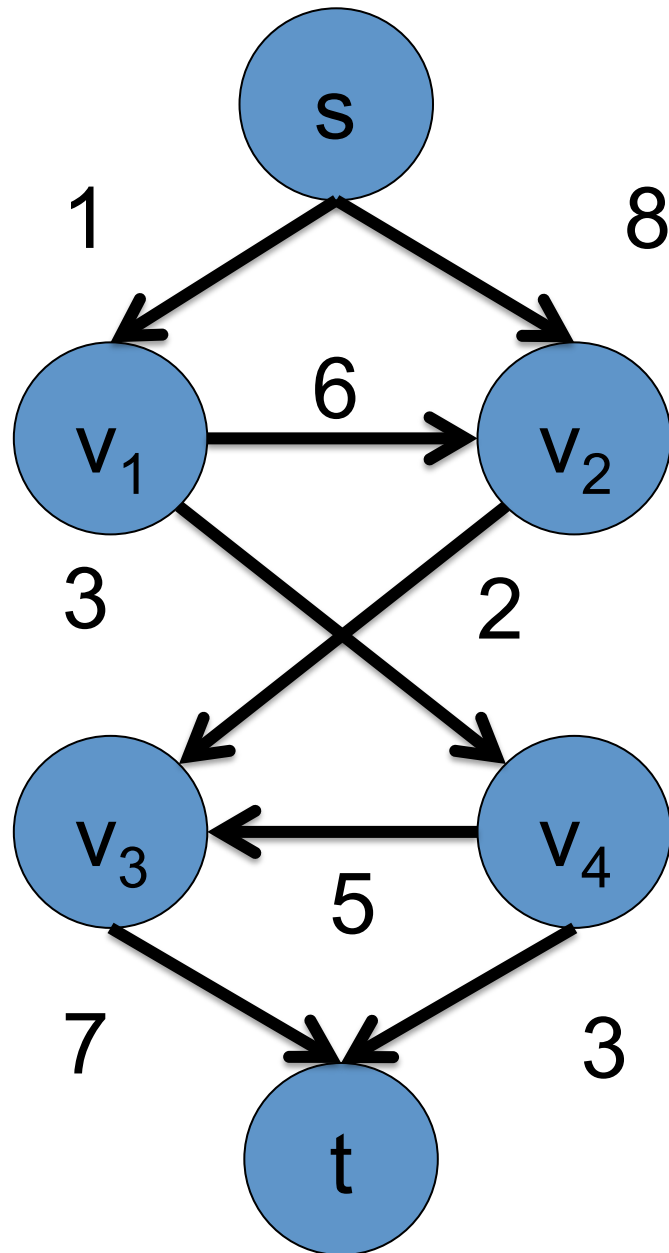
$$\text{flow}(a) \leq c(a)$$

$$\text{flow}(a) \geq 0$$

For all  $v \in V \setminus \{s, t\}$

Incoming flow  
= Outgoing flow

# s-t Flow



Function flow:  $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

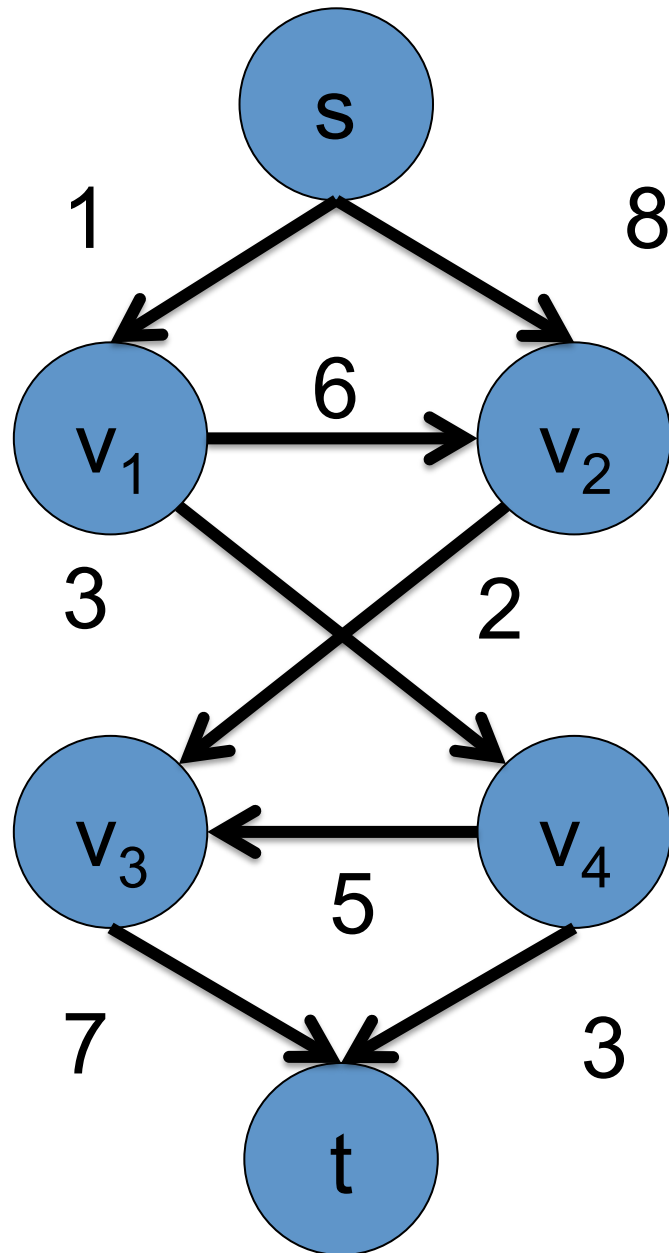
$$\text{flow}(a) \geq 0$$

For all  $v \in V \setminus \{s, t\}$

$$\sum_{(u,v) \in A} \text{flow}((u,v))$$

= Outgoing flow

# s-t Flow



Function flow:  $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

$$\text{flow}(a) \geq 0$$

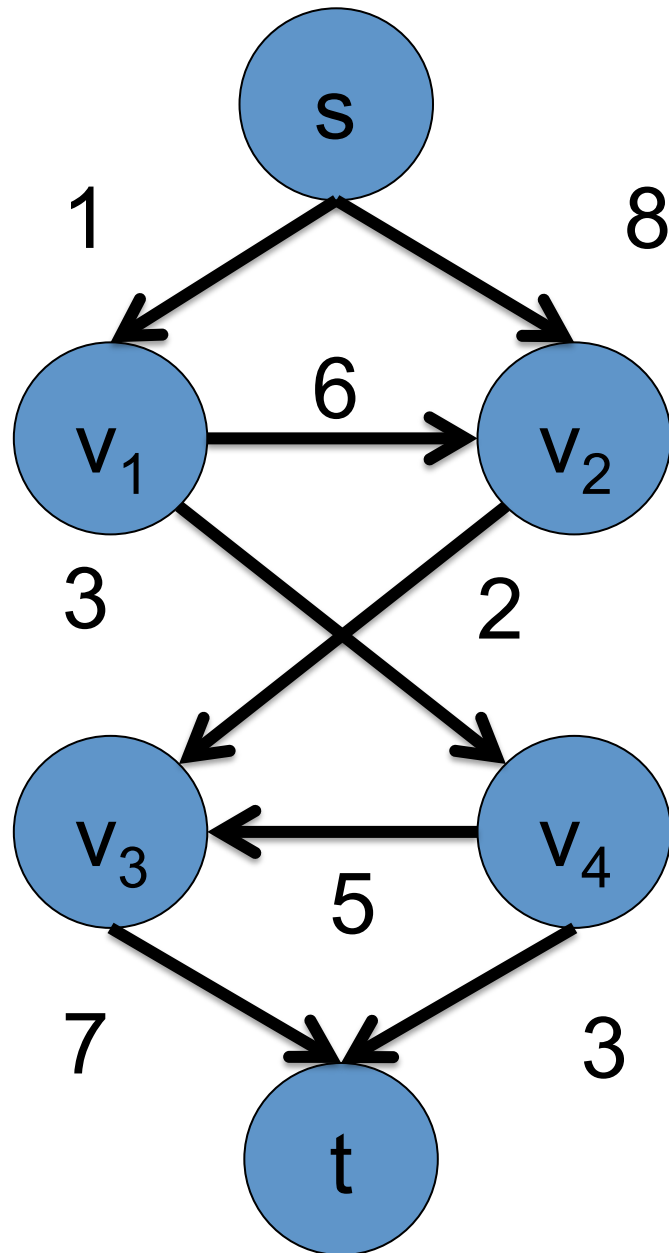
For all  $v \in V \setminus \{s, t\}$

$$\sum_{(u,v) \in A} \text{flow}((u,v))$$

$$= \sum_{(v,u) \in A} \text{flow}((v,u))$$



# s-t Flow



Function flow:  $A \rightarrow R$

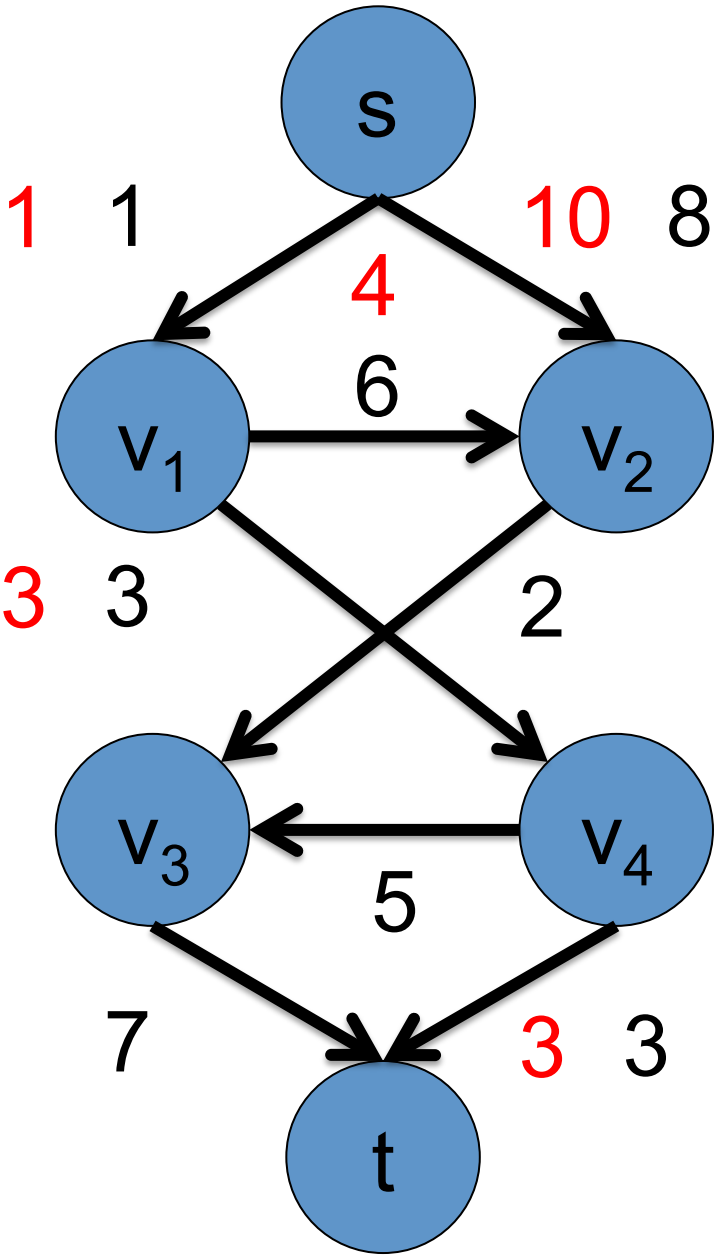
$$\text{flow}(a) \leq c(a)$$

$$\text{flow}(a) \geq 0$$

For all  $v \in V \setminus \{s, t\}$

$$E_{\text{flow}}(v) = 0$$

# s-t Flow



Function flow:  $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

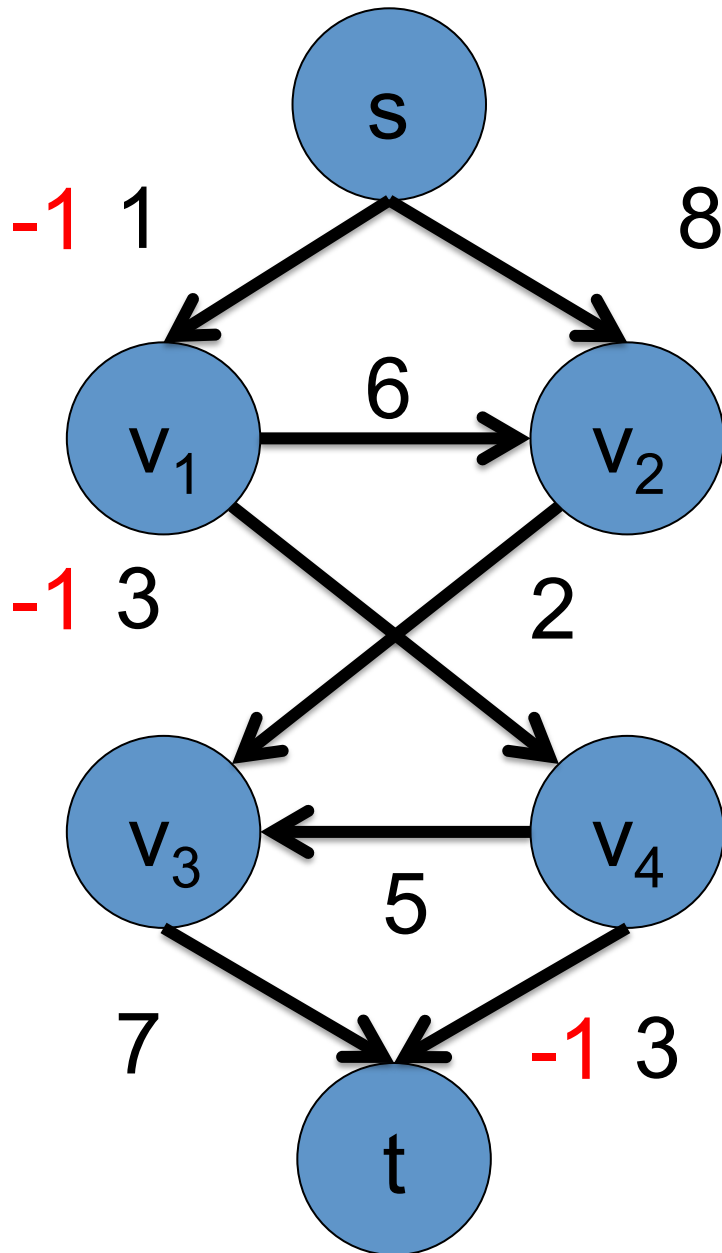
$$\text{flow}(a) \geq 0$$

For all  $v \in V \setminus \{s, t\}$

$$E_{\text{flow}}(v) = 0$$



# s-t Flow



Function flow:  $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

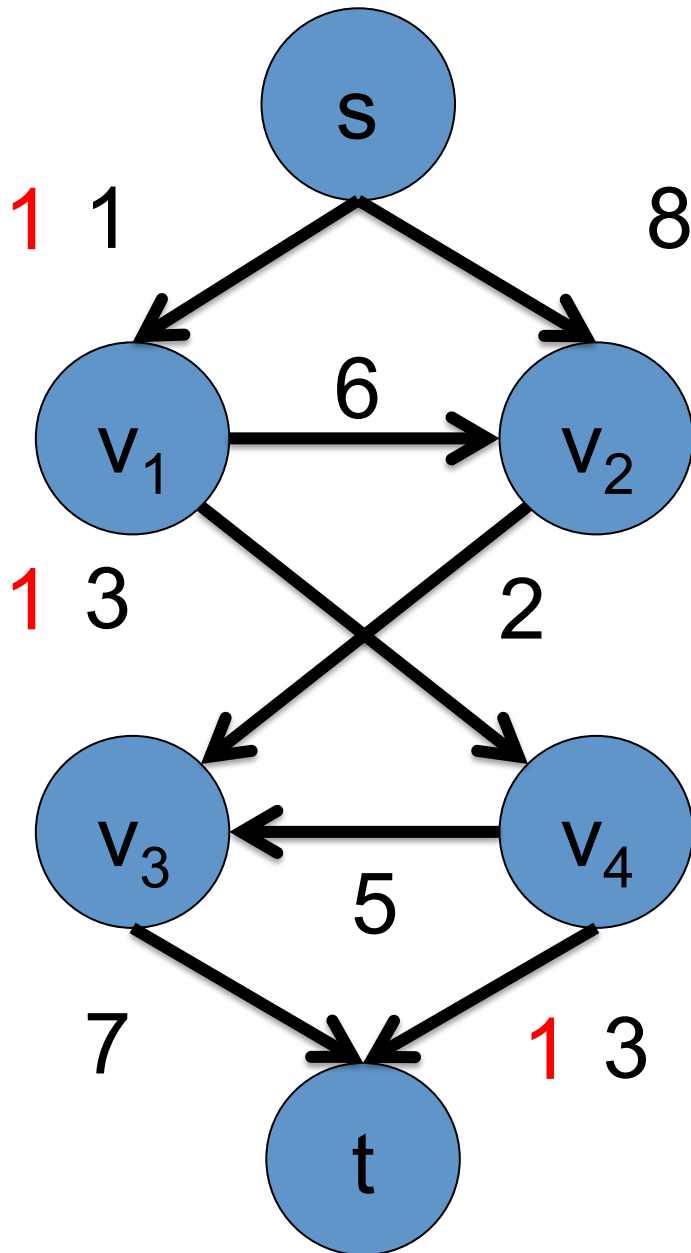
$$\text{flow}(a) \geq 0$$

For all  $v \in V \setminus \{s, t\}$

$$E_{\text{flow}}(v) = 0$$

**X**

# s-t Flow



Function flow:  $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

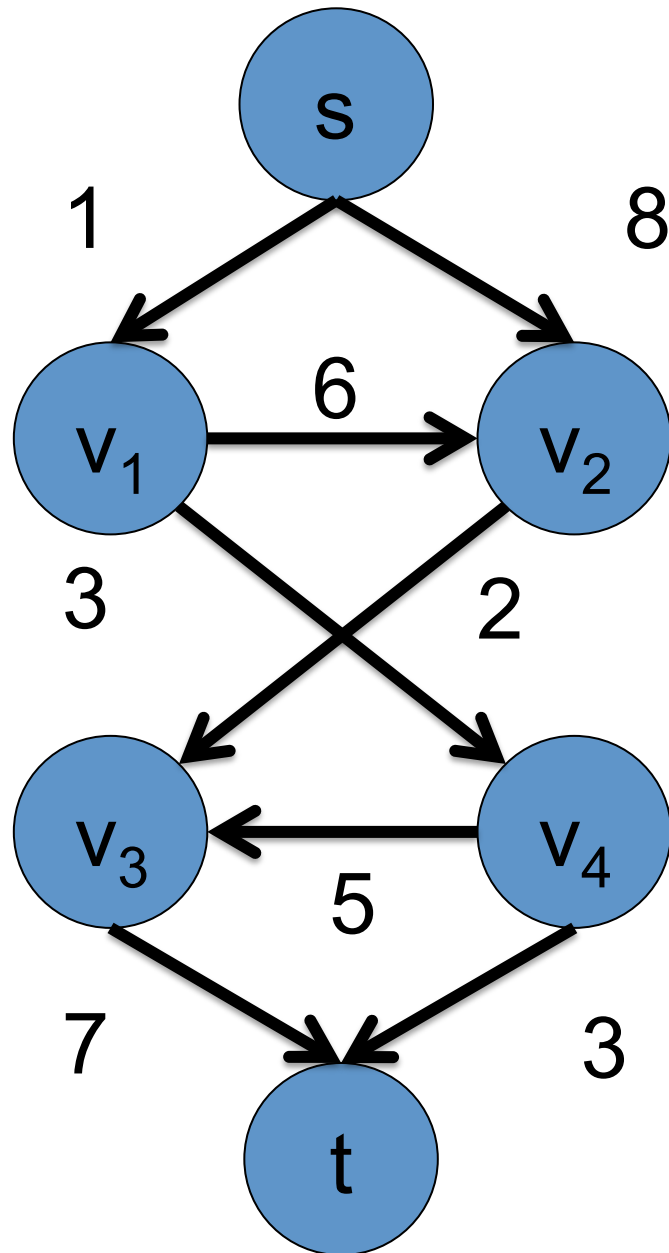
$$\text{flow}(a) \geq 0$$

For all  $v \in V \setminus \{s, t\}$

$$E_{\text{flow}}(v) = 0$$

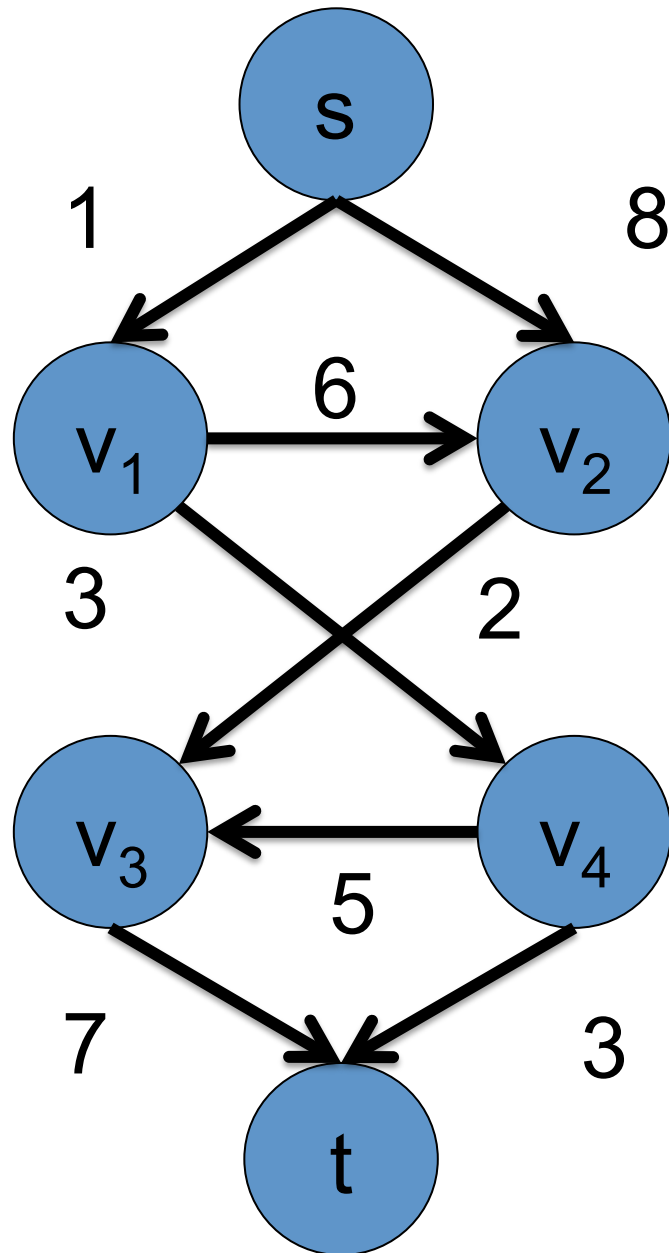


# Value of s-t Flow



Outgoing flow of s  
- Incoming flow of s

# Value of s-t Flow

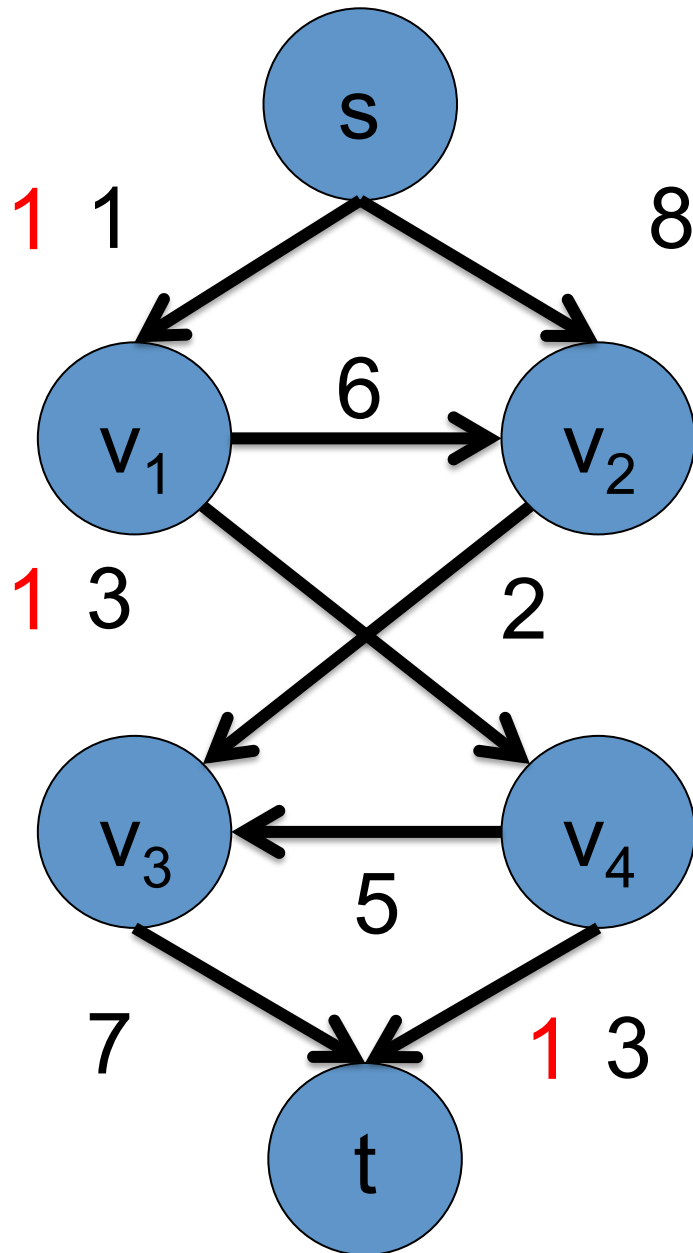


$$-E_{\text{flow}}(s) \quad E_{\text{flow}}(t)$$

$$\sum_{(s,v) \in A} \text{flow}((s,v))$$

$$- \sum_{(u,s) \in A} \text{flow}((u,s))$$

# Value of s-t Flow



$$-E_{\text{flow}}(s) \quad E_{\text{flow}}(t)$$

$$\sum_{(s,v) \in A} \text{flow}((s,v))$$

$$- \sum_{(u,s) \in A} \text{flow}((u,s))$$

Value = 1

# Outline

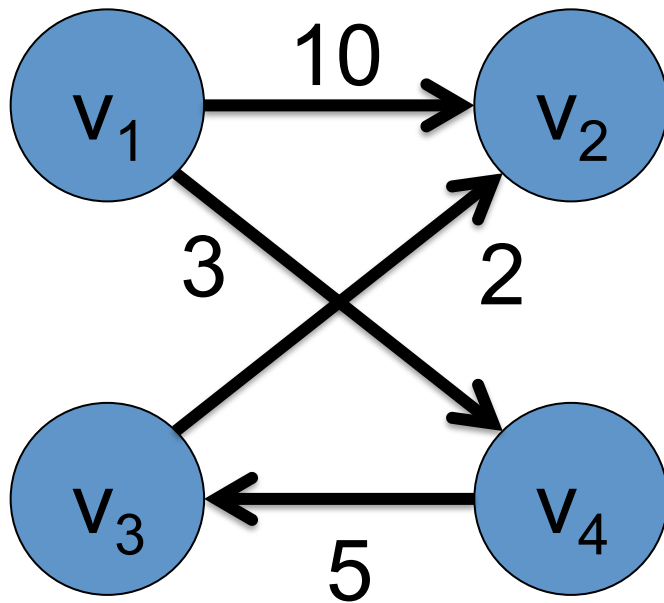
- Preliminaries
  - Functions and Excess Functions
  - s-t Flow
  - **s-t Cut**
  - Flows vs. Cuts
- Maximum Flow
- Algorithms
- Energy minimization with max flow/min cut



# Cut

$$D = (V, A)$$

Let  $U$  be a subset of  $V$



$C$  is a set of arcs such that

- $(u, v) \in A$
- $u \in U$
- $v \in V \setminus U$

$C$  is a cut in the digraph  $D$

# Cut

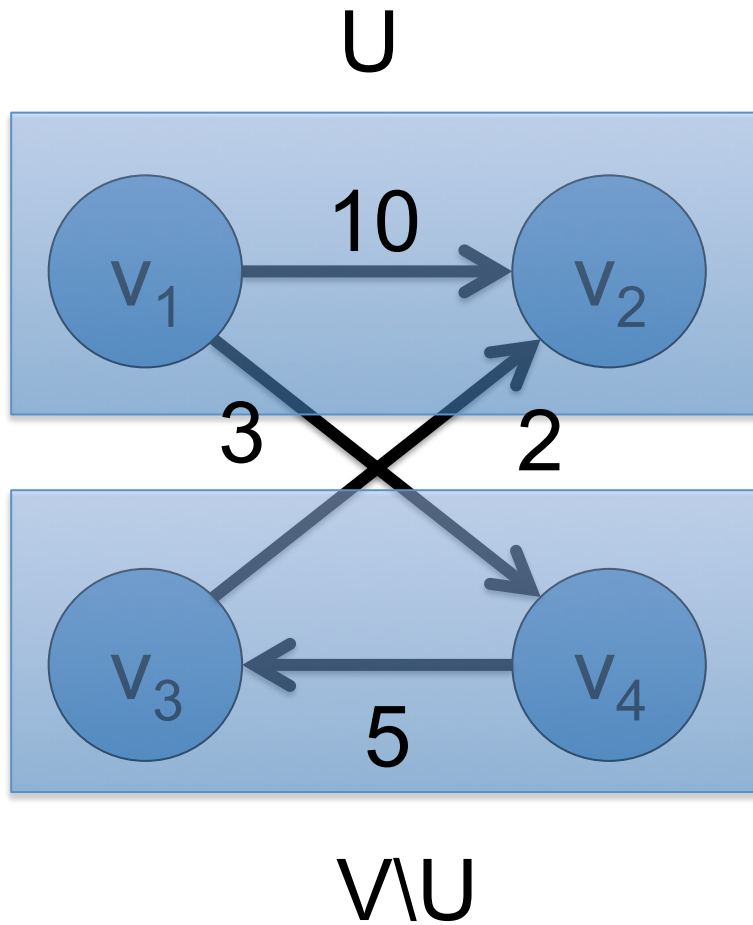
$$D = (V, A)$$

What is C?

$$\{(v_1, v_2), (v_1, v_4)\} ?$$

$$\{(v_1, v_4), (v_3, v_2)\} ?$$

$$\{(v_1, v_4)\} ?$$



# Cut

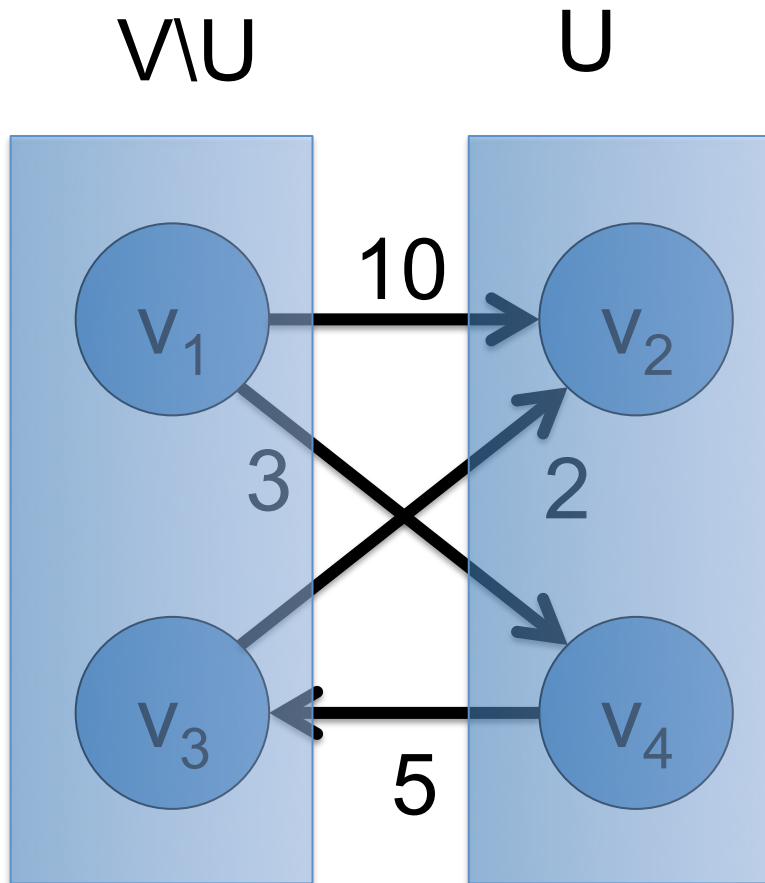
$$D = (V, A)$$

What is C?

$$\{(v_1, v_2), (v_1, v_4), (v_3, v_2)\} ?$$

✓  $\{(v_4, v_3)\} ?$

$$\{(v_1, v_4), (v_3, v_2)\} ?$$



# Cut

$$D = (V, A)$$

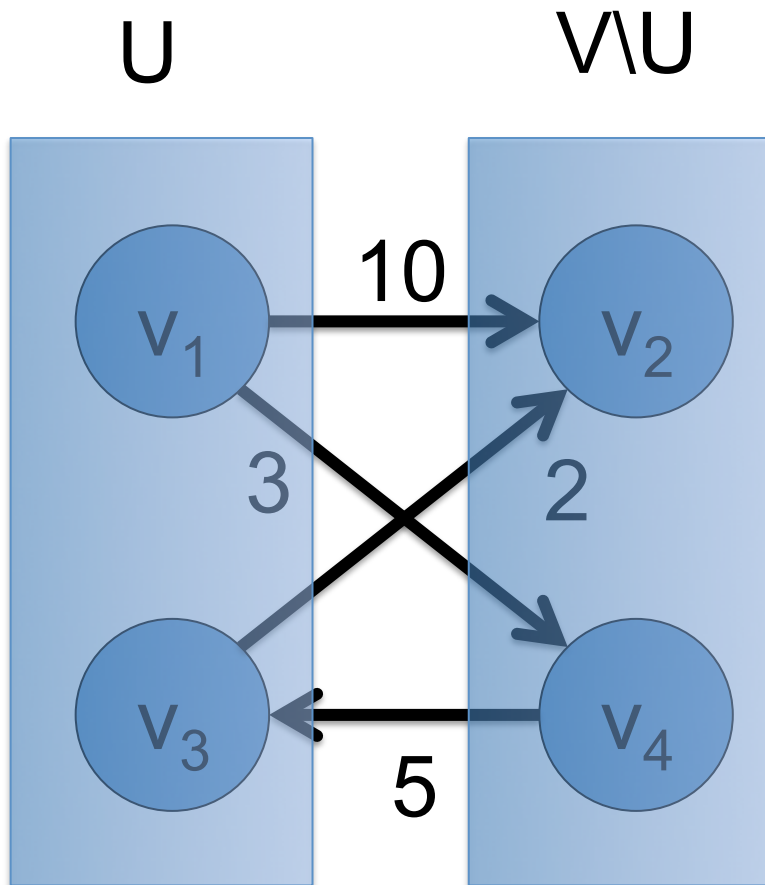
What is C?



$$\{(v_1, v_2), (v_1, v_4), (v_3, v_2)\} ?$$

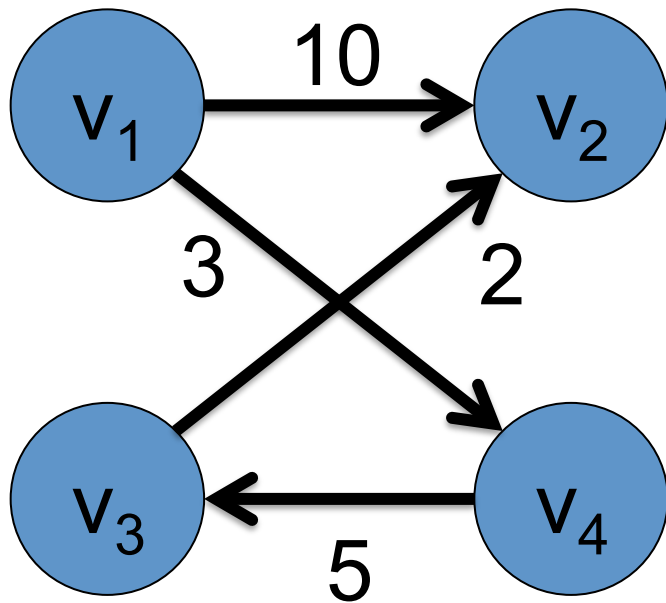
$$\{(v_3, v_2)\} ?$$

$$\{(v_1, v_4), (v_3, v_2)\} ?$$



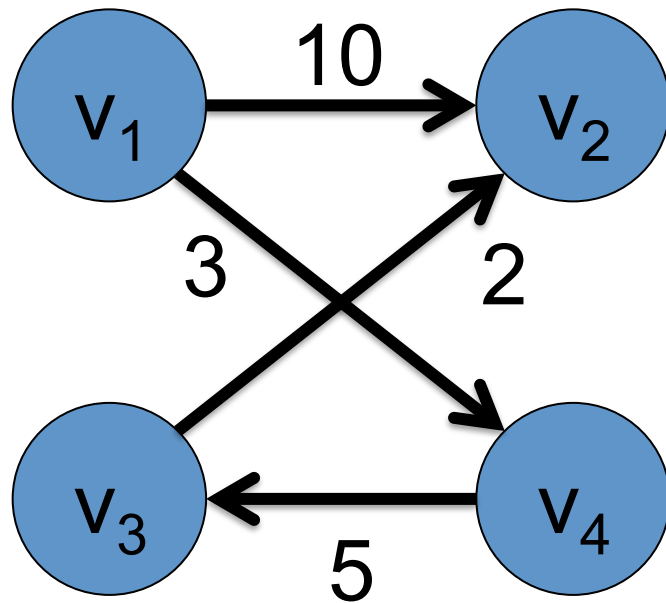
# Cut

$$D = (V, A)$$



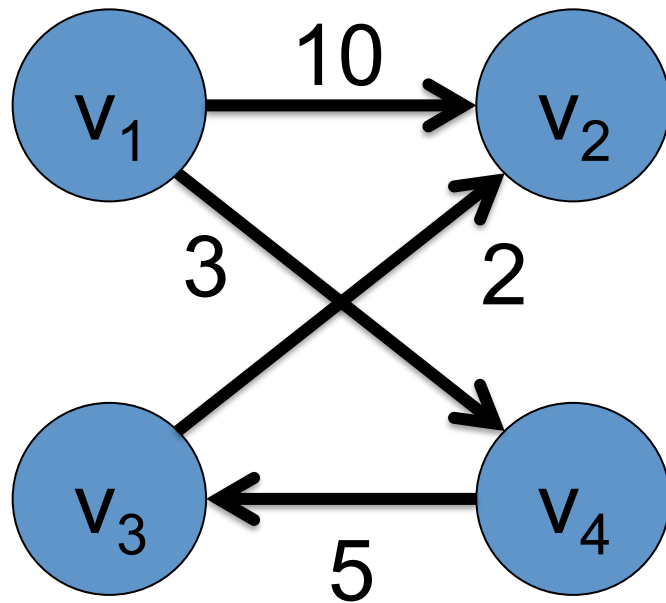
$$C = \text{out-arcs}(U)$$

# Capacity of Cut



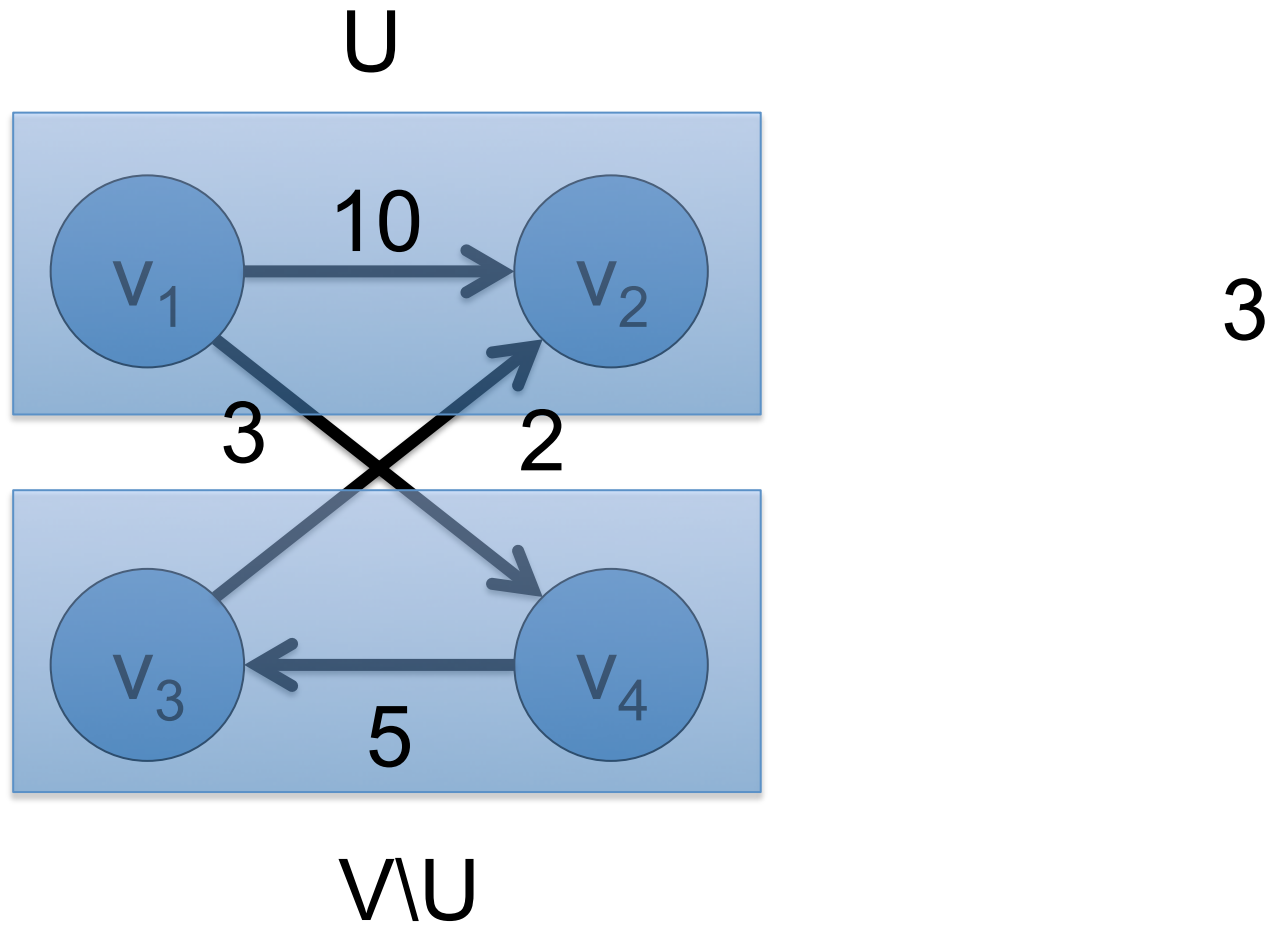
Sum of capacity of all arcs in C

# Capacity of Cut



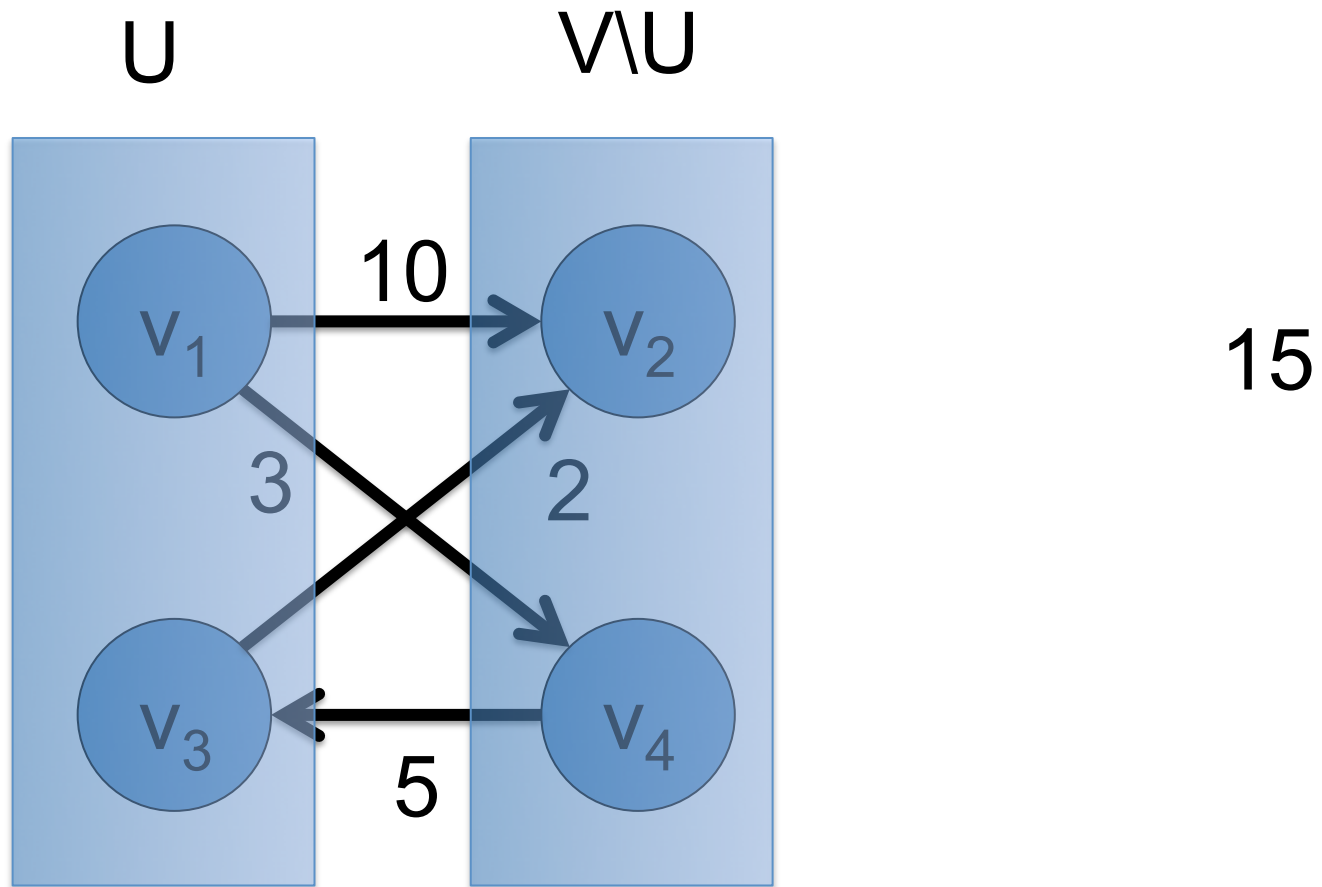
$$\sum_{a \in C} c(a)$$

# Capacity of Cut



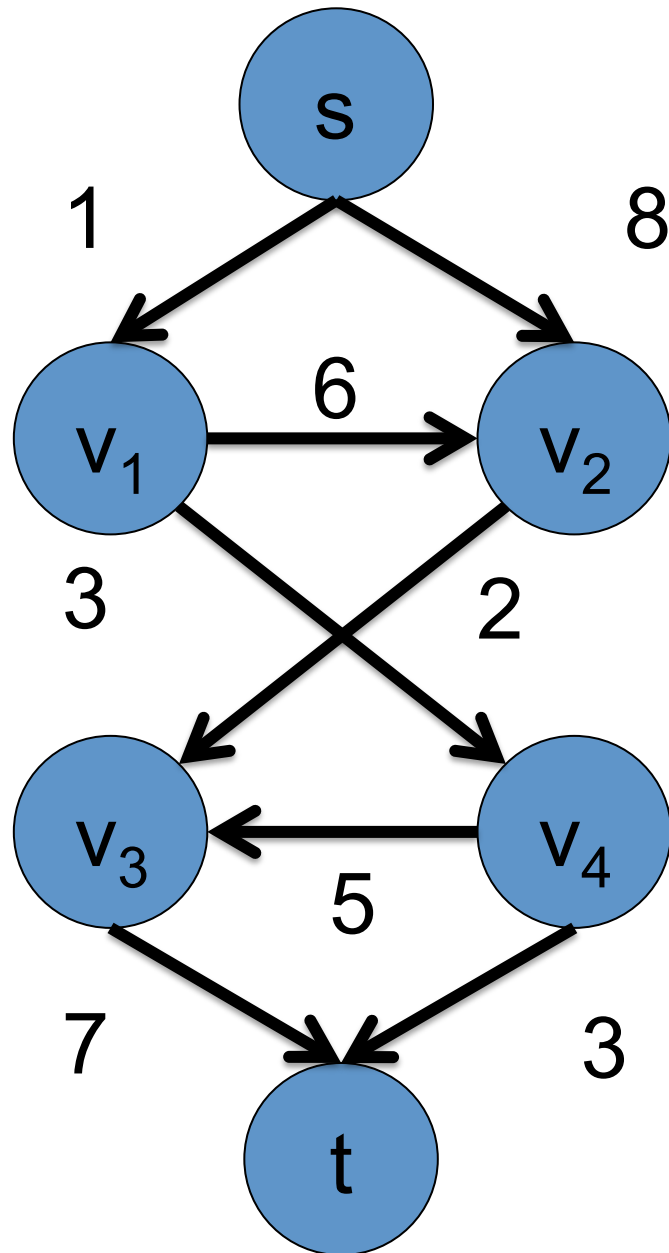


# Capacity of Cut



# s-t Cut

$$D = (V, A)$$



A source vertex “s”

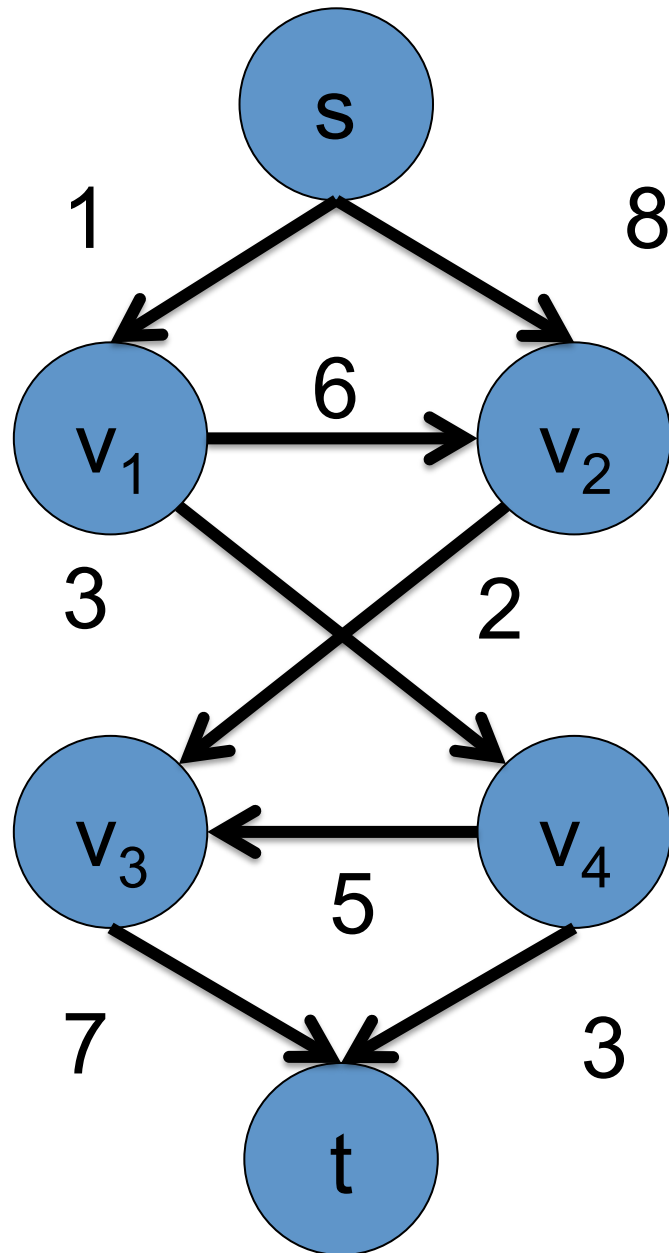
A sink vertex “t”

C is a cut such that

- $s \in U$
- $t \in V \setminus U$

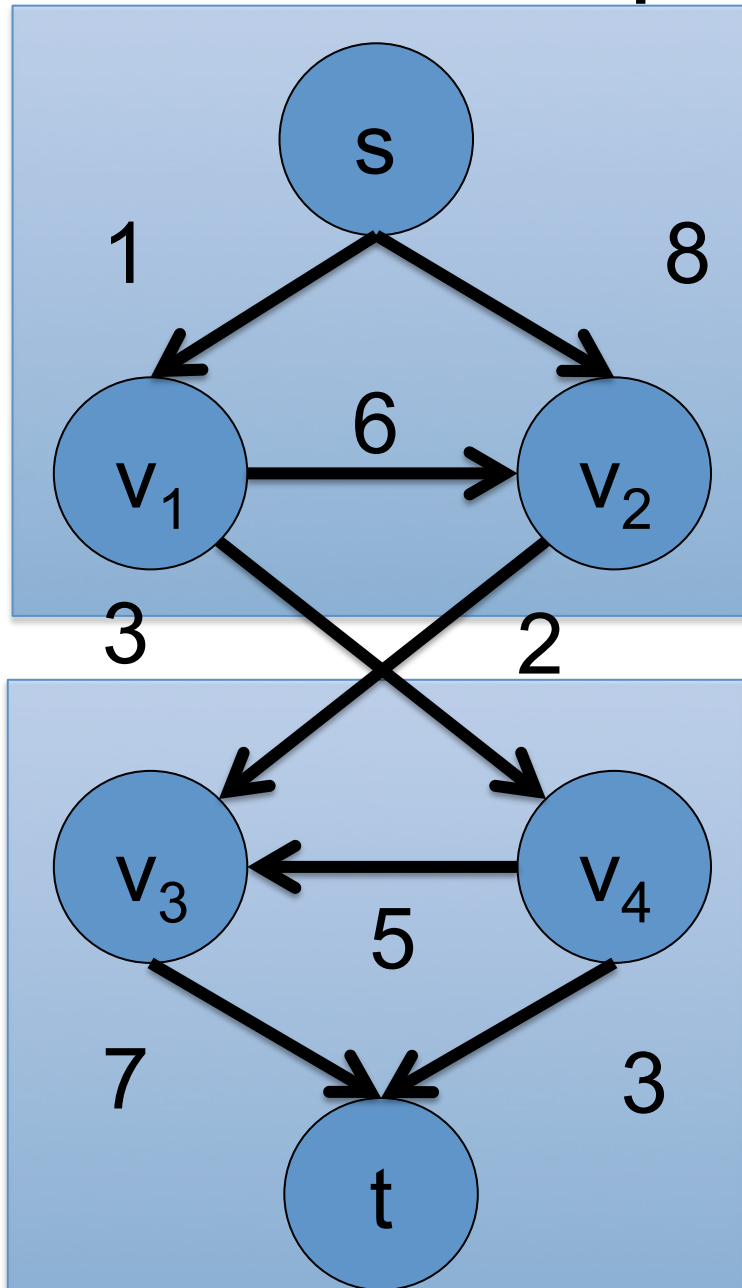
C is an s-t cut

# Capacity of s-t Cut



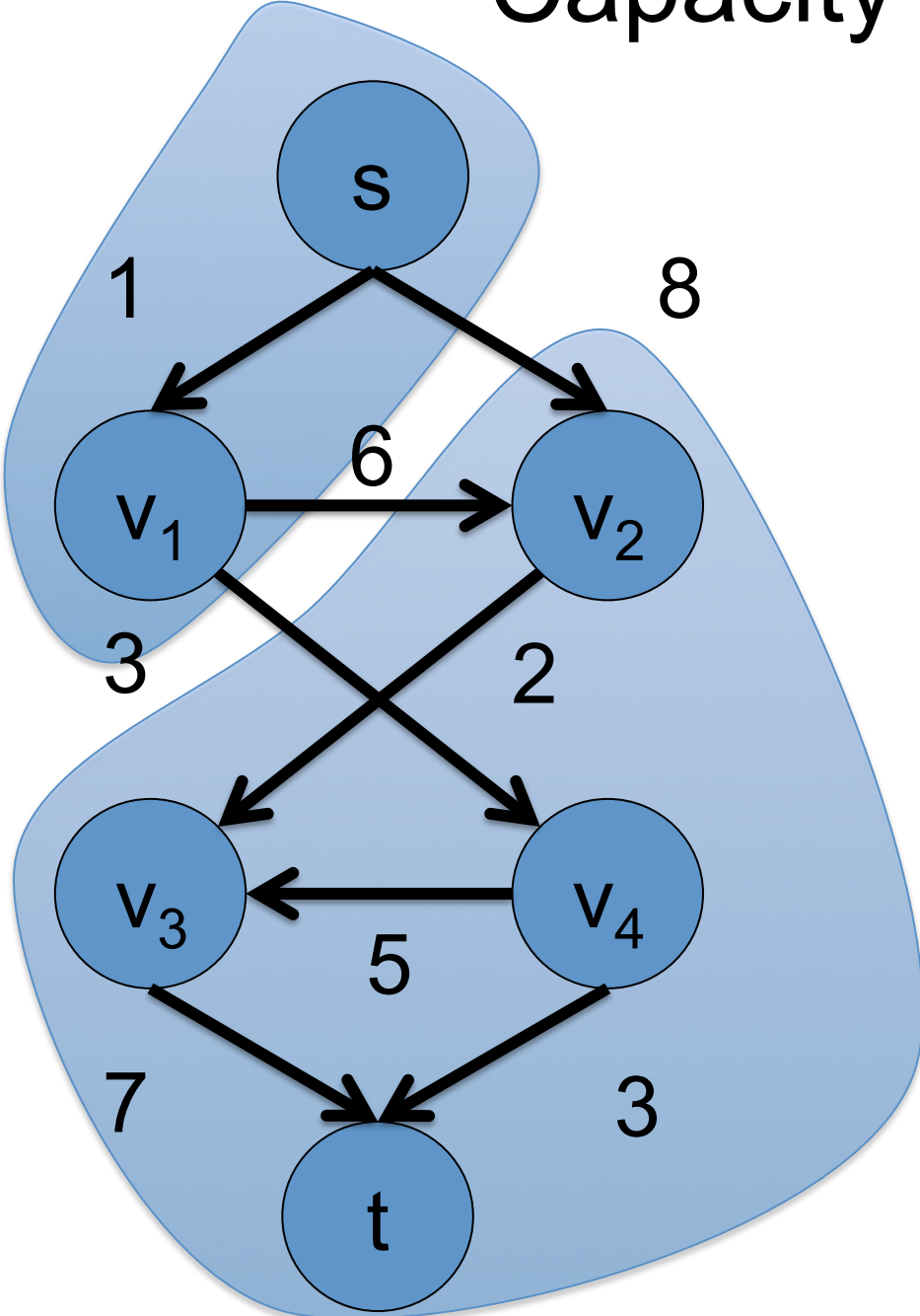
$$\sum_{a \in C} c(a)$$

# Capacity of s-t Cut



5

# Capacity of s-t Cut



17

# Outline

- Preliminaries
  - s-t Flow
  - s-t Cut
  - **Flows vs. Cuts**
- Maximum Flow
- Algorithms
- Energy minimization with max flow/min cut

# Flows vs. Cuts

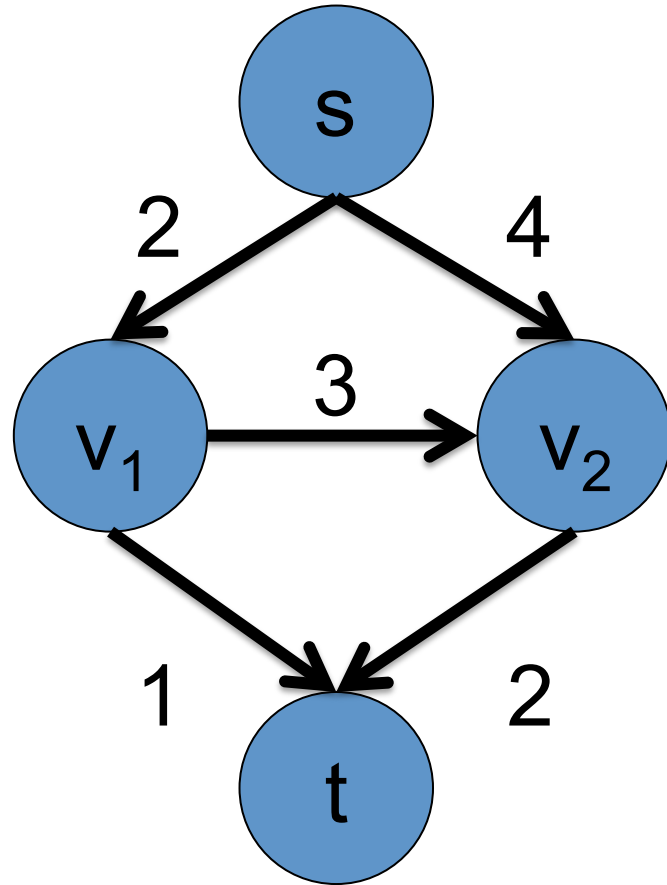
(See next lecture)

# Outline

- Preliminaries
- **Maximum Flow**
  - Residual Graph
  - Max-Flow Min-Cut Theorem
- Algorithms
- Energy minimization with max flow/min cut



# Maximum Flow Problem



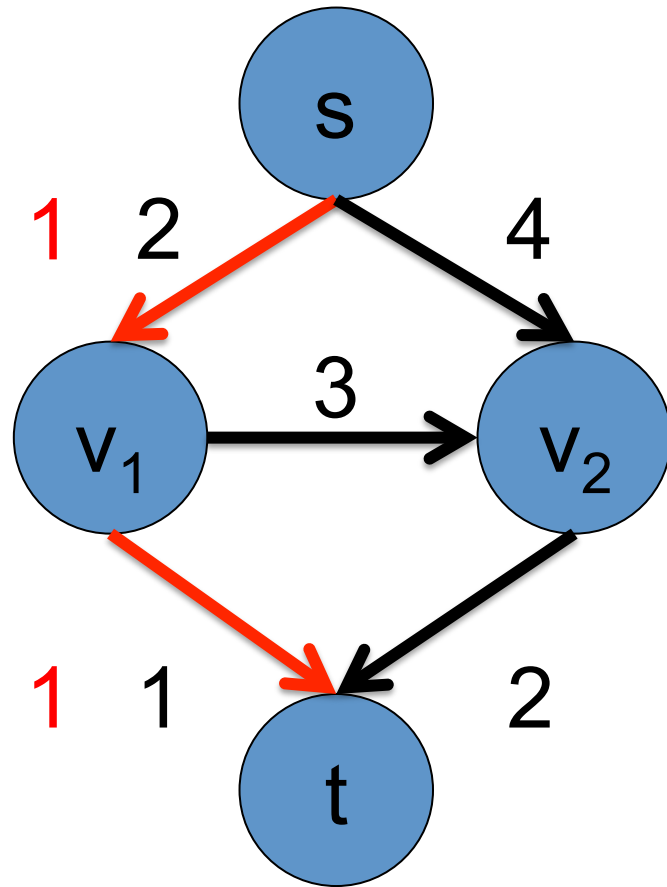
Find the flow with the maximum value !!

$$\sum_{(s,v) \in A} \text{flow}((s,v))$$

$$- \sum_{(u,s) \in A} \text{flow}((u,s))$$

**First suggestion to solve this problem !!**

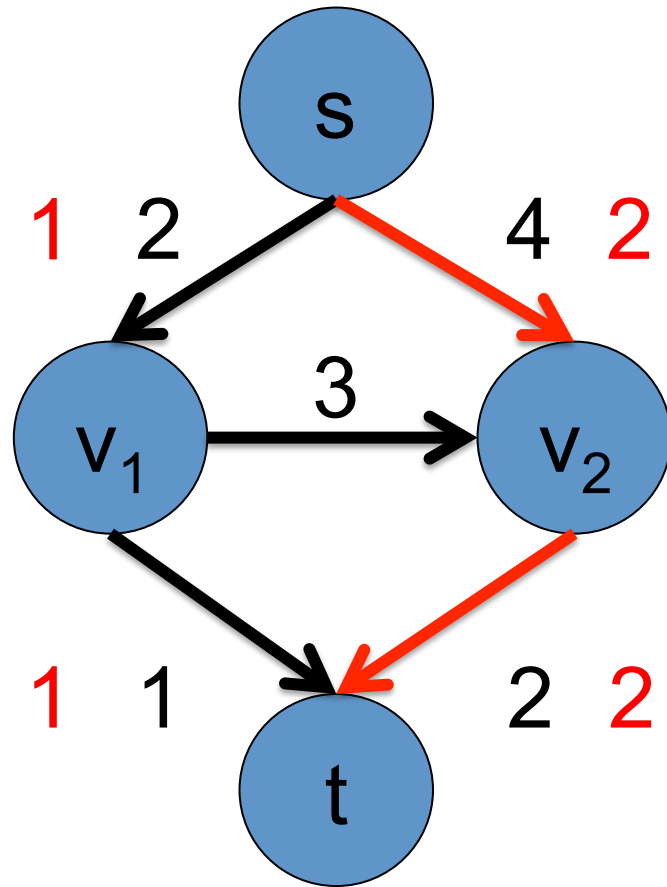
# Passing Flow through s-t Paths



Find an s-t path where  
 $\text{flow}(a) < c(a)$  for all arcs

Pass maximum allowable  
flow through the arcs

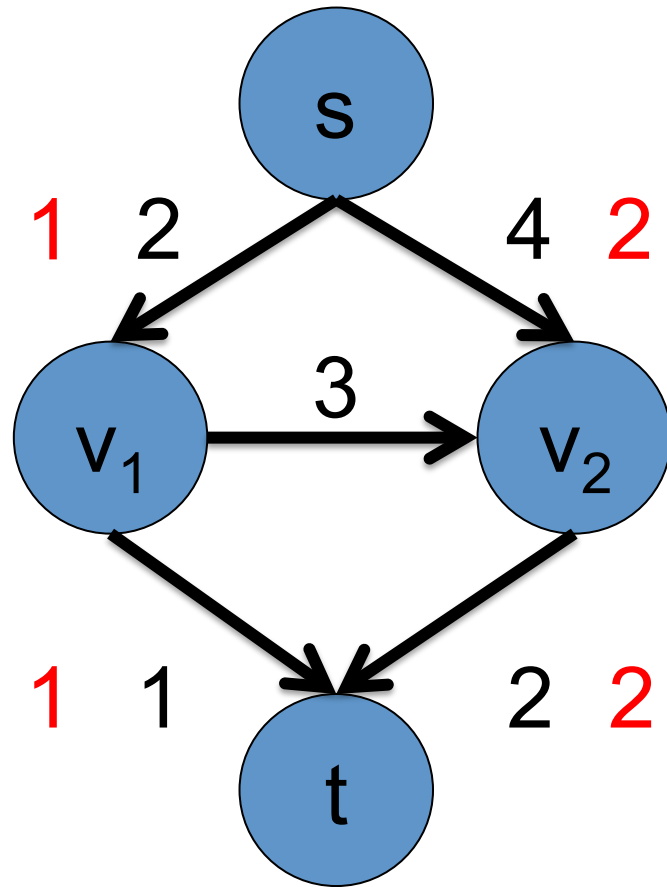
# Passing Flow through s-t Paths



Find an s-t path where  
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# Passing Flow through s-t Paths



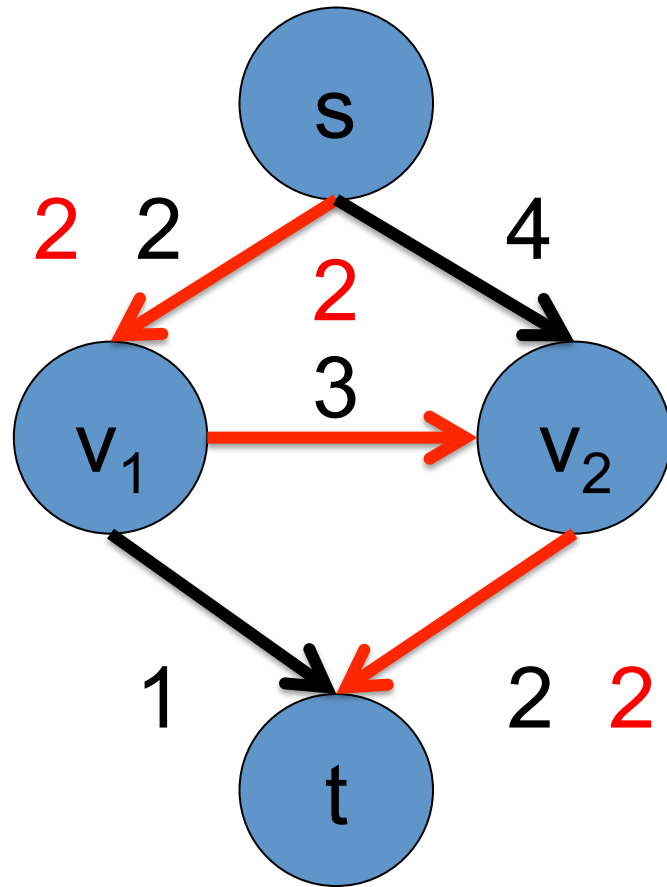
Find an s-t path where  
 $\text{flow}(a) < c(a)$  for all arcs

No more paths. Stop.

Will this give us maximum flow?

**NO !!!**

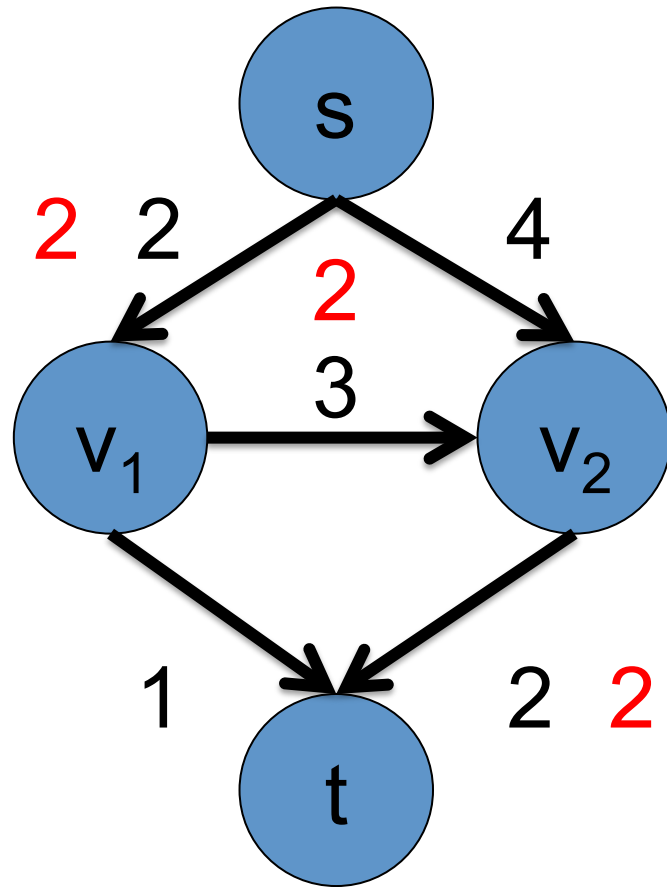
# Passing Flow through s-t Paths



Find an s-t path where  
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# Passing Flow through s-t Paths



Find an s-t path where  
 $\text{flow}(a) < c(a)$  for all arcs

No more paths. Stop.

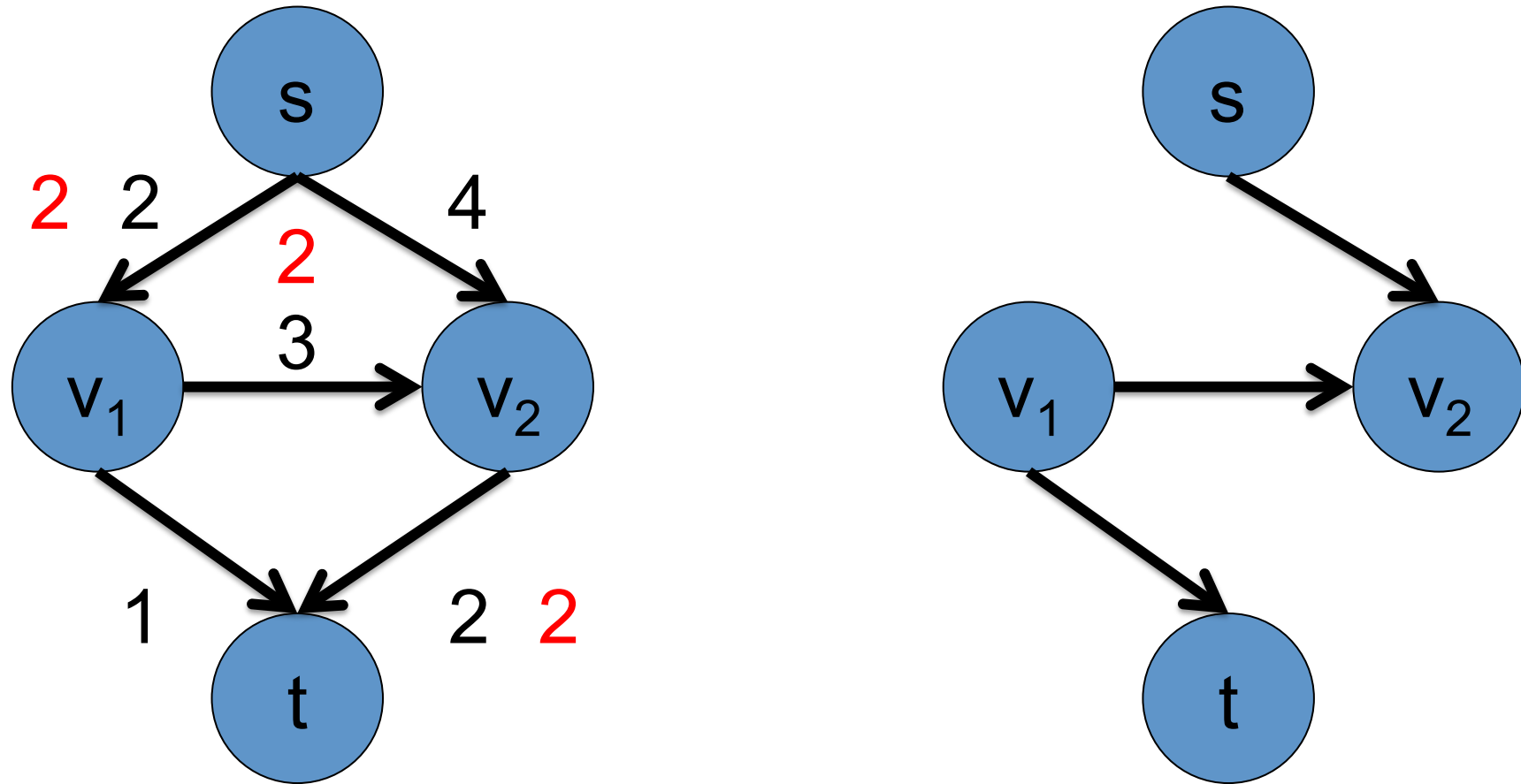
**Another method?**

**Incorrect Answer !!**

# Outline

- Preliminaries
- Maximum Flow
  - **Residual Graph**
  - Max-Flow Min-Cut Theorem
- Algorithms
- Energy minimization with max flow/min cut

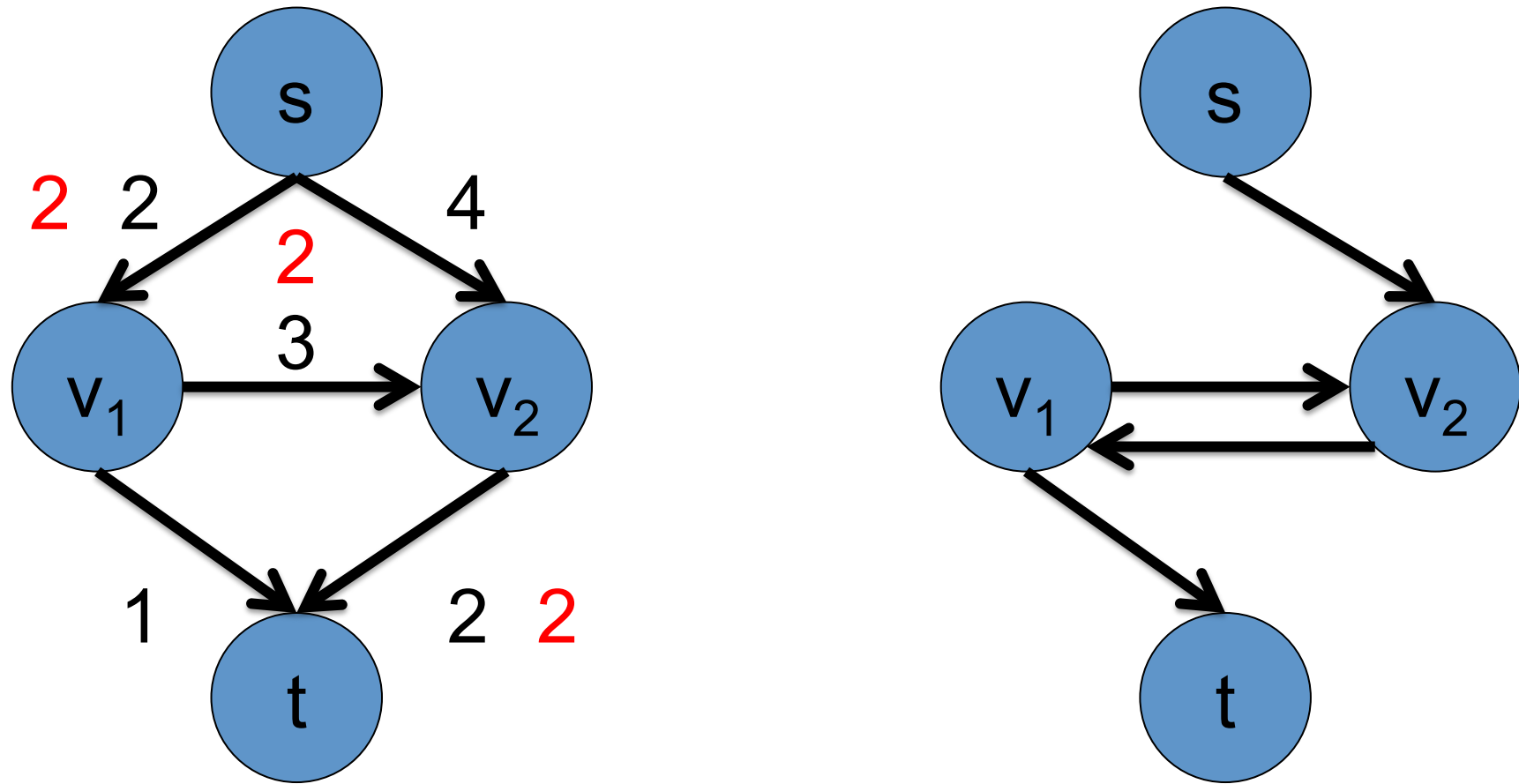
# Residual Graph



Arcs where  $\text{flow}(a) < c(a)$



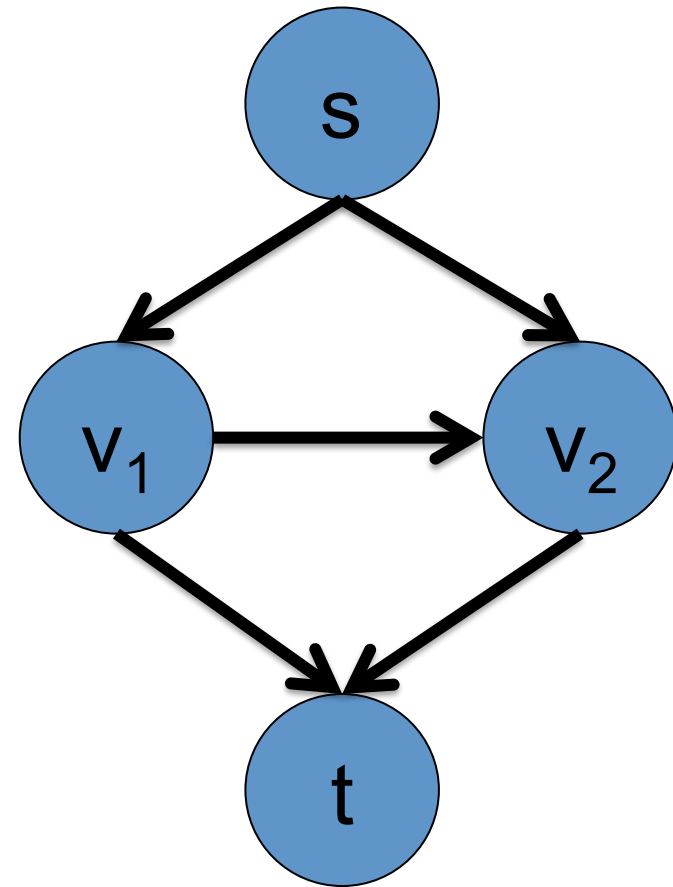
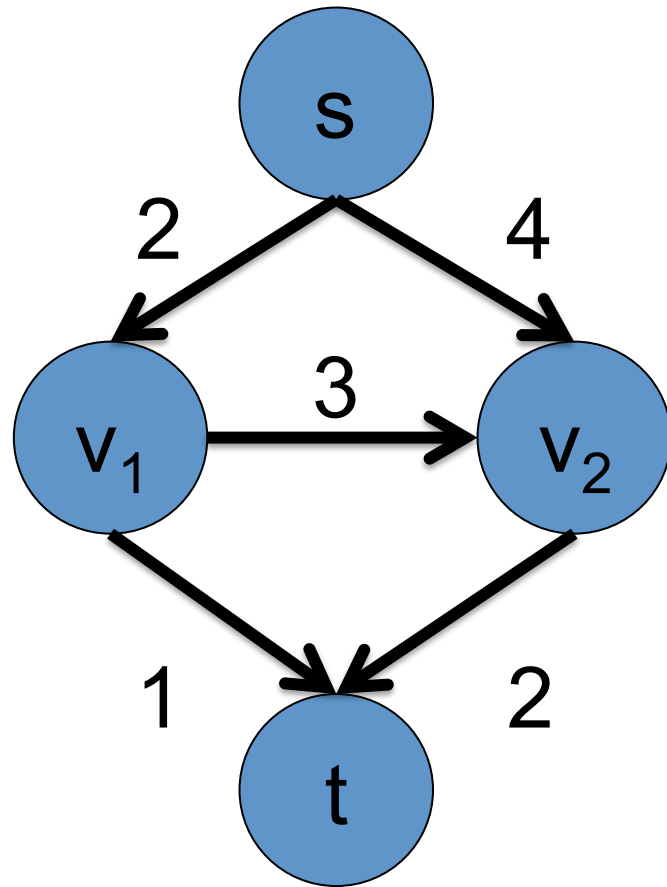
# Residual Graph



Including arcs to  $s$  and from  $t$  is not necessary

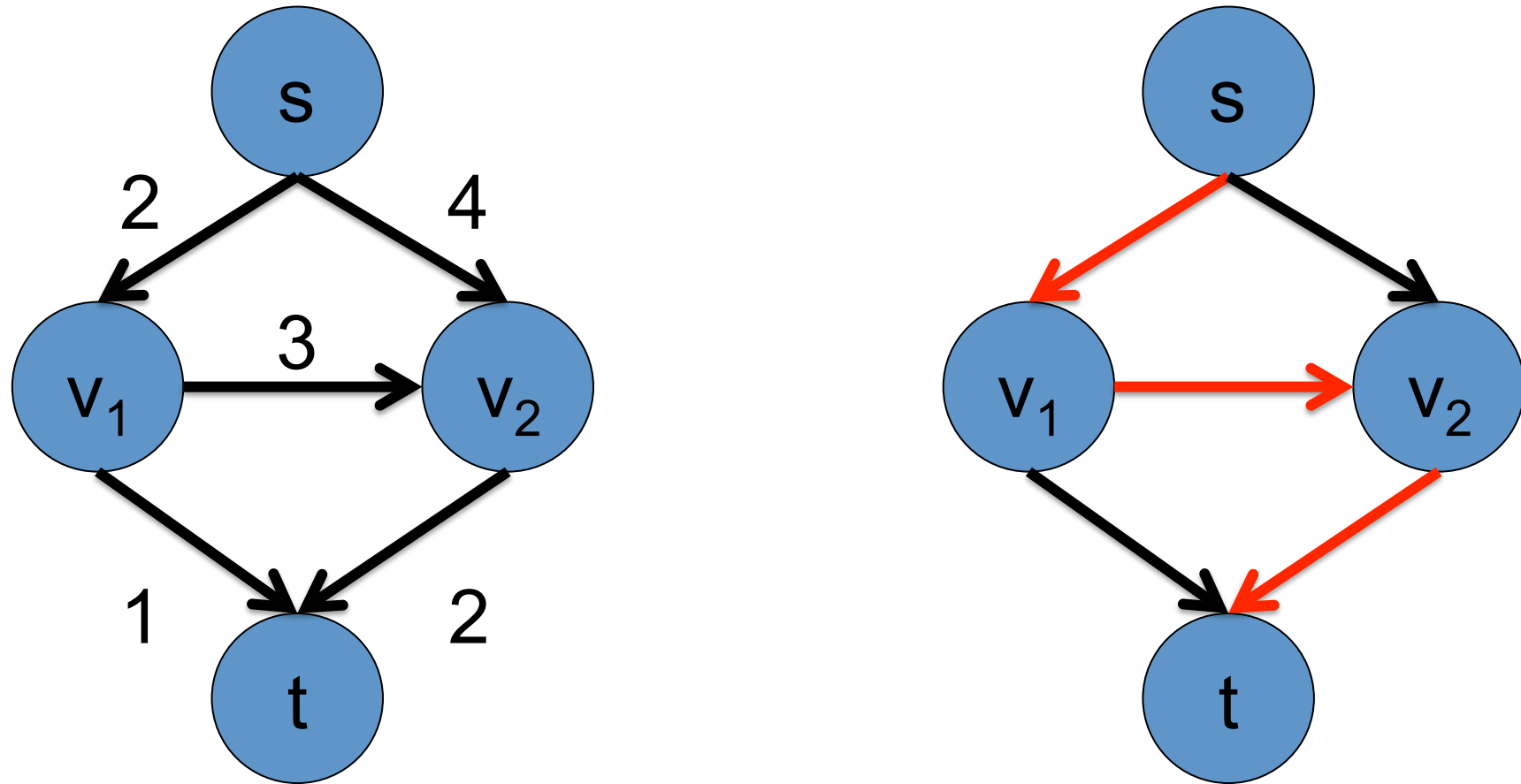
Inverse of arcs where  $\text{flow}(a) > 0$

# Maximum Flow using Residual Graphs



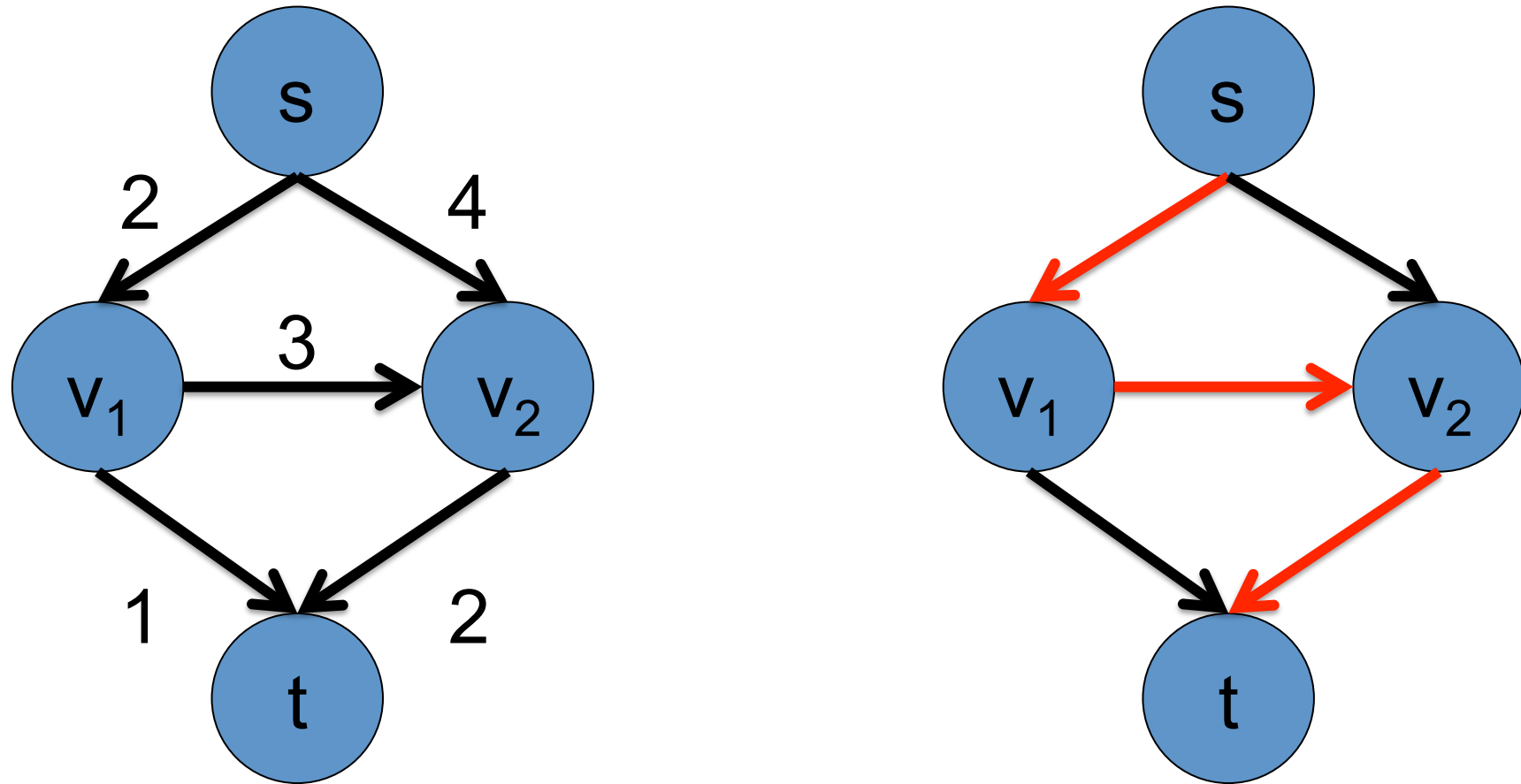
Start with zero flow.

# Maximum Flow using Residual Graphs



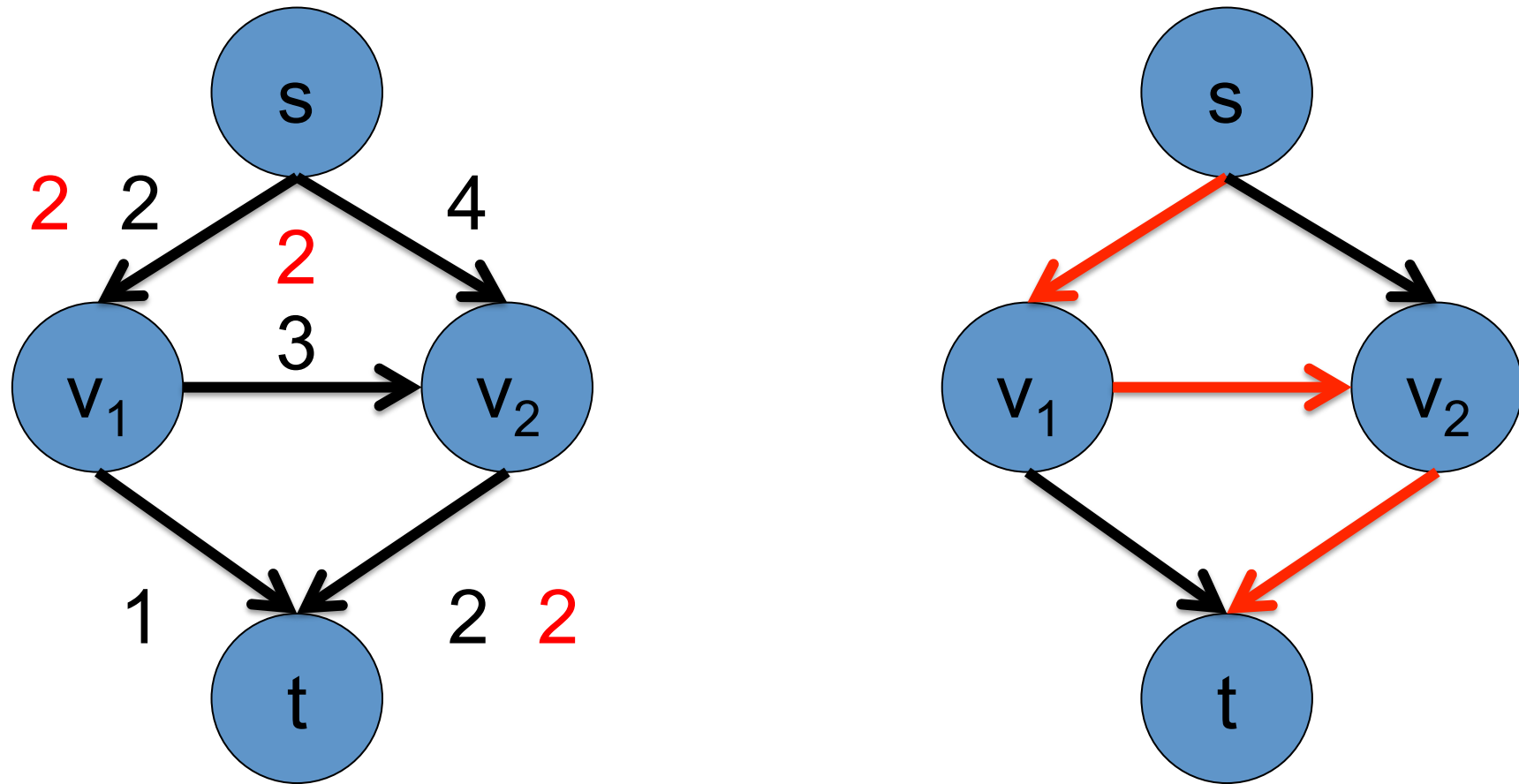
Find an  $s$ - $t$  path in the residual graph.

# Maximum Flow using Residual Graphs



For inverse arcs in path, subtract flow  $K$ .

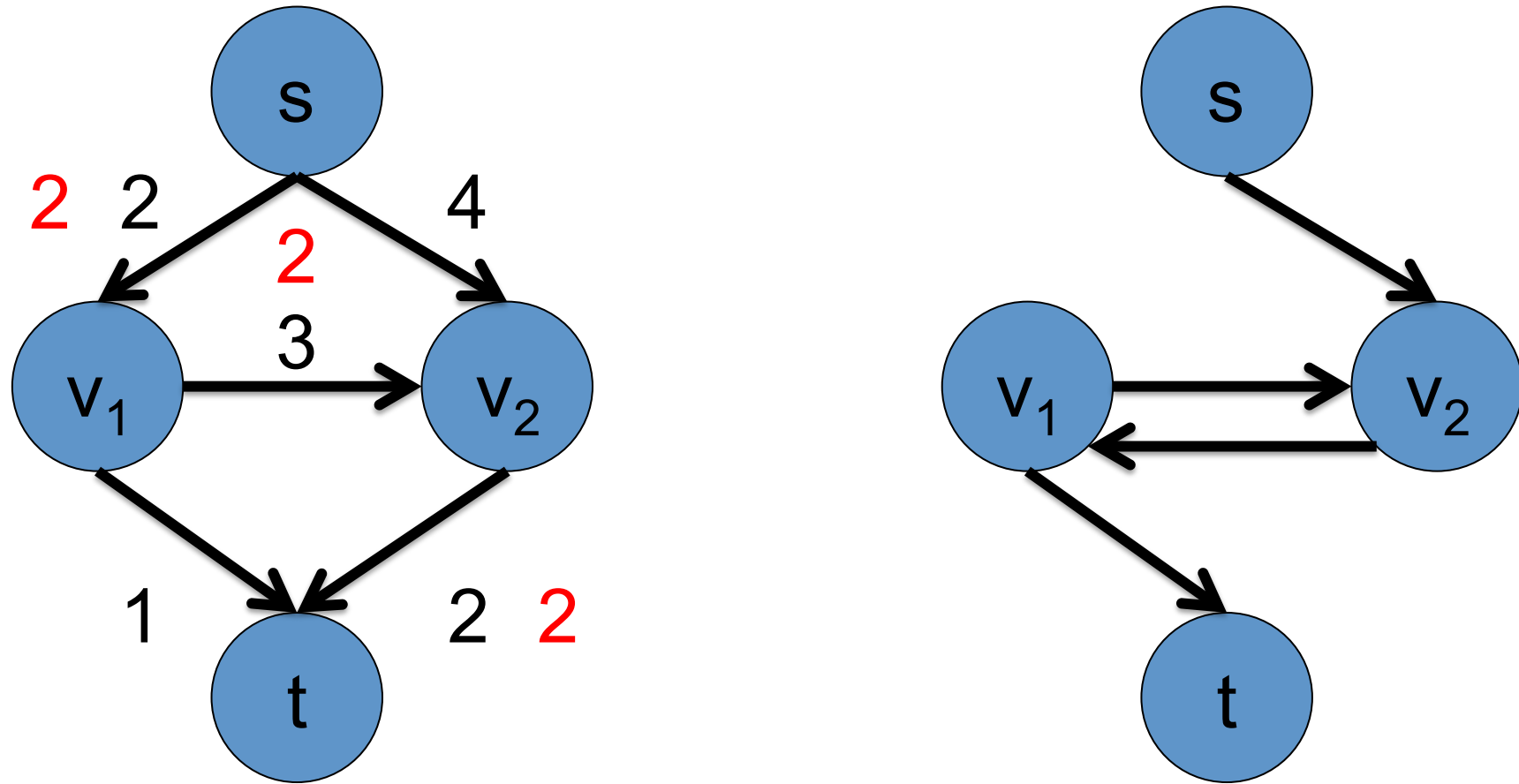
# Maximum Flow using Residual Graphs



Choose maximum allowable value of  $K$ .

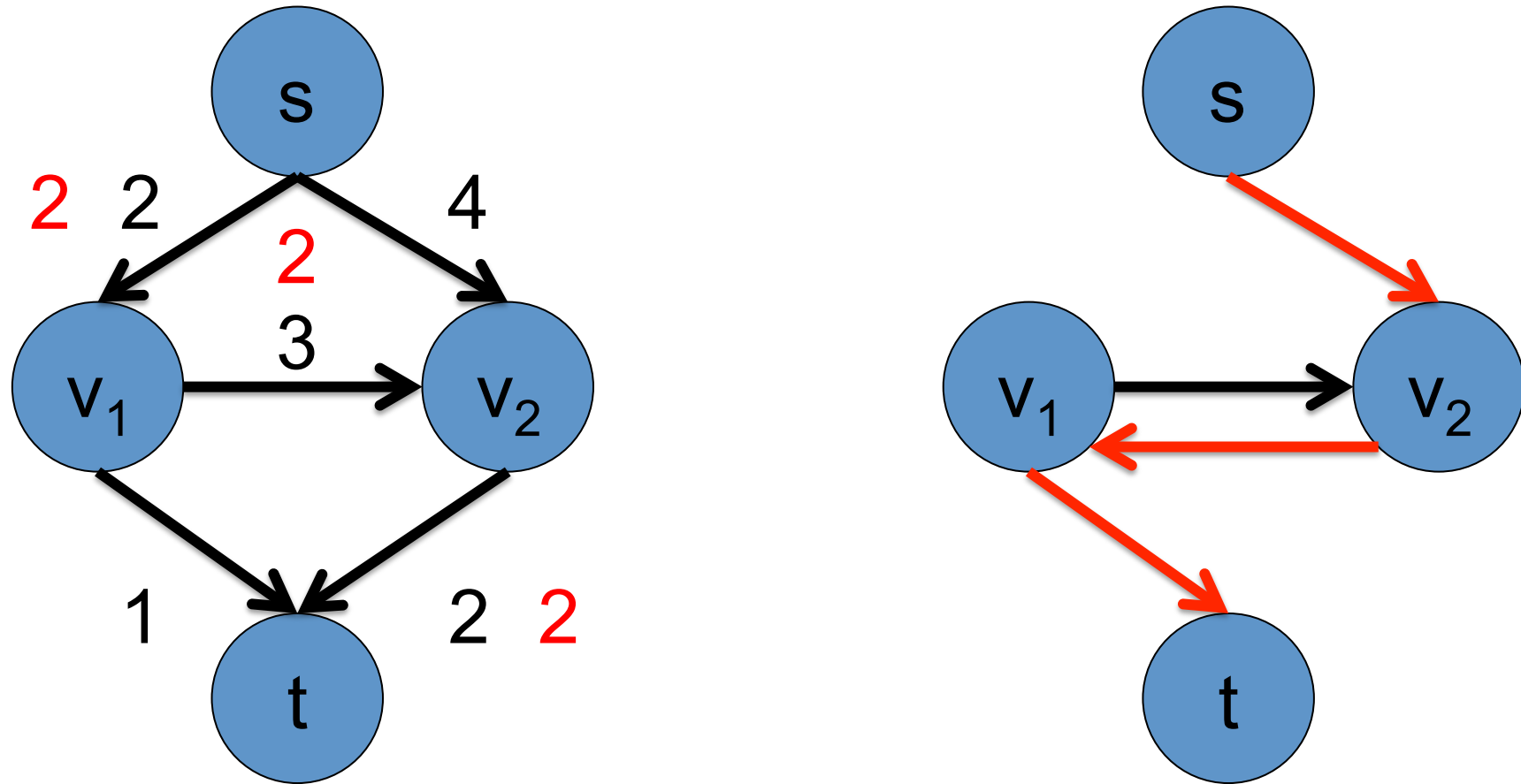
For forward arcs in path, add flow  $K$ .

# Maximum Flow using Residual Graphs



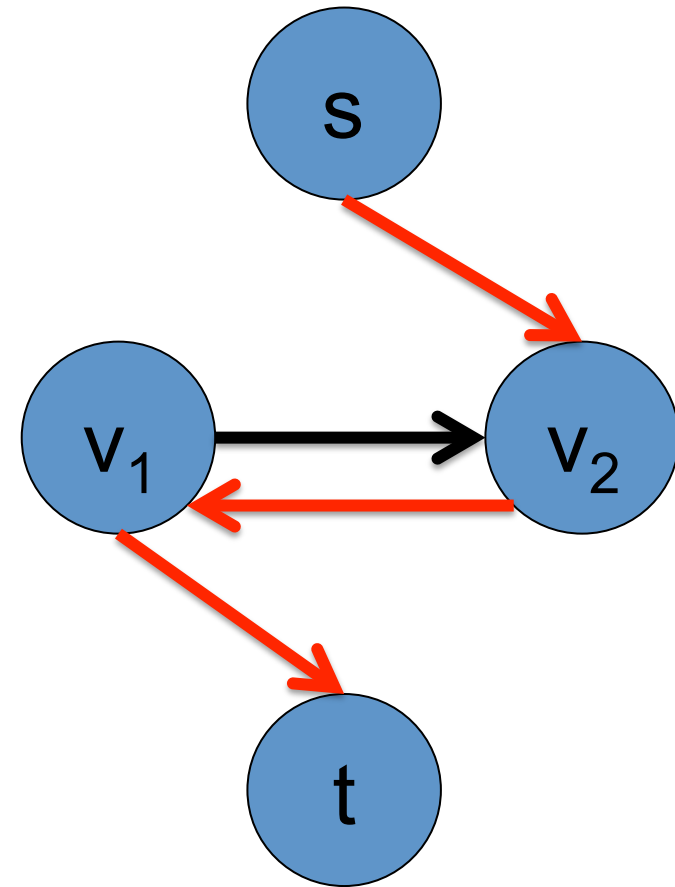
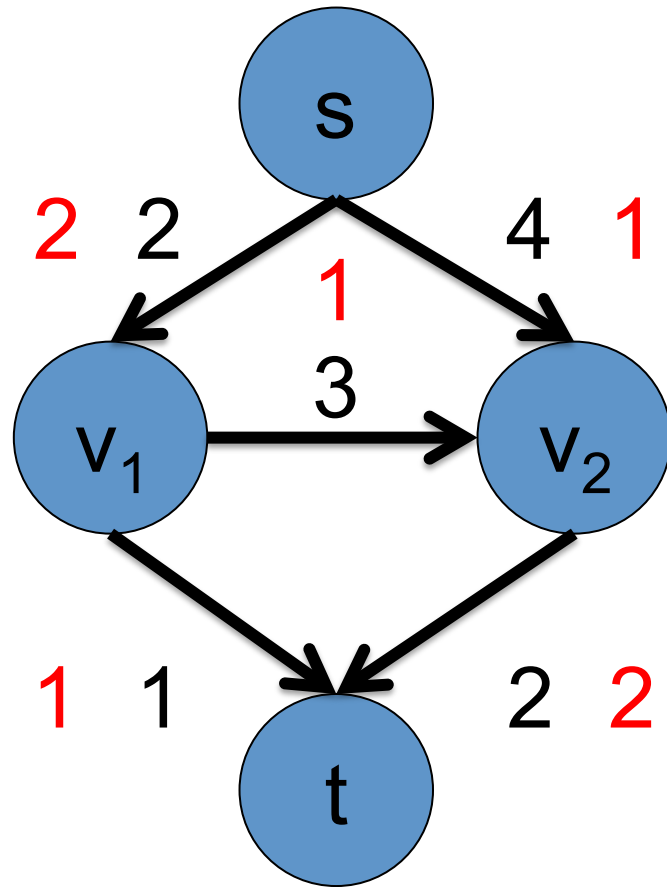
Update the residual graph.

# Maximum Flow using Residual Graphs



Find an s-t path in the residual graph.

# Maximum Flow using Residual Graphs

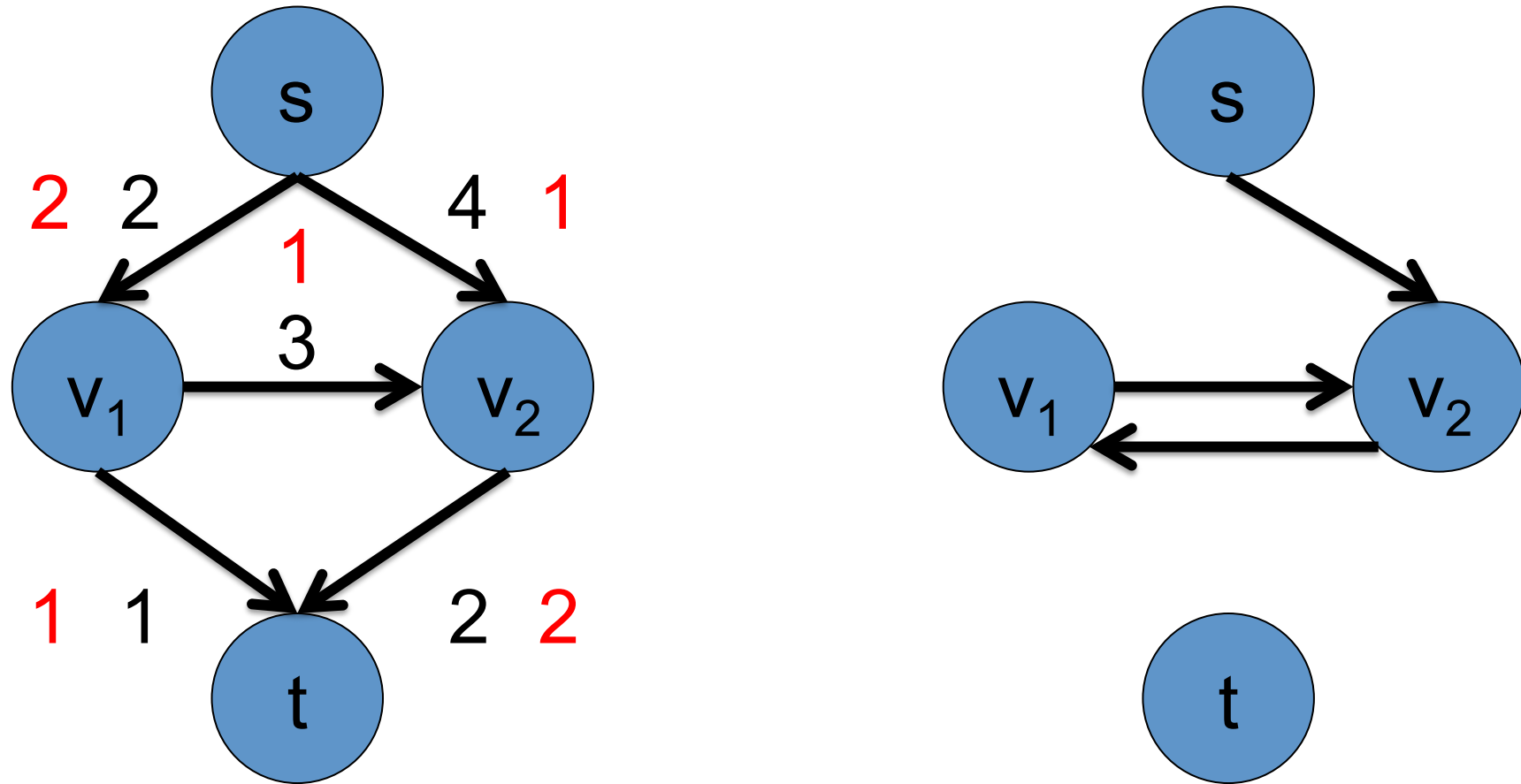


Choose maximum allowable value of  $K$ .

Add  $K$  to  $(s, v_2)$  and  $(v_1, t)$ . Subtract  $K$  from  $(v_1, v_2)$ .

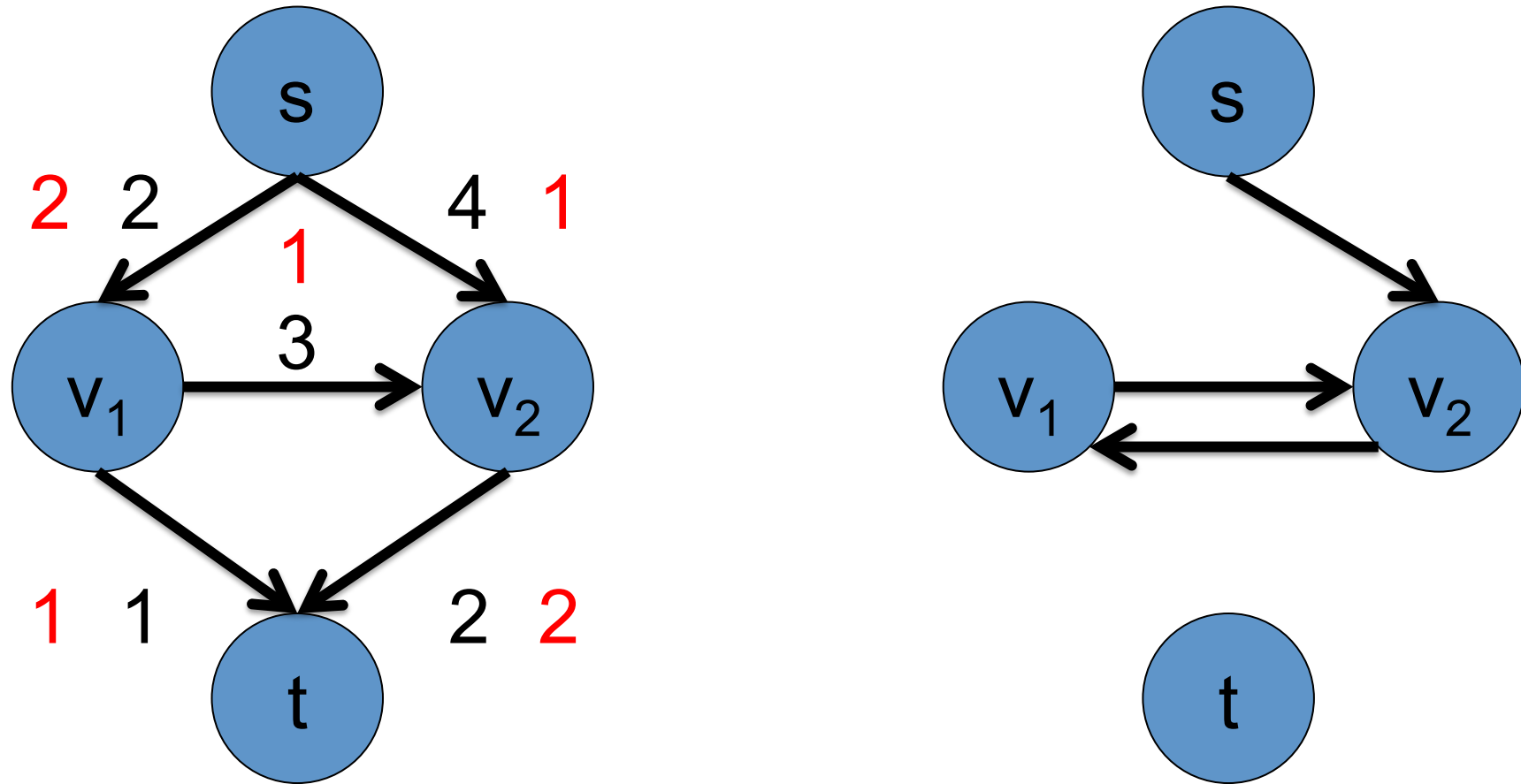


# Maximum Flow using Residual Graphs



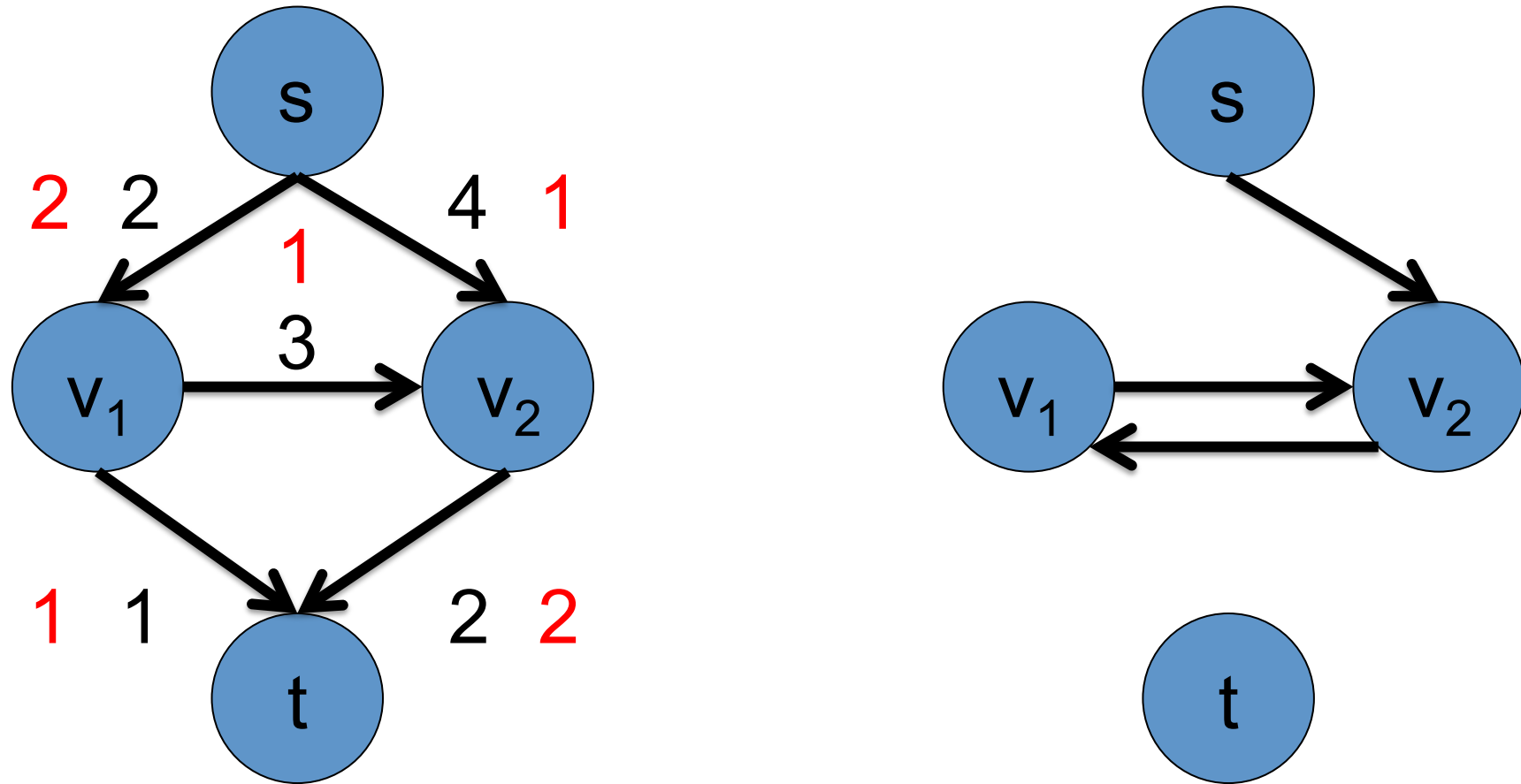
Update the residual graph.

# Maximum Flow using Residual Graphs



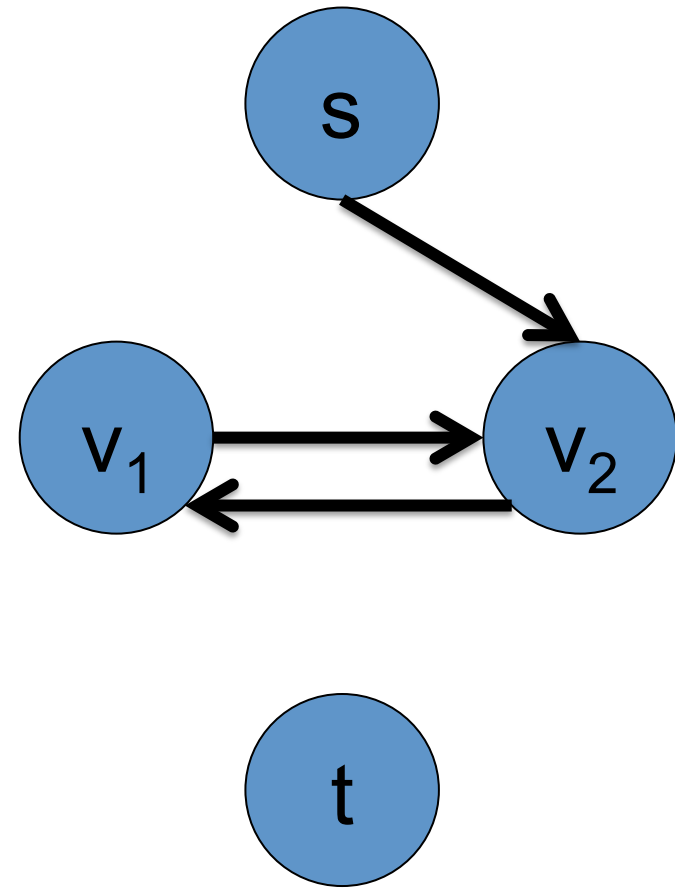
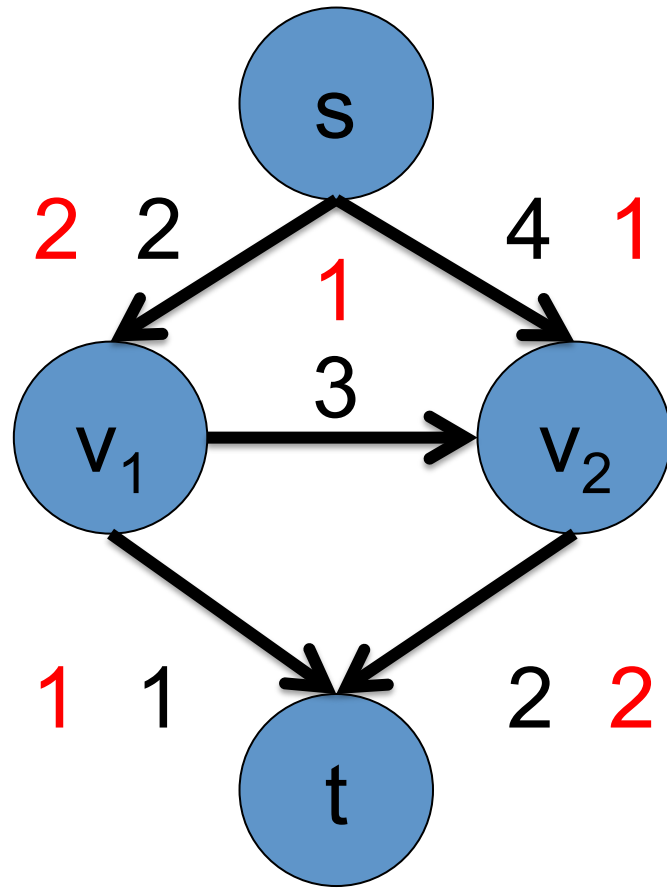
Find an s-t path in the residual graph.

# Maximum Flow using Residual Graphs



No more s-t paths. Stop.

# Maximum Flow using Residual Graphs



Correct Answer.