

Graphical Models

Discrete Inference and Learning

Lecture 3 (contd.)

MVA

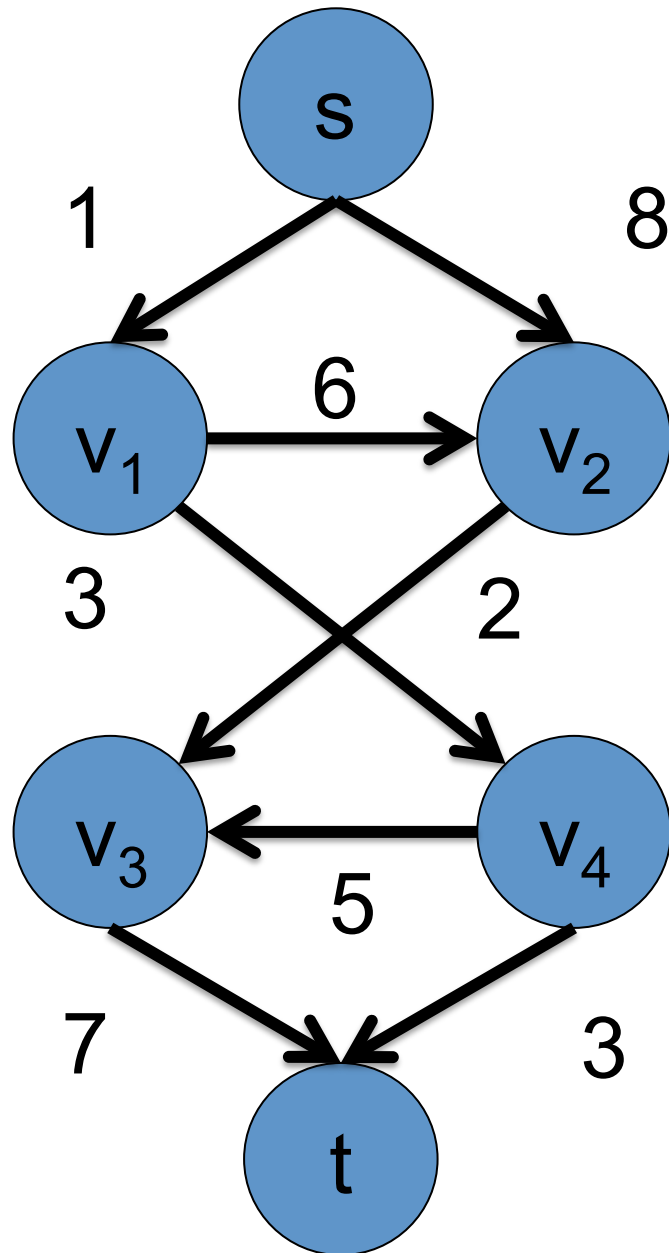
2020 – 2021

<http://thoth.inrialpes.fr/~alahari/disinflern>

Outline

- Preliminaries
 - **s-t Flow**
 - s-t Cut
 - Flows vs. Cuts
- Maximum Flow
- Algorithms
- Energy minimization with max flow/min cut

s-t Flow



Function flow: $A \rightarrow R$

Flow of arc \leq arc capacity

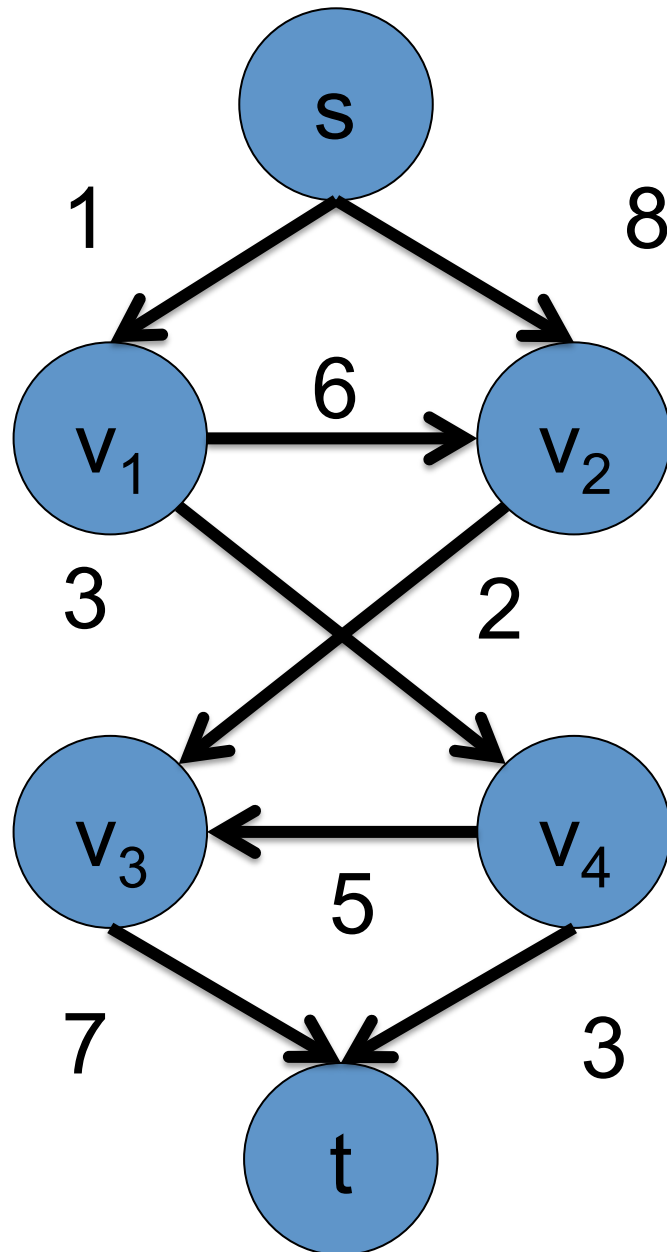
Flow is non-negative

For all vertex except s, t

Incoming flow

= Outgoing flow

s-t Flow



Function flow: $A \rightarrow R$

$\text{flow}(a) \leq c(a)$

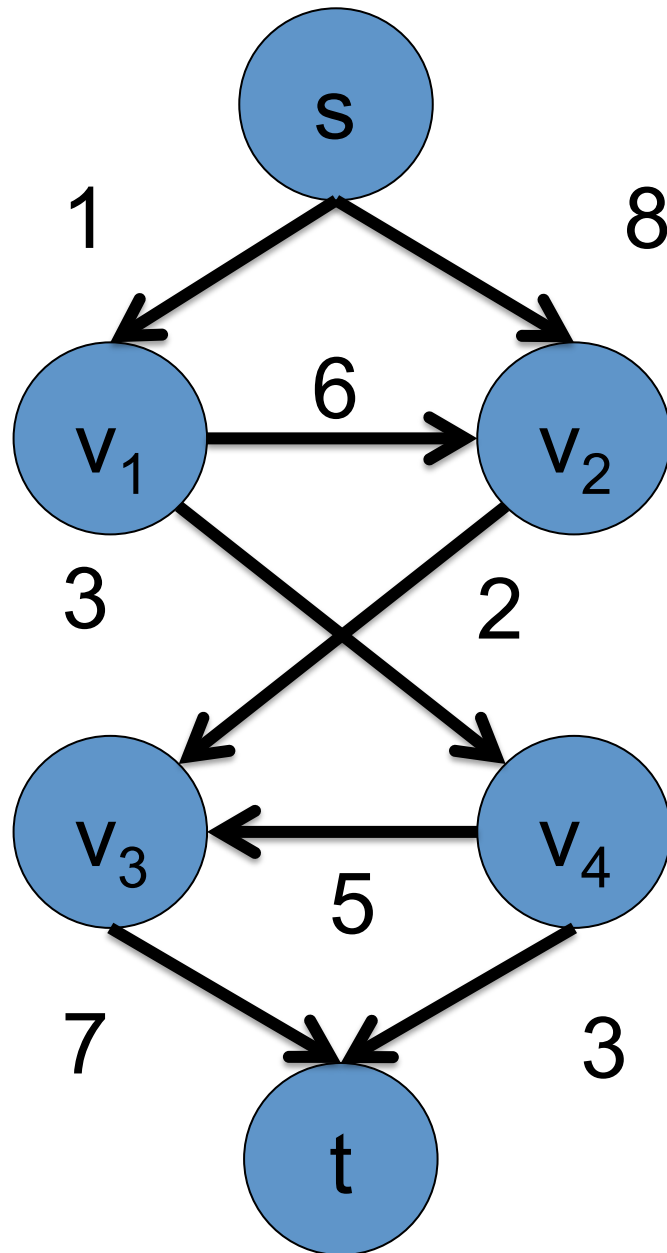
Flow is non-negative

For all vertex except s, t

Incoming flow

= Outgoing flow

s-t Flow



Function flow: $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

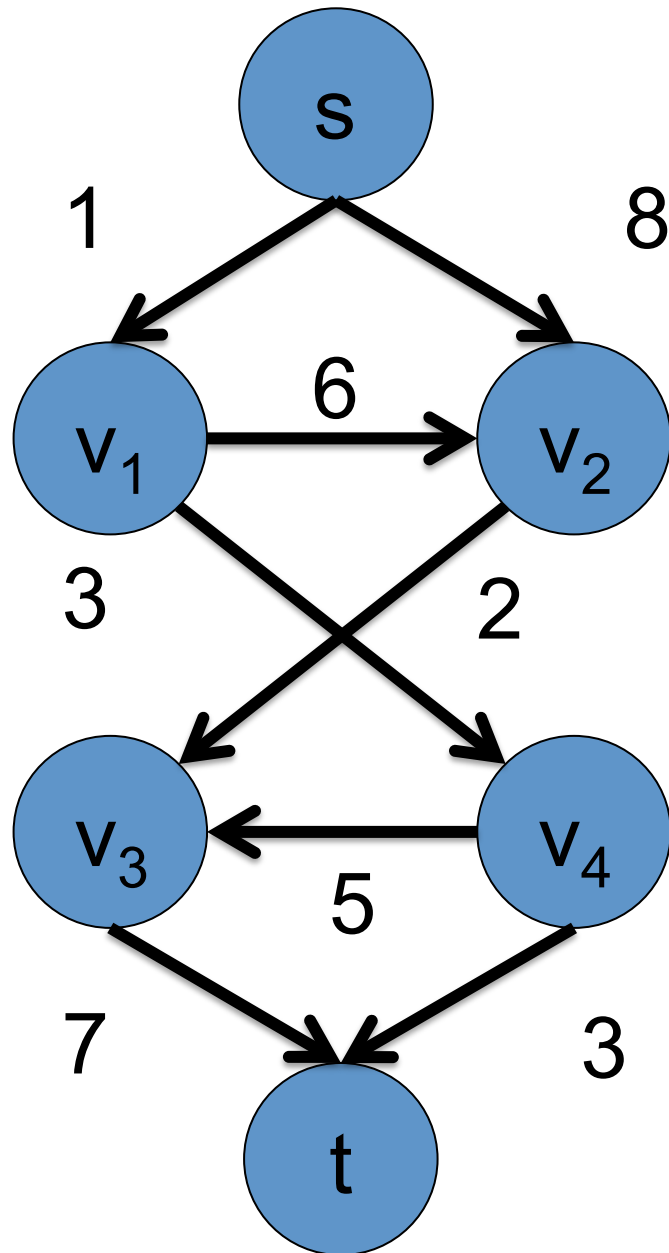
$$\text{flow}(a) \geq 0$$

For all vertex except s, t

Incoming flow

= Outgoing flow

s-t Flow



Function flow: $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

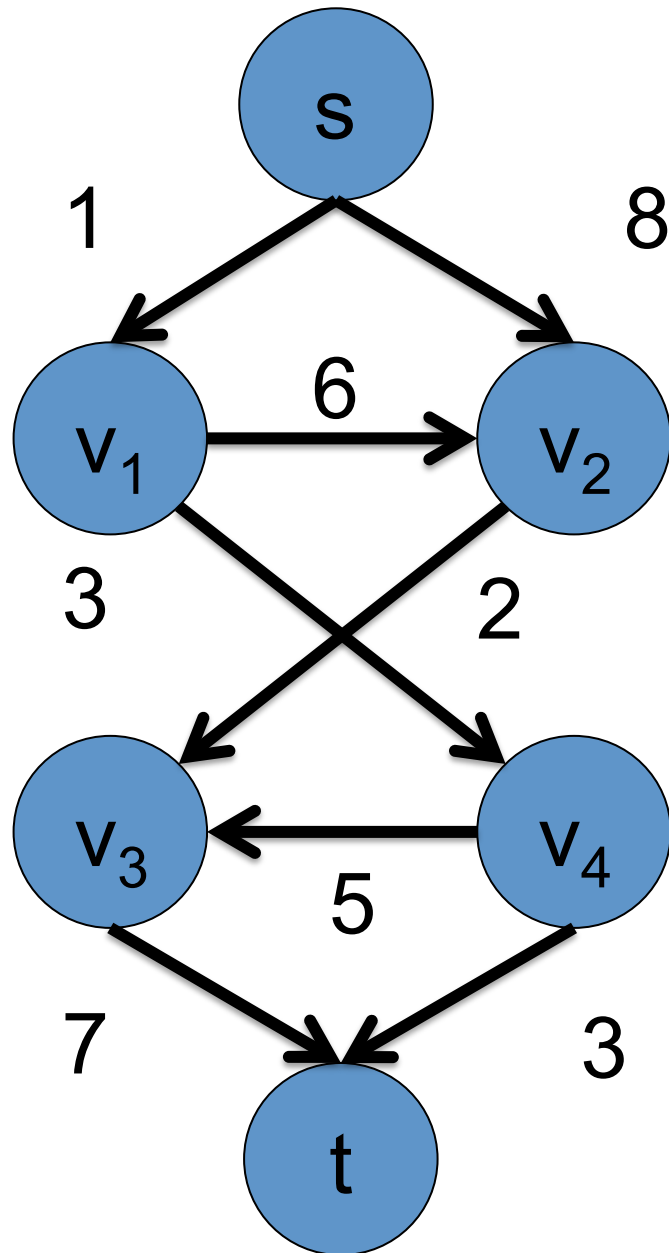
$$\text{flow}(a) \geq 0$$

For all $v \in V \setminus \{s, t\}$

Incoming flow

= Outgoing flow

s-t Flow



Function flow: $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

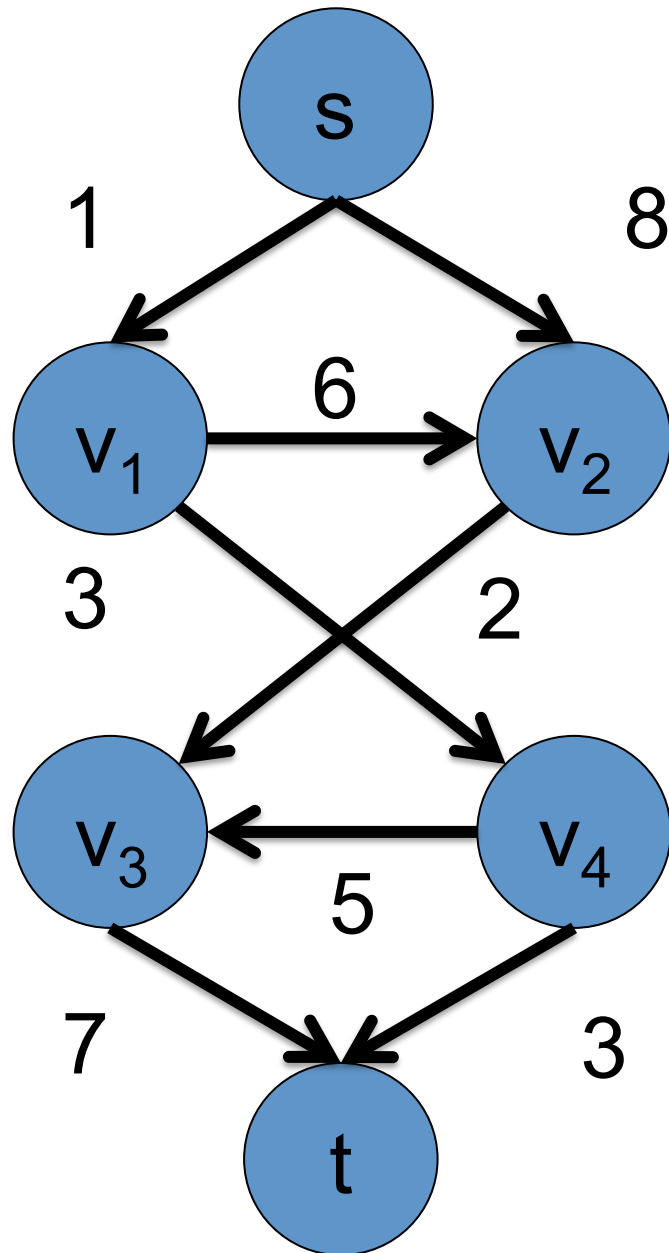
$$\text{flow}(a) \geq 0$$

For all $v \in V \setminus \{s, t\}$

$$\sum_{(u,v) \in A} \text{flow}((u,v))$$

= Outgoing flow

s-t Flow



Function flow: $A \rightarrow R$

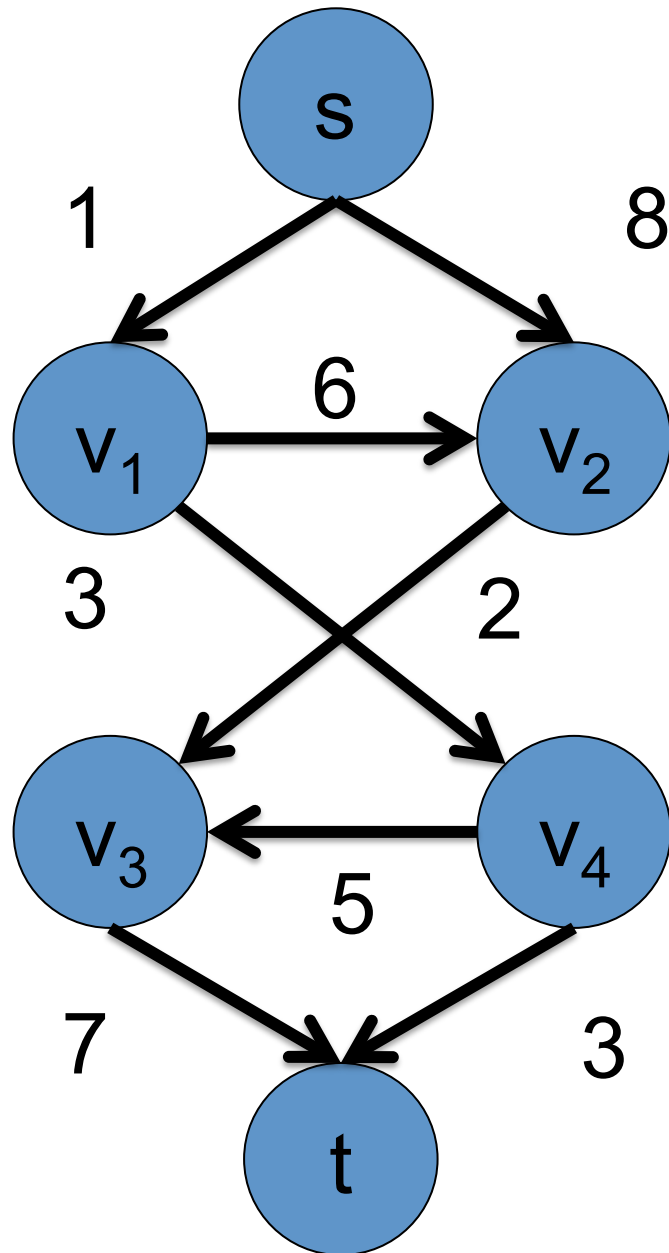
$$\text{flow}(a) \leq c(a)$$

$$\text{flow}(a) \geq 0$$

For all $v \in V \setminus \{s, t\}$

$$\begin{aligned} & \sum_{(u,v) \in A} \text{flow}((u,v)) \\ &= \sum_{(v,u) \in A} \text{flow}((v,u)) \end{aligned}$$

s-t Flow



Function flow: $A \rightarrow R$

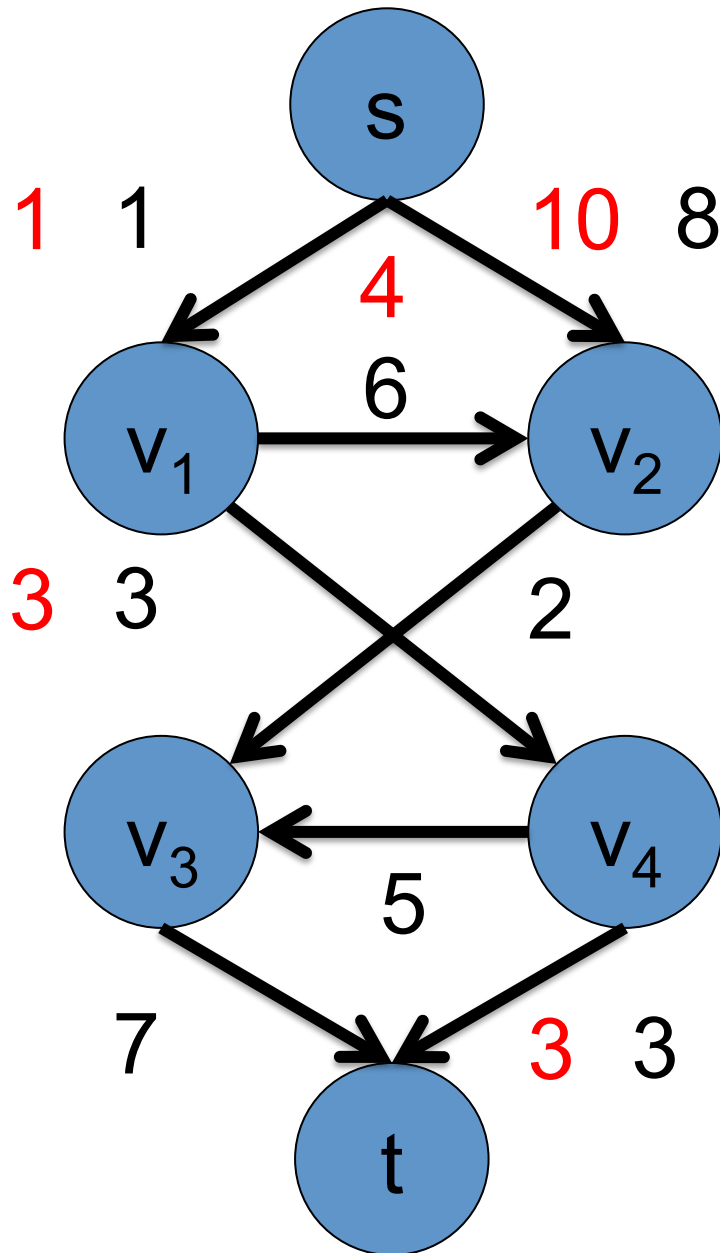
$$\text{flow}(a) \leq c(a)$$

$$\text{flow}(a) \geq 0$$

For all $v \in V \setminus \{s, t\}$

$$E_{\text{flow}}(v) = 0$$

s-t Flow



Function flow: $A \rightarrow R$

$\text{flow}(a) \leq c(a)$

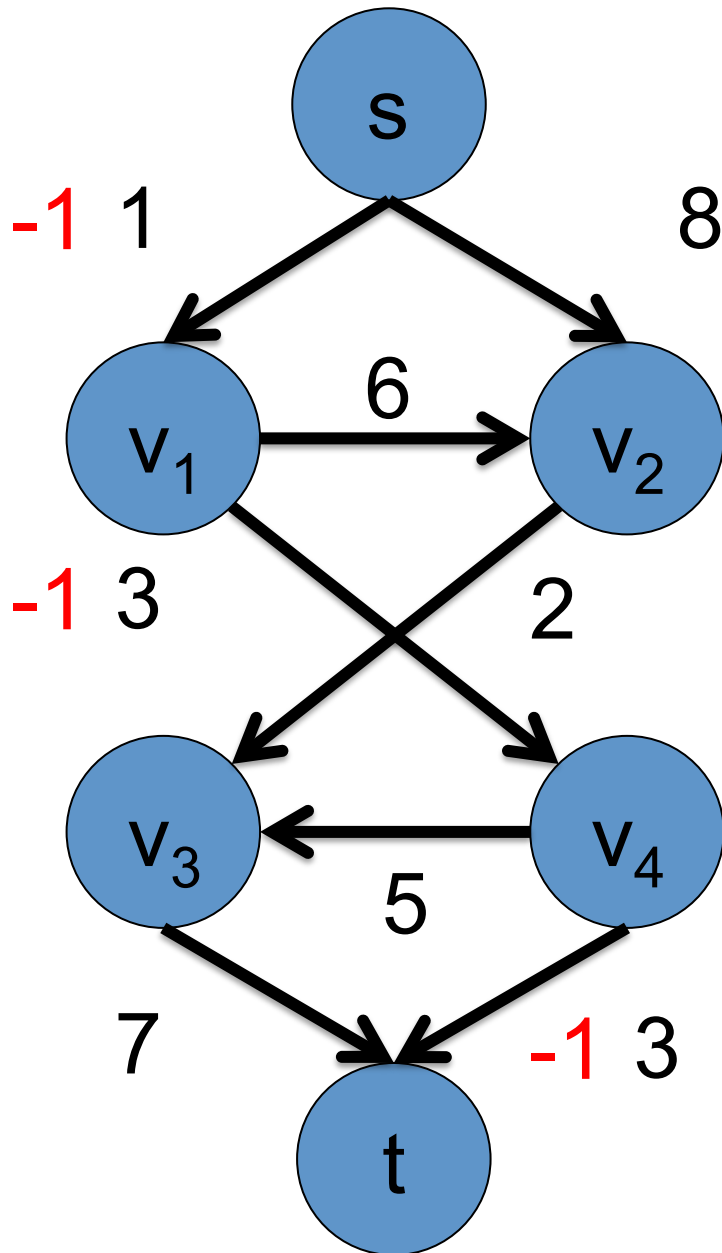
$\text{flow}(a) \geq 0$

For all $v \in V \setminus \{s, t\}$

$$E_{\text{flow}}(v) = 0$$

X

s-t Flow



Function flow: $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

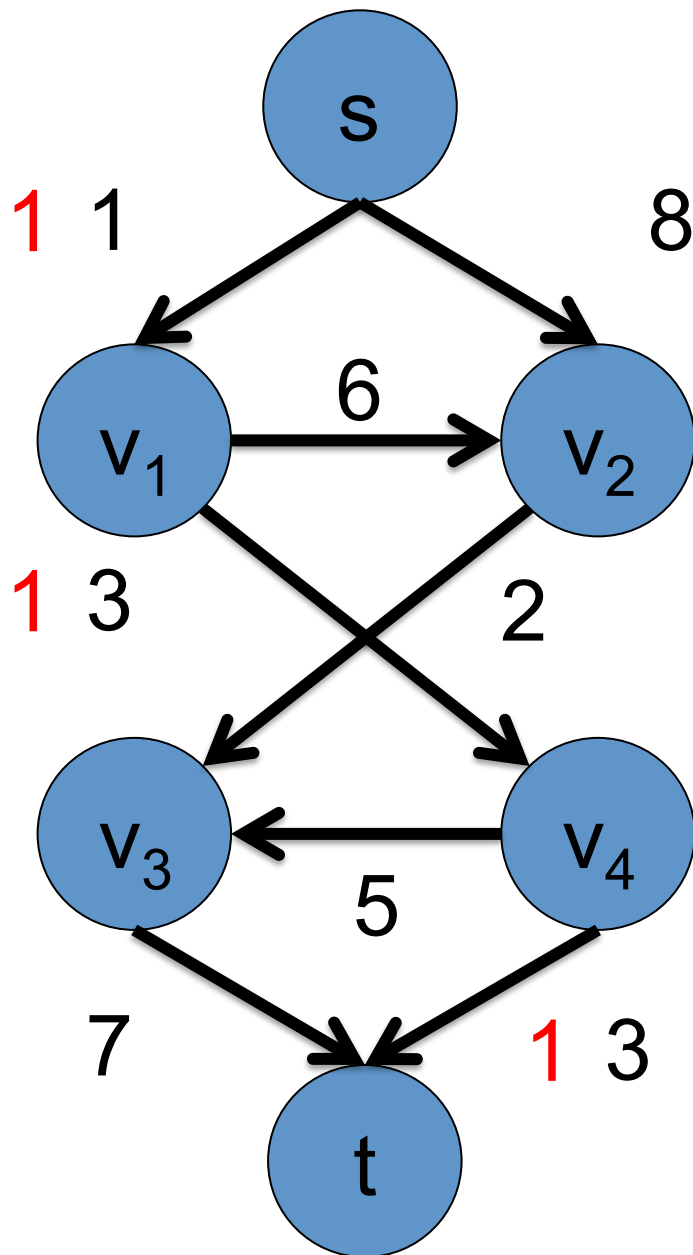
$$\text{flow}(a) \geq 0$$

For all $v \in V \setminus \{s, t\}$

$$E_{\text{flow}}(v) = 0$$

X

s-t Flow



Function flow: $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

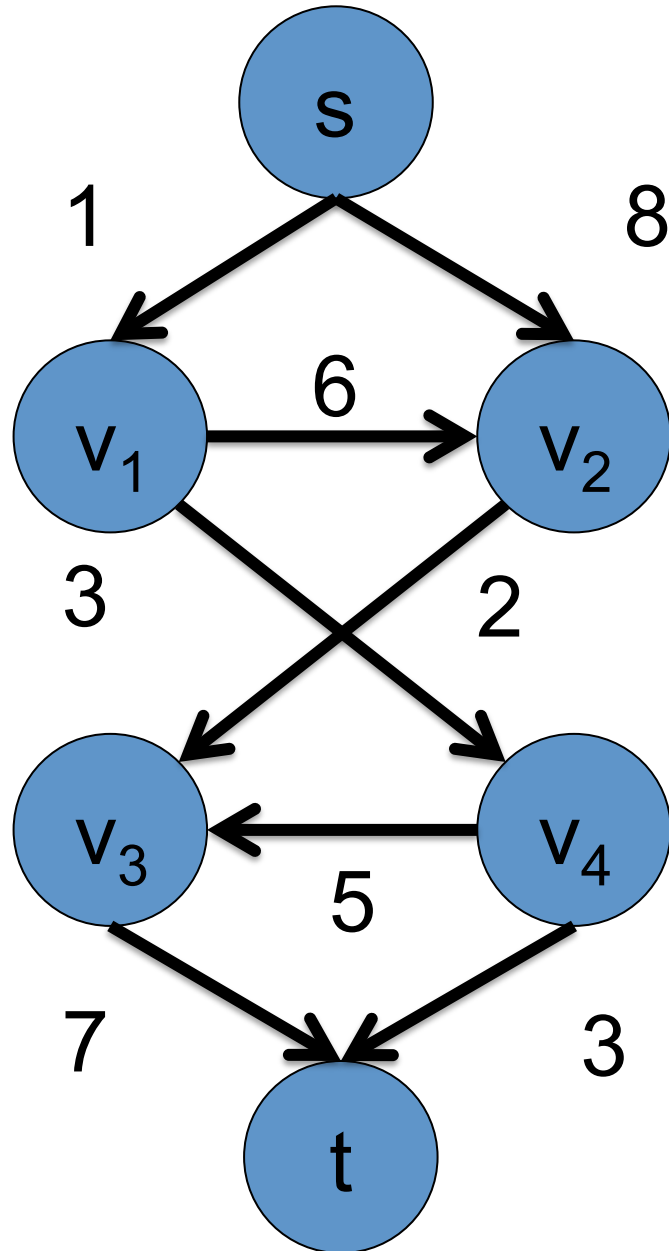
$$\text{flow}(a) \geq 0$$

For all $v \in V \setminus \{s, t\}$

$$E_{\text{flow}}(v) = 0$$

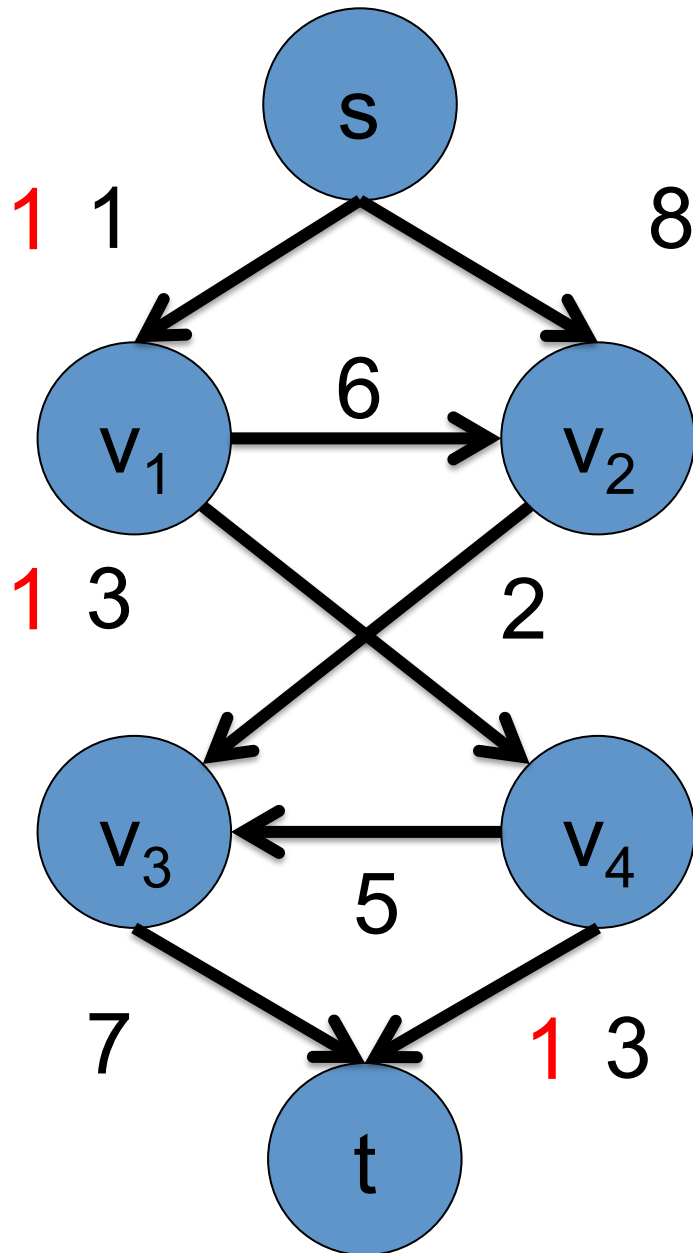


Value of s-t Flow



Outgoing flow of s
- Incoming flow of s

Value of s-t Flow



$$-E_{\text{flow}}(s) \quad E_{\text{flow}}(t)$$

$$\sum_{(s,v) \in A} \text{flow}((s,v))$$

$$- \sum_{(u,s) \in A} \text{flow}((u,s))$$

Value = 1

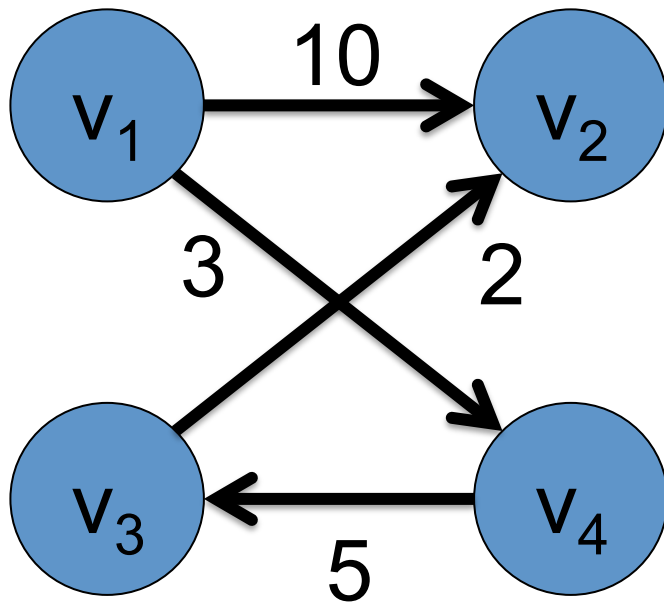
Outline

- Preliminaries
 - Functions and Excess Functions
 - s-t Flow
 - **s-t Cut**
 - Flows vs. Cuts
- Maximum Flow
- Algorithms
- Energy minimization with max flow/min cut

Cut

$$D = (V, A)$$

Let U be a subset of V



C is a set of arcs such that

- $(u, v) \in A$
- $u \in U$
- $v \in V \setminus U$

C is a cut in the digraph D

Cut

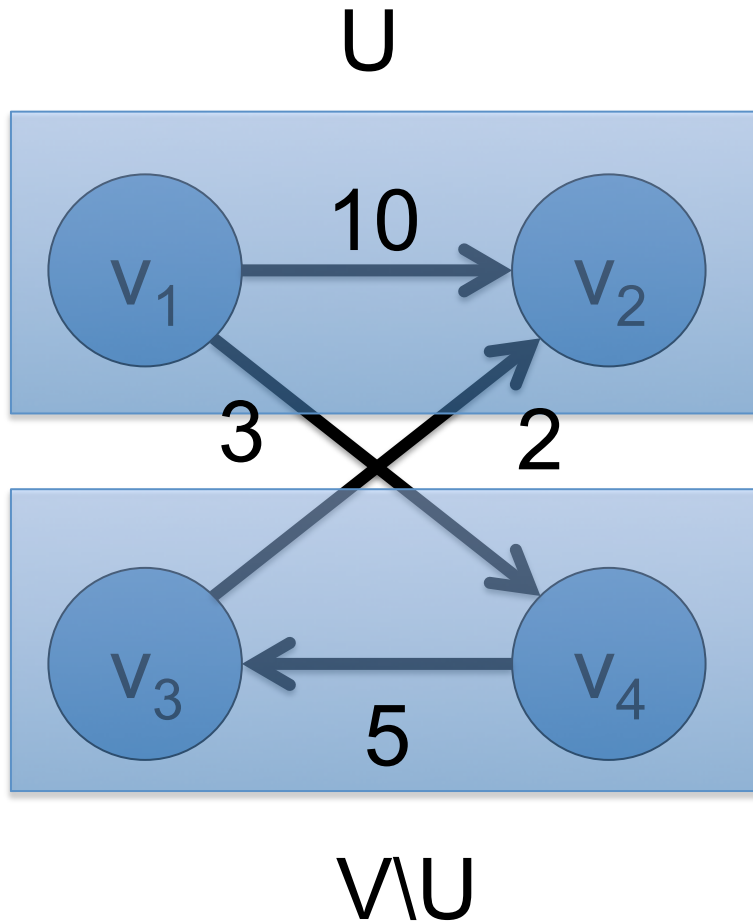
$$D = (V, A)$$

What is C?

$$\{(v_1, v_2), (v_1, v_4)\} ?$$

$$\{(v_1, v_4), (v_3, v_2)\} ?$$

$$\{(v_1, v_4)\} ?$$



Cut

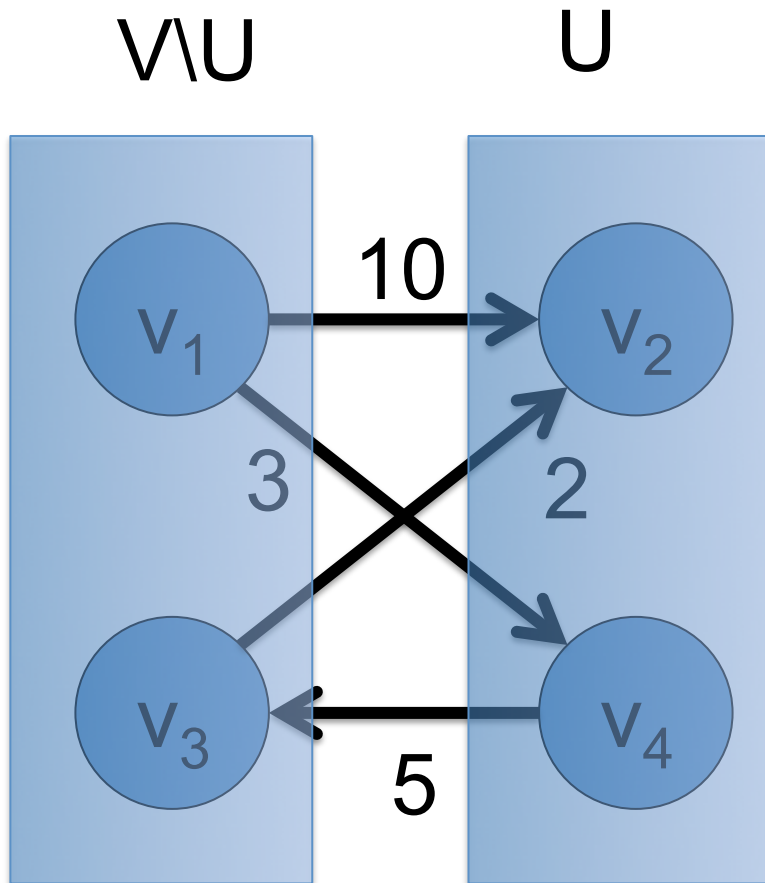
$$D = (V, A)$$

What is C?

$$\{(v_1, v_2), (v_1, v_4), (v_3, v_2)\} ?$$

✓ $\{(v_4, v_3)\} ?$

$$\{(v_1, v_4), (v_3, v_2)\} ?$$



Cut

$$D = (V, A)$$

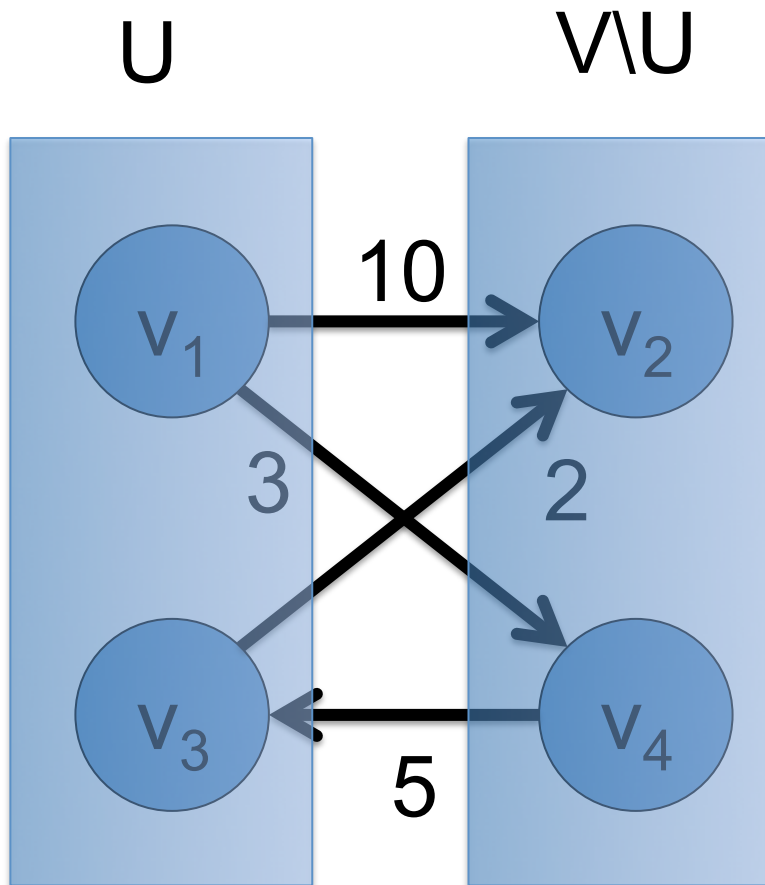
What is C?



$\{(v_1, v_2), (v_1, v_4), (v_3, v_2)\}$?

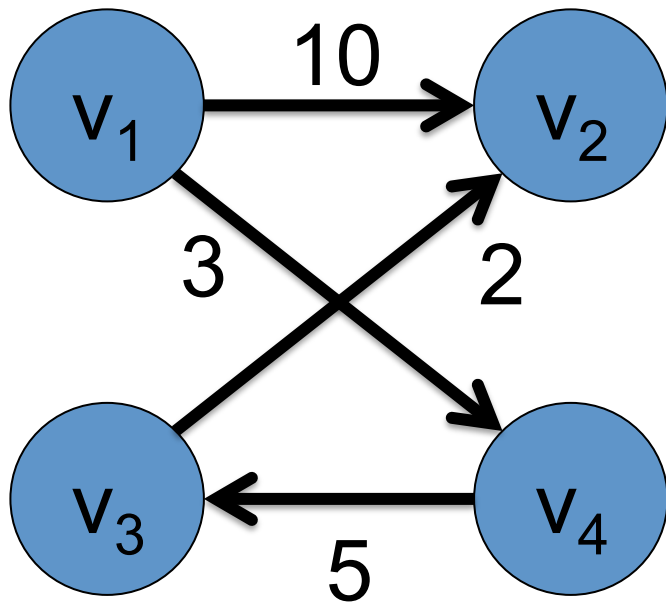
$\{(v_3, v_2)\}$?

$\{(v_1, v_4), (v_3, v_2)\}$?



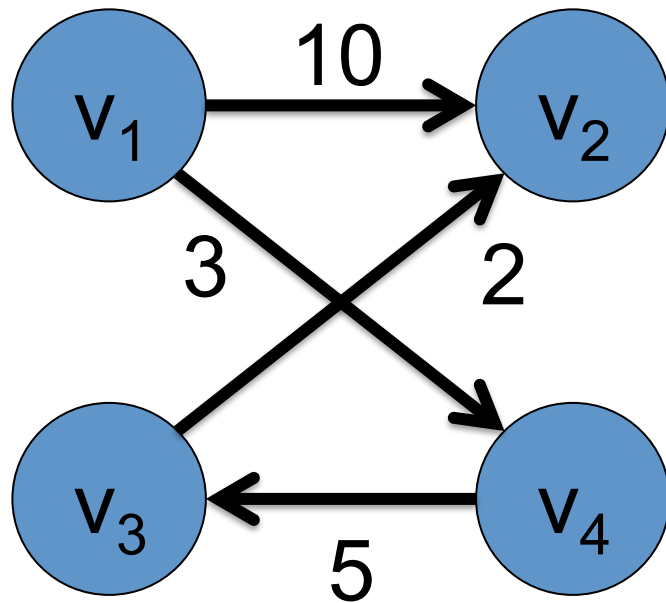
Cut

$$D = (V, A)$$



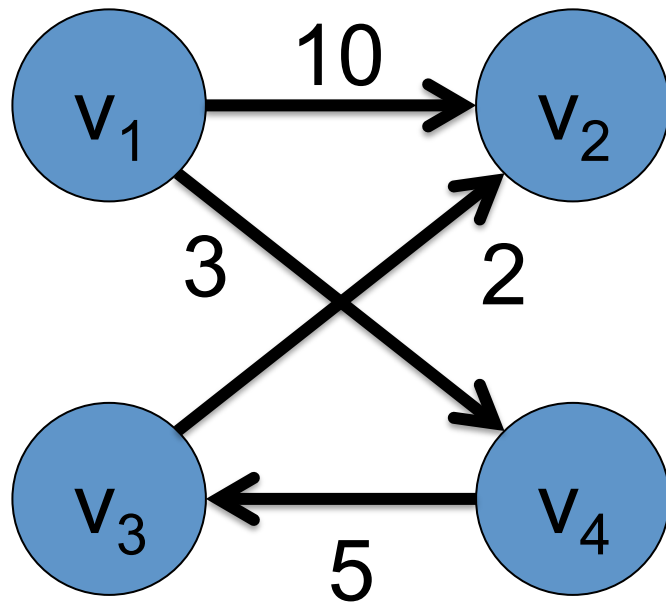
$$C = \text{out-arcs}(U)$$

Capacity of Cut



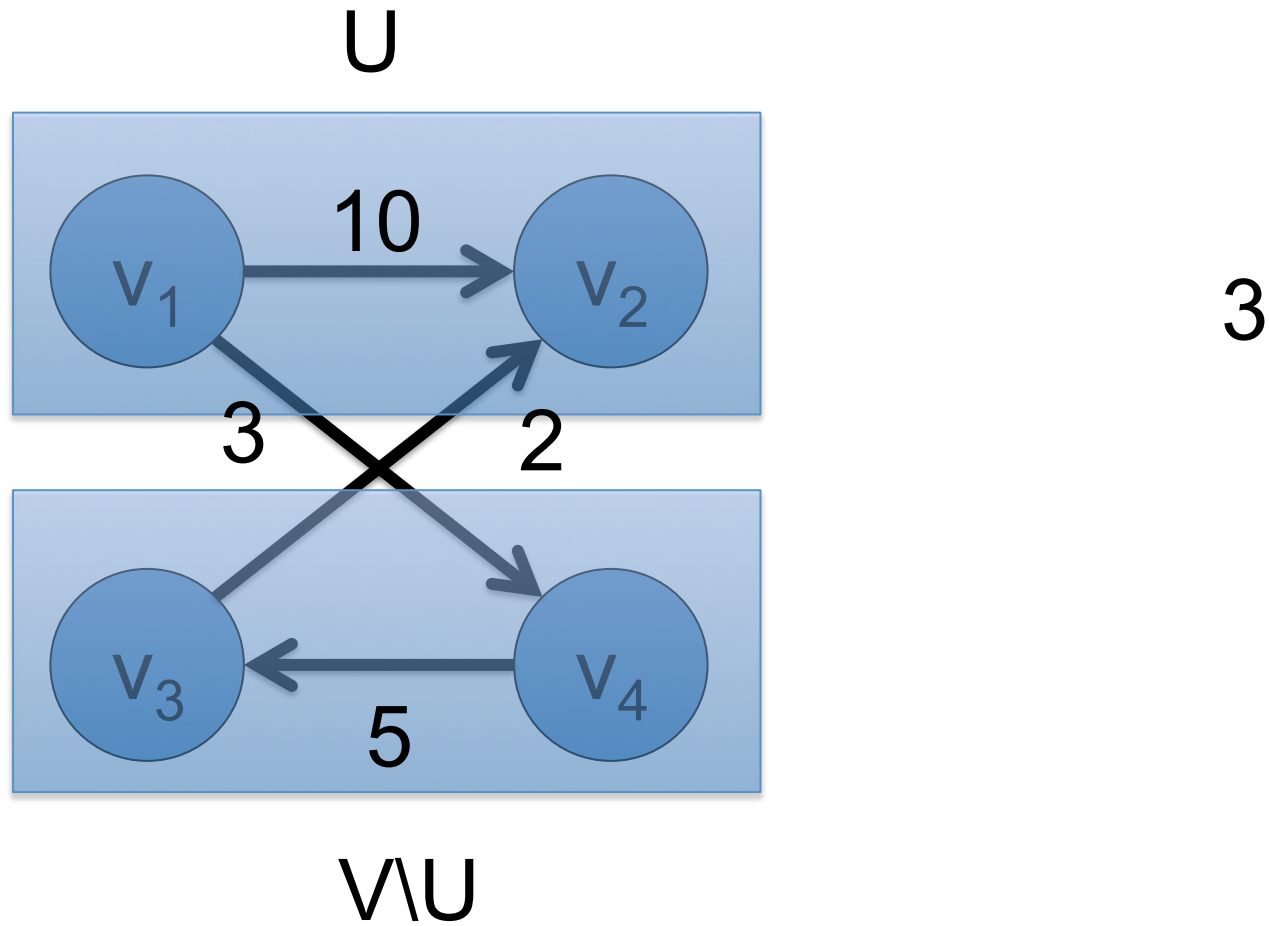
Sum of capacity of all arcs in C

Capacity of Cut

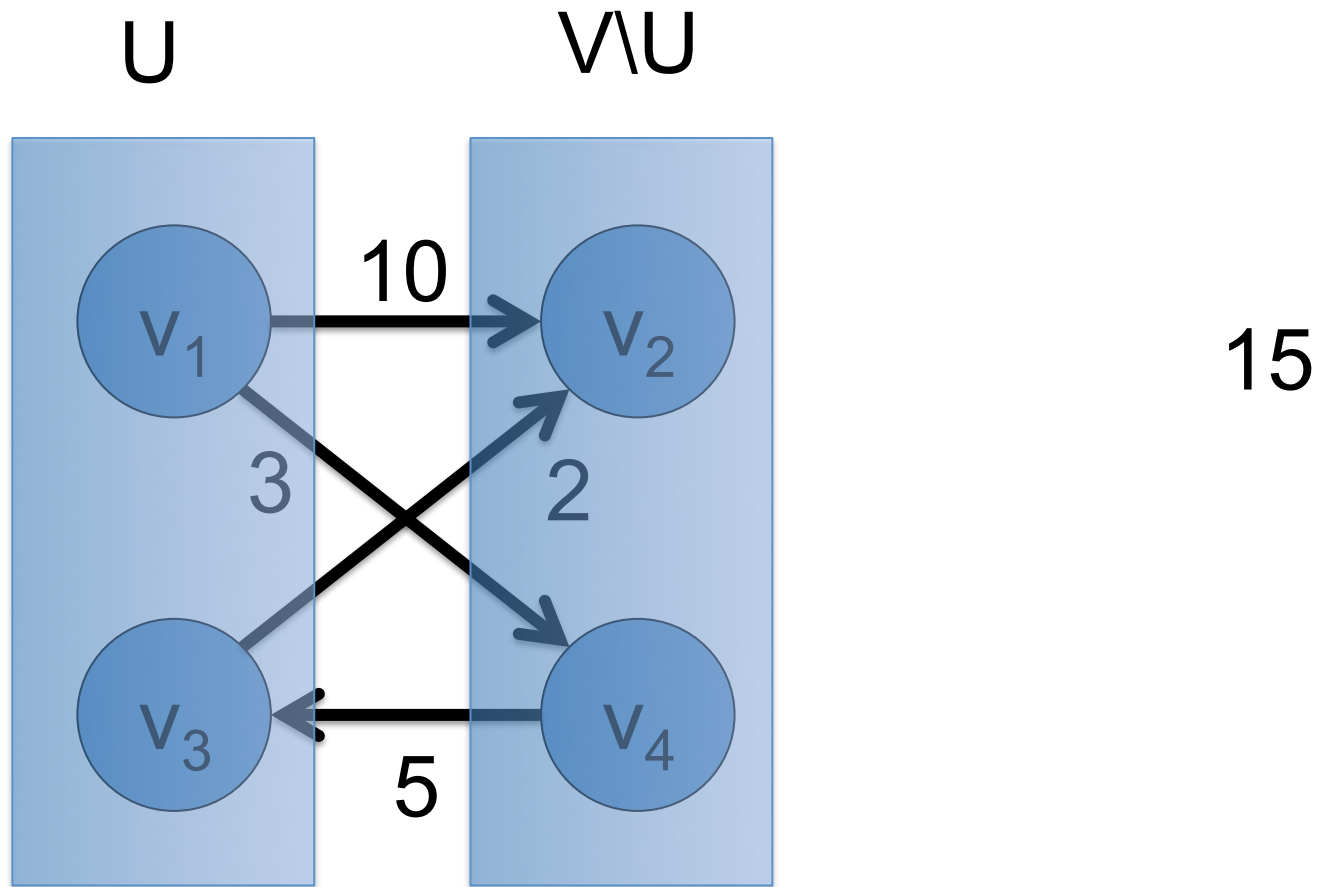


$$\sum_{a \in C} c(a)$$

Capacity of Cut

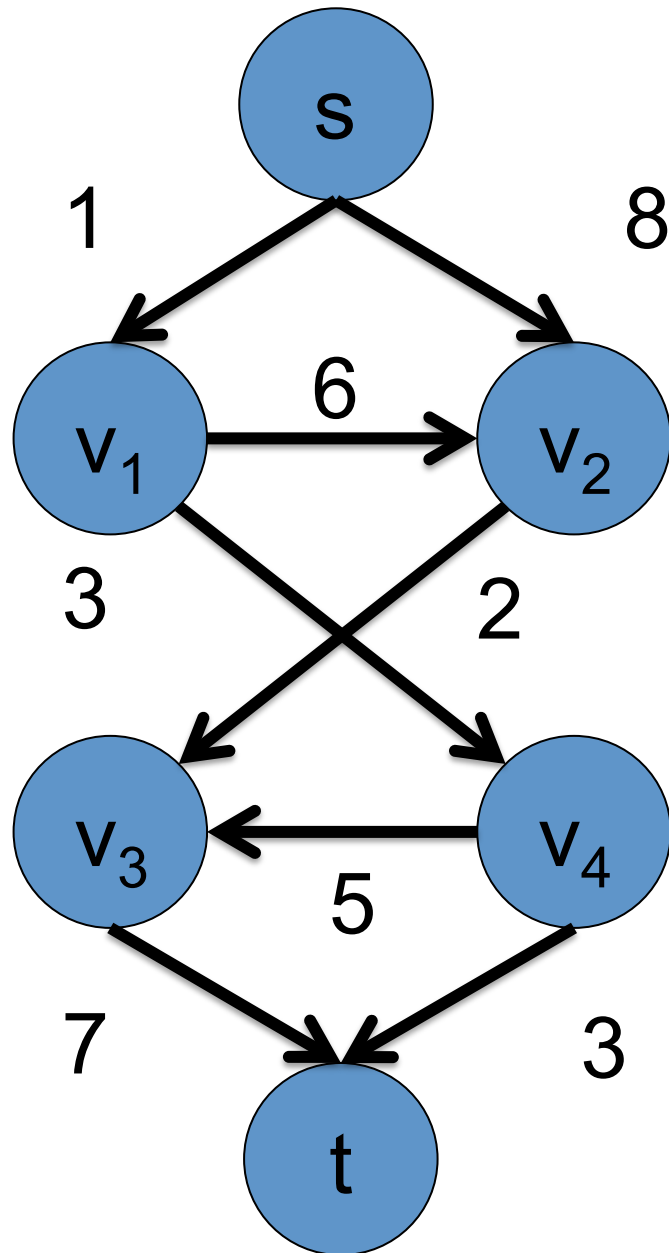


Capacity of Cut



s-t Cut

$$D = (V, A)$$



A source vertex “s”

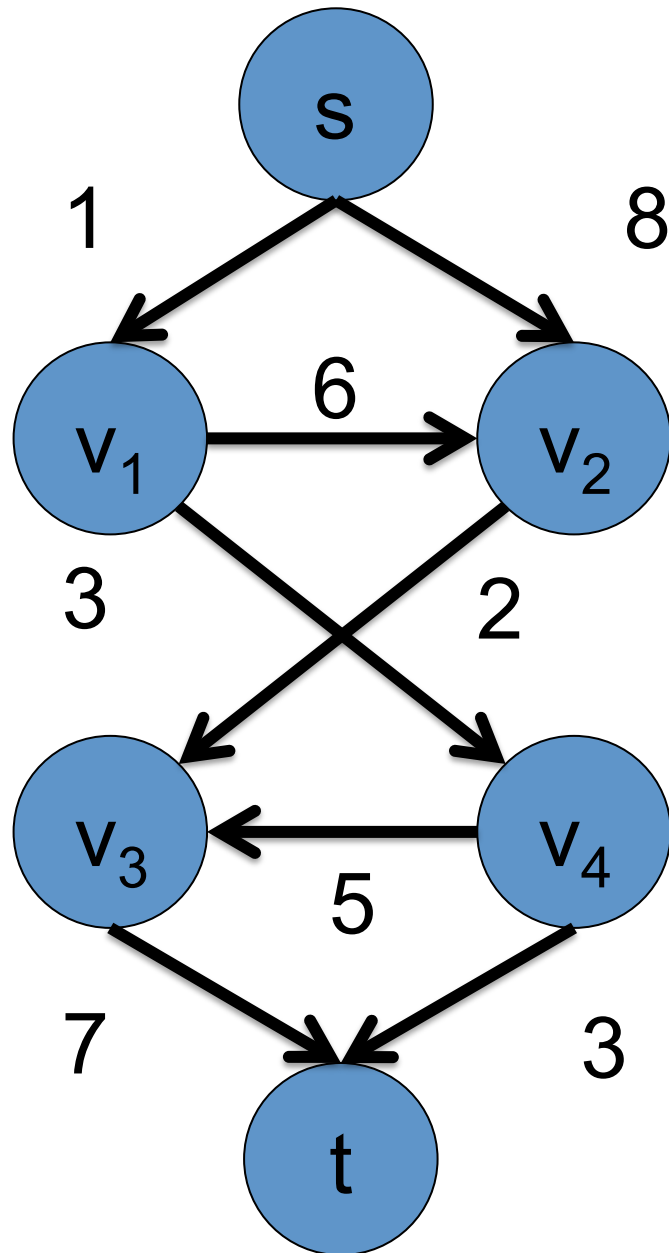
A sink vertex “t”

C is a cut such that

- $s \in U$
- $t \in V \setminus U$

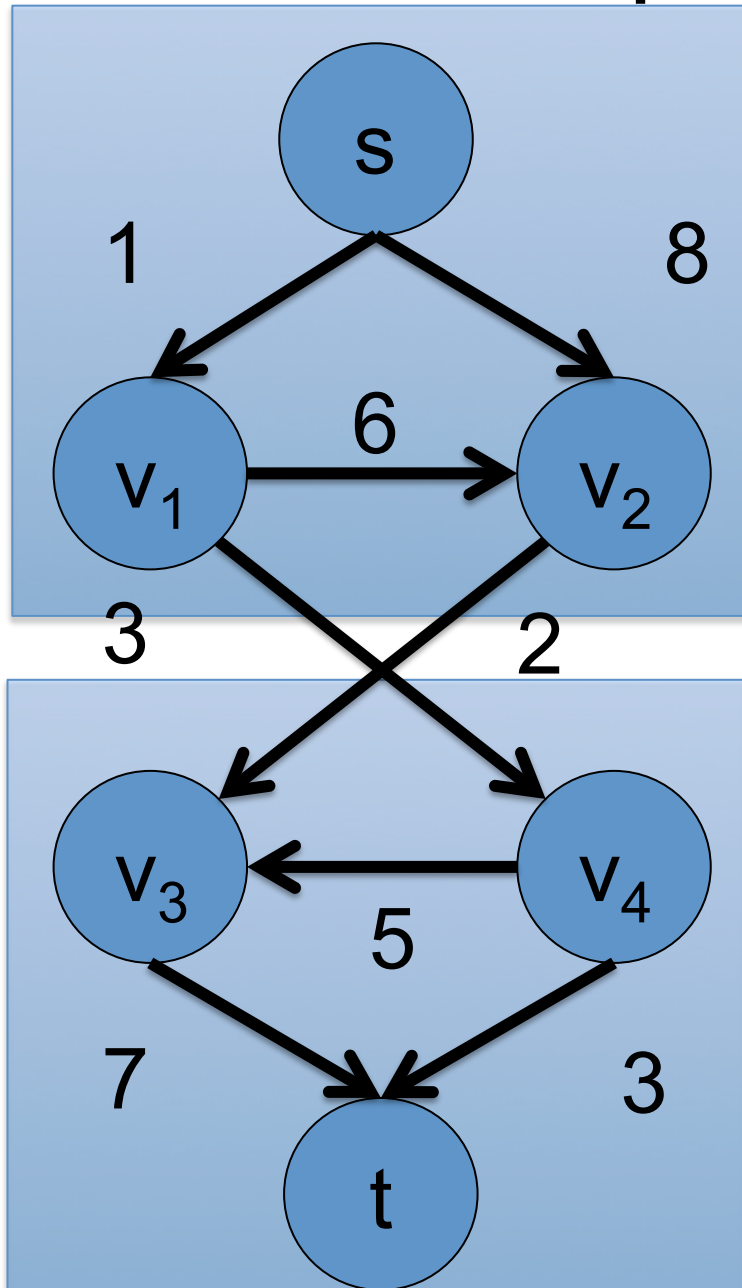
C is an s-t cut

Capacity of s-t Cut



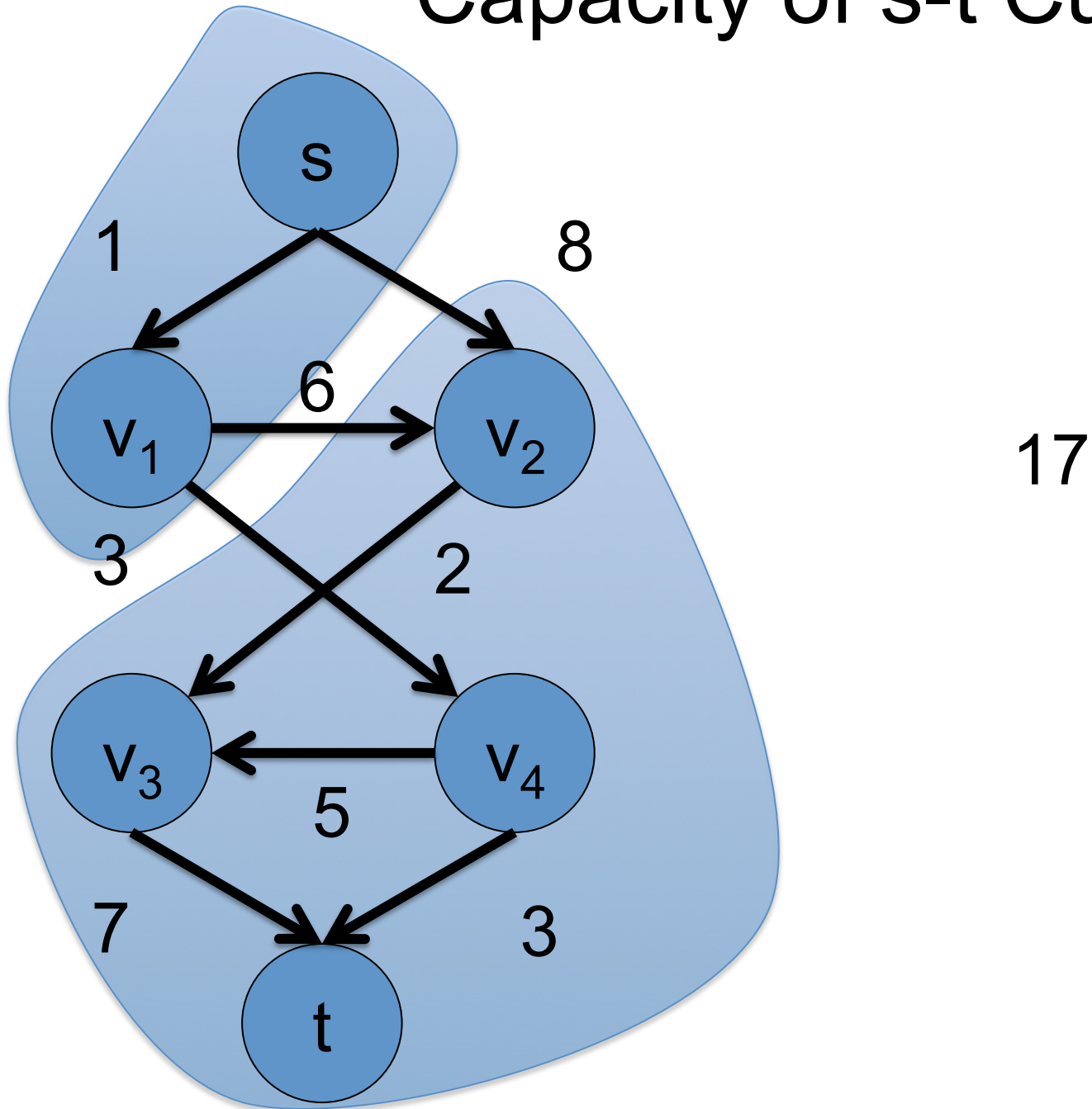
$$\sum_{a \in C} c(a)$$

Capacity of s-t Cut



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Capacity of s-t Cut



Outline

- Preliminaries
 - s-t Flow
 - s-t Cut
 - **Flows vs. Cuts**
- Maximum Flow
- Algorithms
- Energy minimization with max flow/min cut

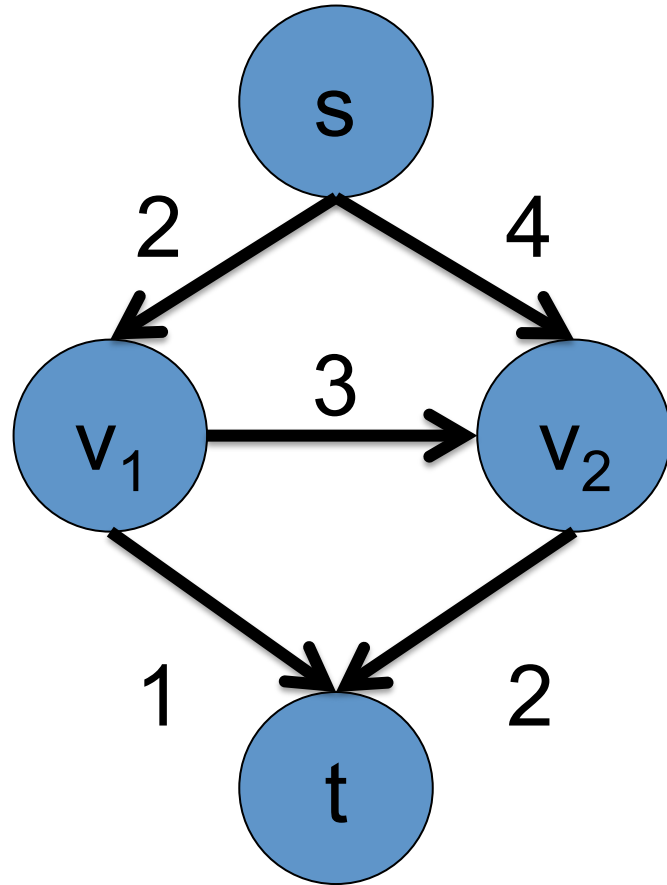
Flows vs. Cuts

(See lecture 5)

Outline

- Preliminaries
- **Maximum Flow**
 - Residual Graph
 - Max-Flow Min-Cut Theorem
- Algorithms
- Energy minimization with max flow/min cut

Maximum Flow Problem



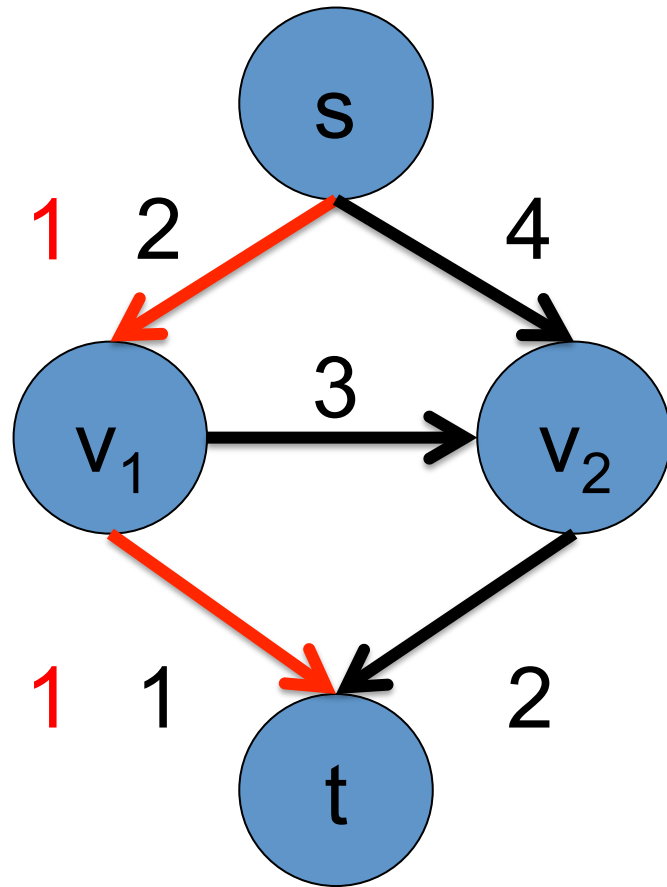
Find the flow with the maximum value !!

$$\sum_{(s,v) \in A} \text{flow}((s,v))$$

$$- \sum_{(u,s) \in A} \text{flow}((u,s))$$

First suggestion to solve this problem !!

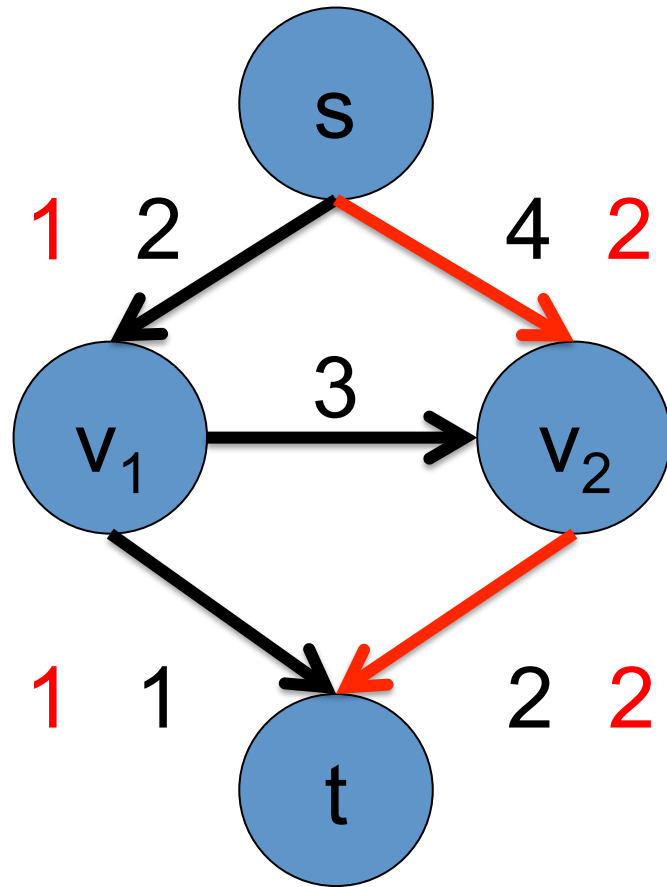
Passing Flow through s-t Paths



Find an s-t path where
 $\text{flow}(a) < c(a)$ for all arcs

Pass maximum allowable
flow through the arcs

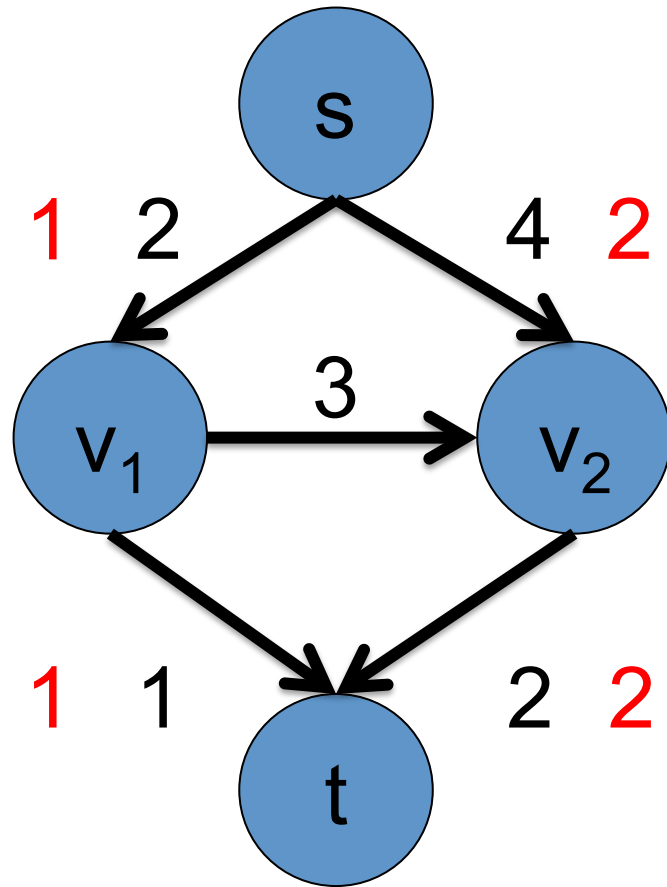
Passing Flow through s-t Paths



Find an s-t path where
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Pass maximum allowable
flow through the arcs

Passing Flow through s-t Paths



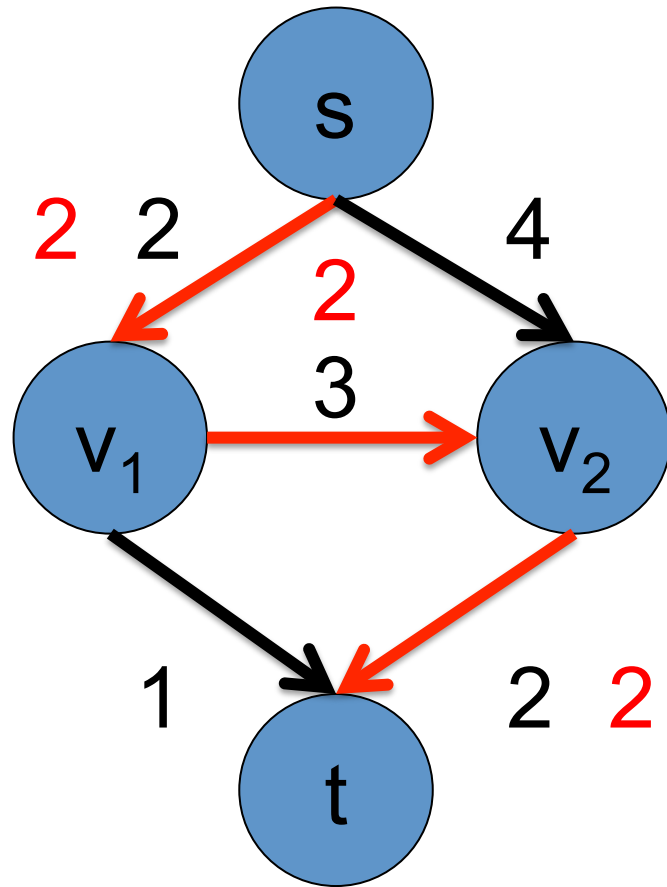
Find an s-t path where
 $\text{flow}(a) < c(a)$ for all arcs

No more paths. Stop.

Will this give us maximum flow?

NO !!!

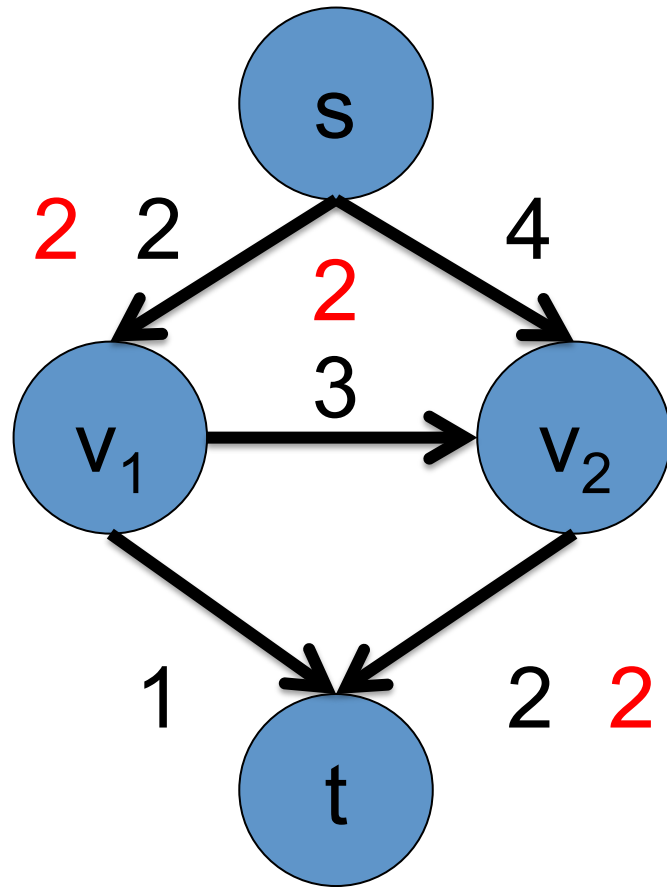
Passing Flow through s-t Paths



Find an s-t path where
 $\text{flow}(a) < c(a)$ for all arcs

Pass maximum allowable
flow through the arcs

Passing Flow through s-t Paths



Find an s-t path where
 $\text{flow}(a) < c(a)$ for all arcs

No more paths. Stop.

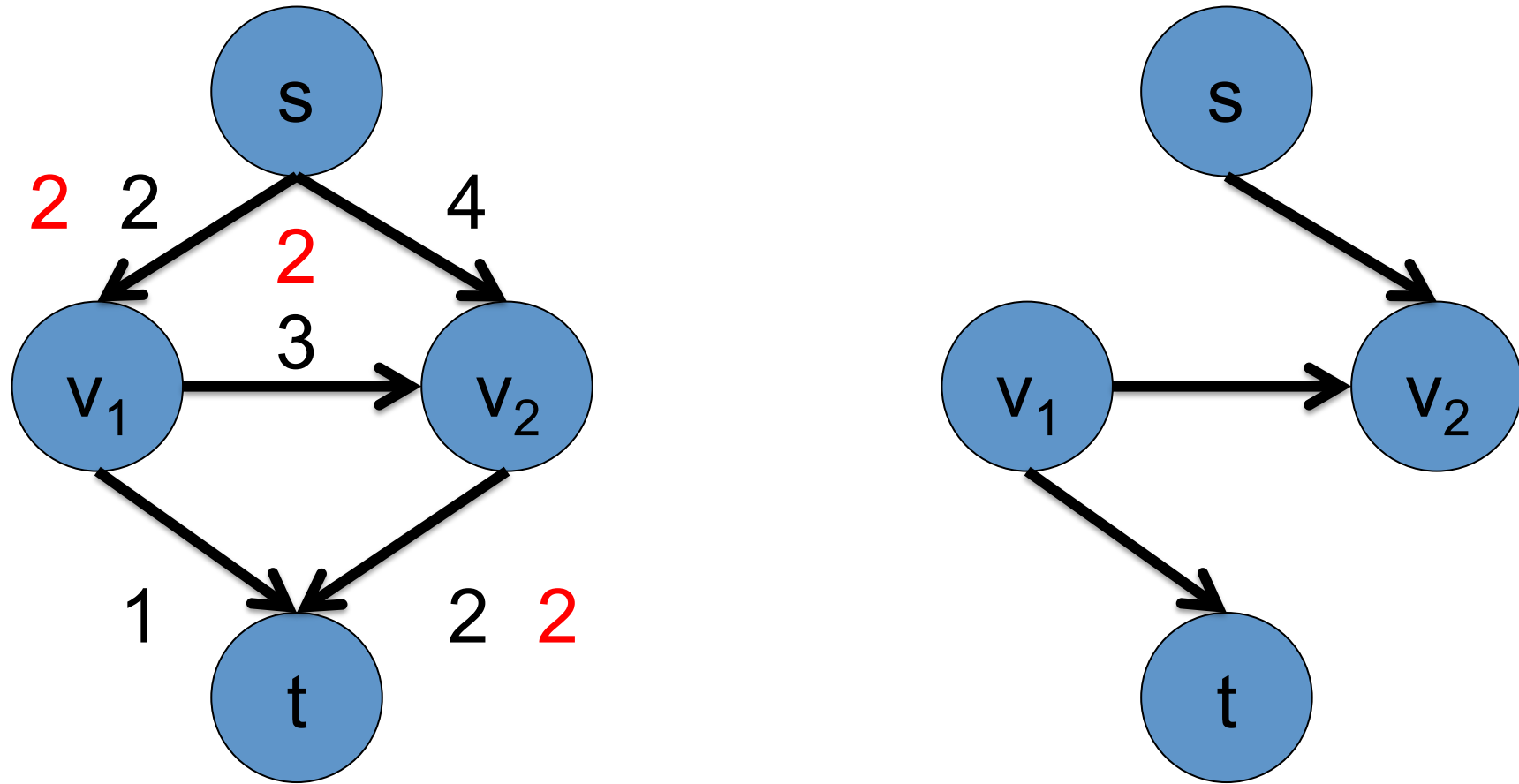
Another method?

Incorrect Answer !!

Outline

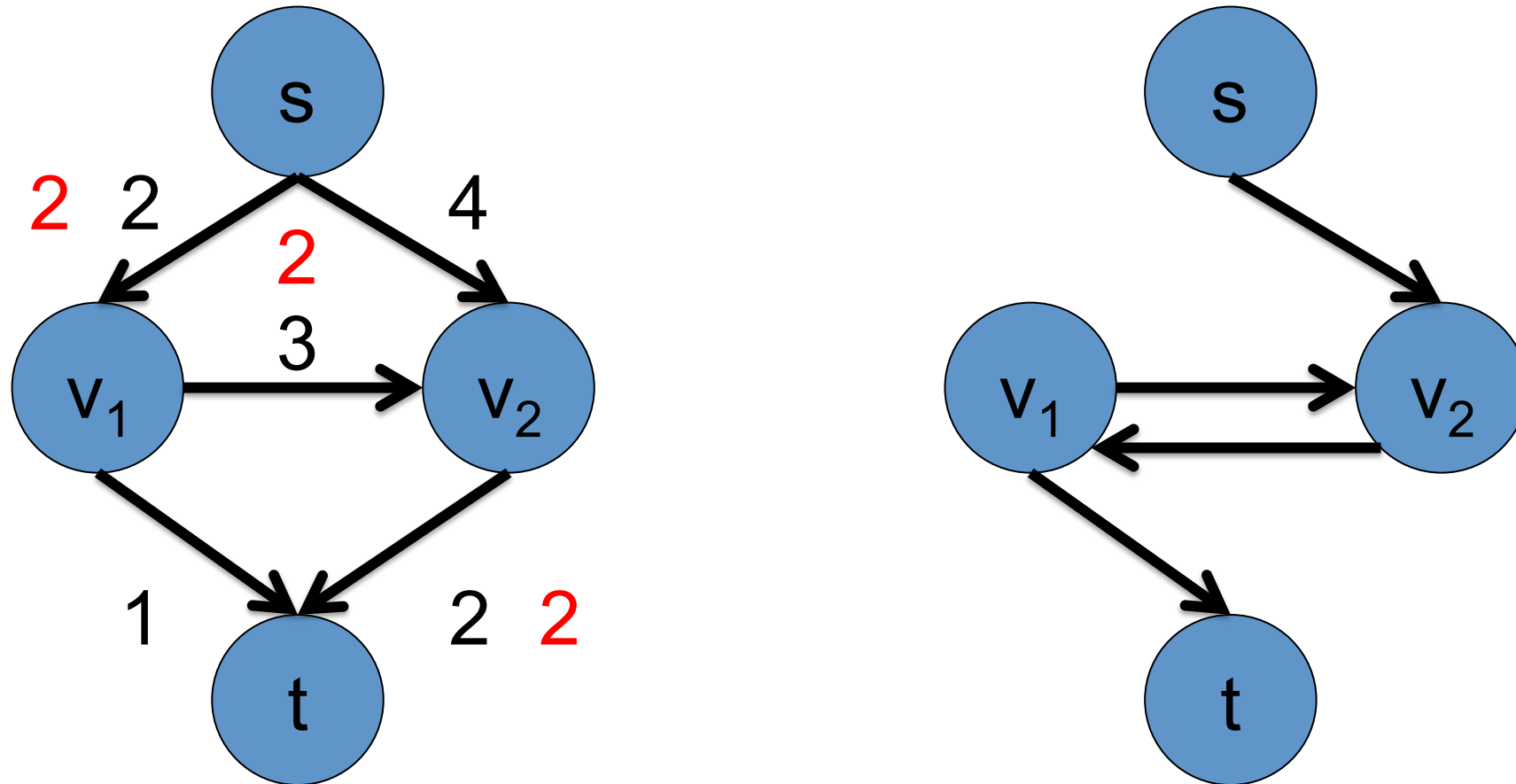
- Preliminaries
- Maximum Flow
 - **Residual Graph**
 - Max-Flow Min-Cut Theorem
- Algorithms
- Energy minimization with max flow/min cut

Residual Graph



Arcs where $\text{flow}(a) < c(a)$

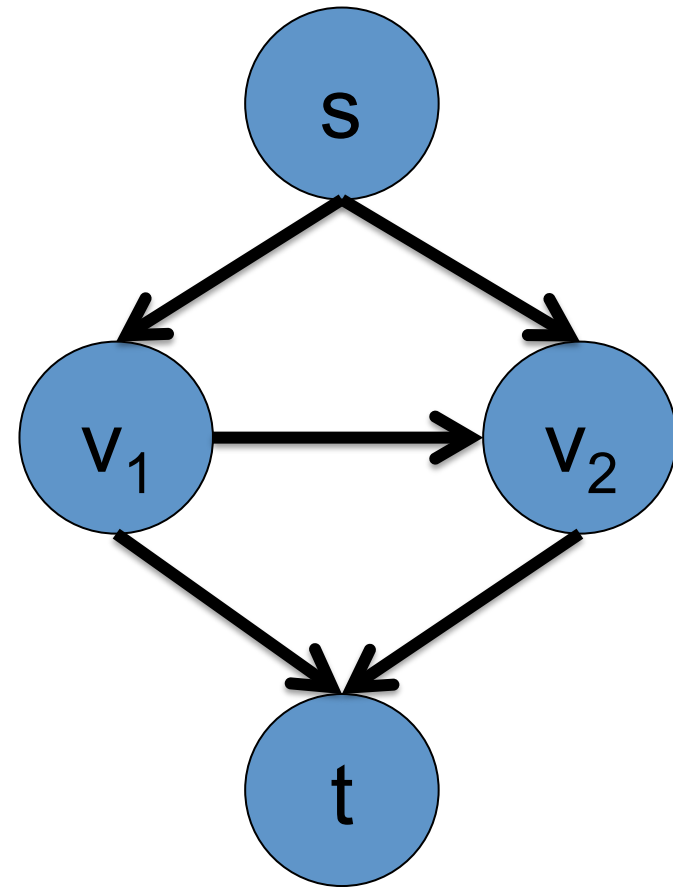
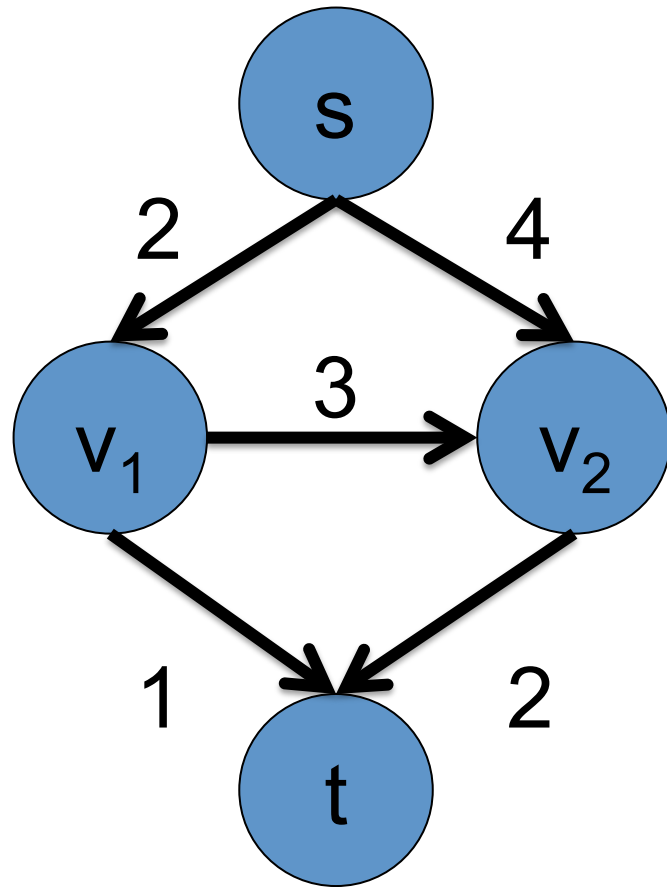
Residual Graph



Including arcs to s and from t is not necessary

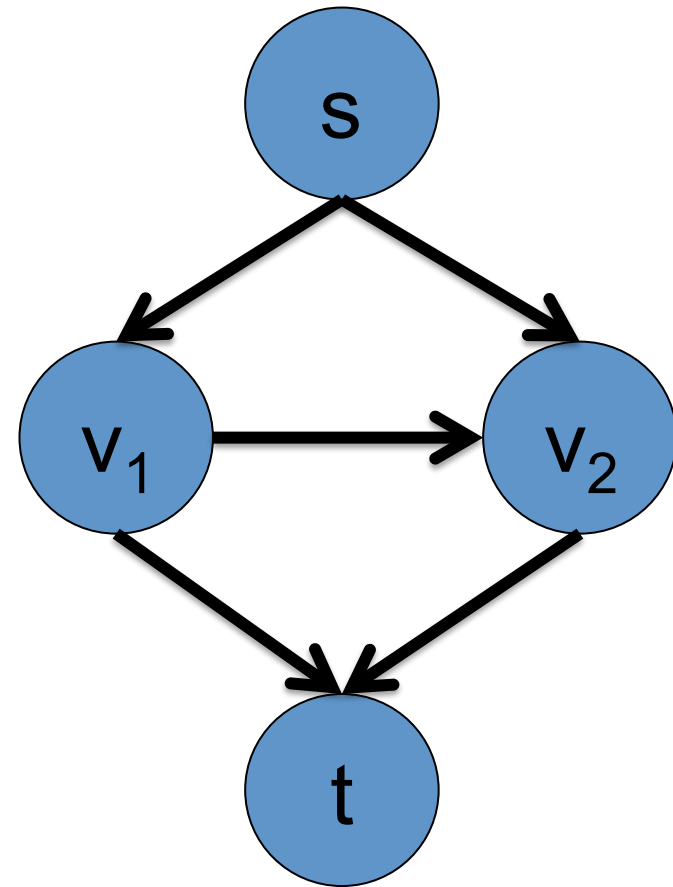
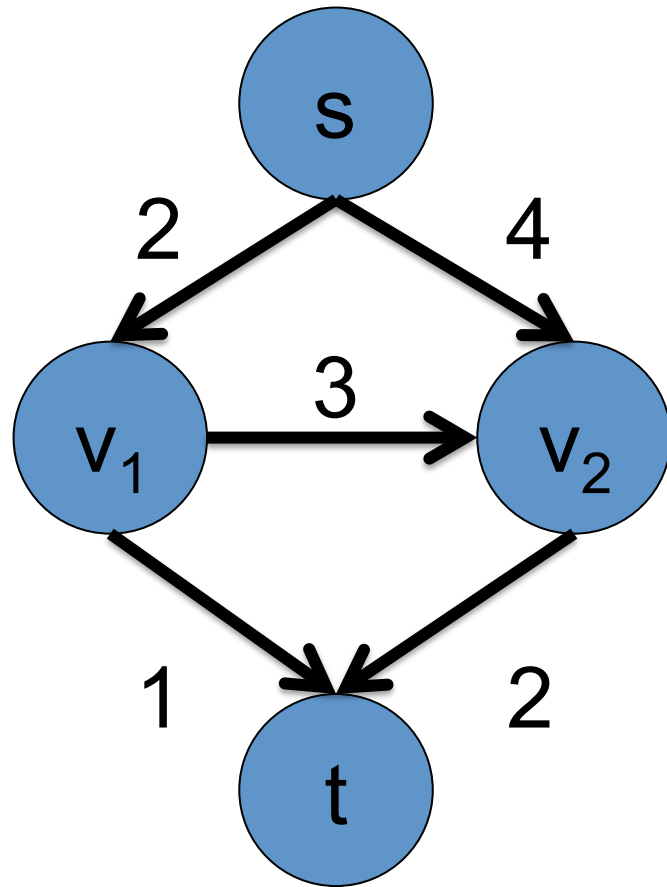
Inverse of arcs where $\text{flow}(a) > 0$

Maximum Flow using Residual Graphs



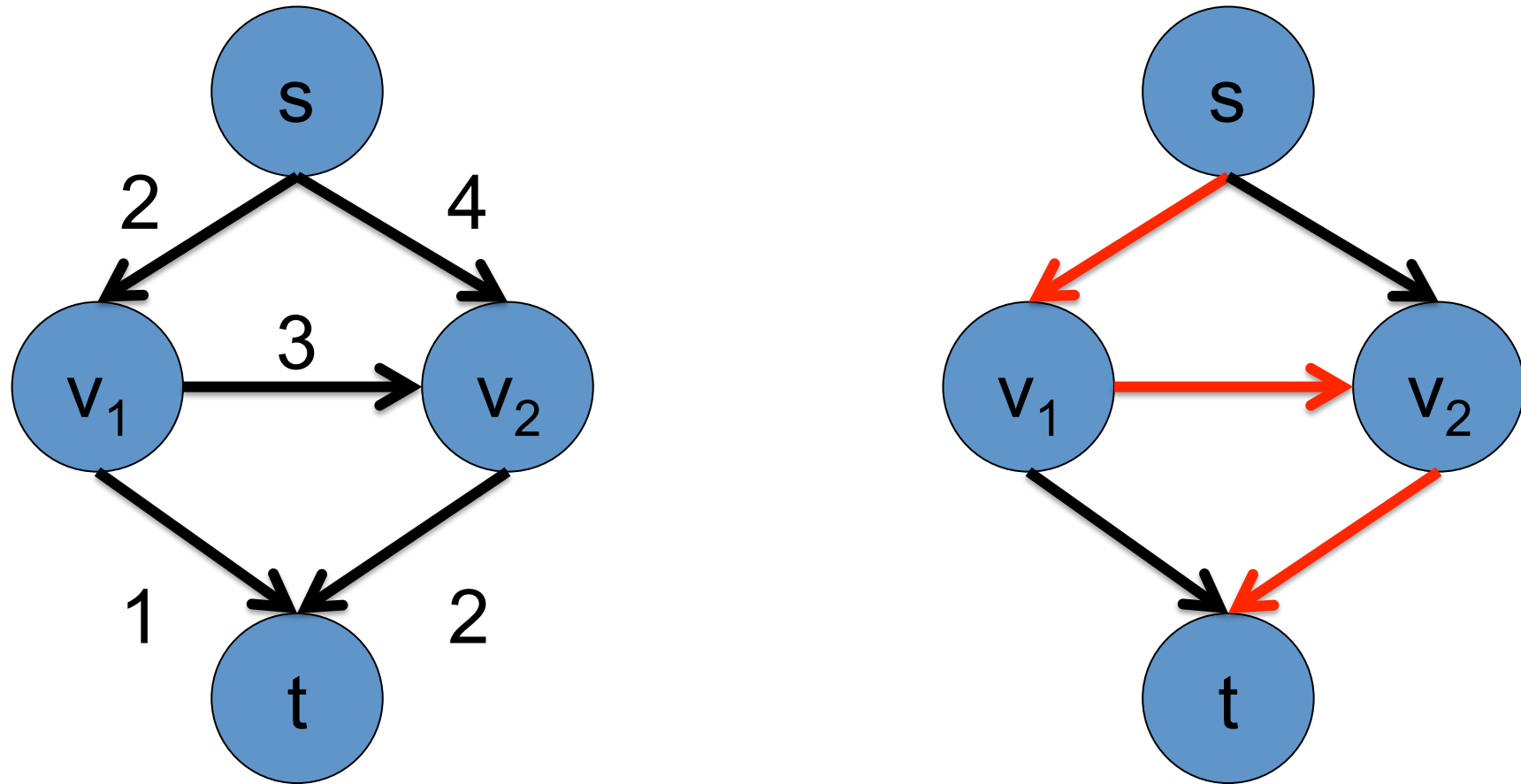
Start with zero flow.

Maximum Flow using Residual Graphs



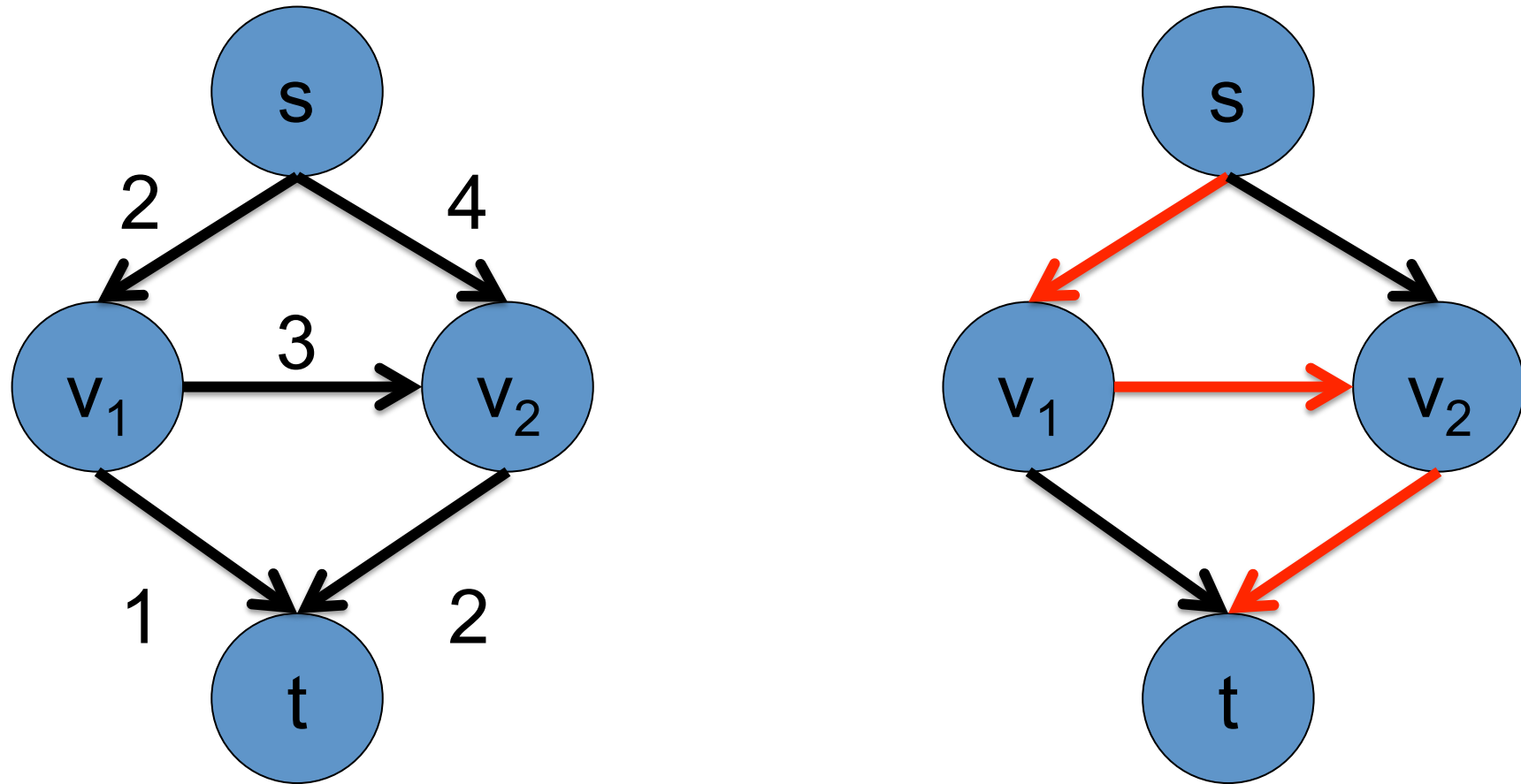
Find an s - t path in the residual graph.

Maximum Flow using Residual Graphs



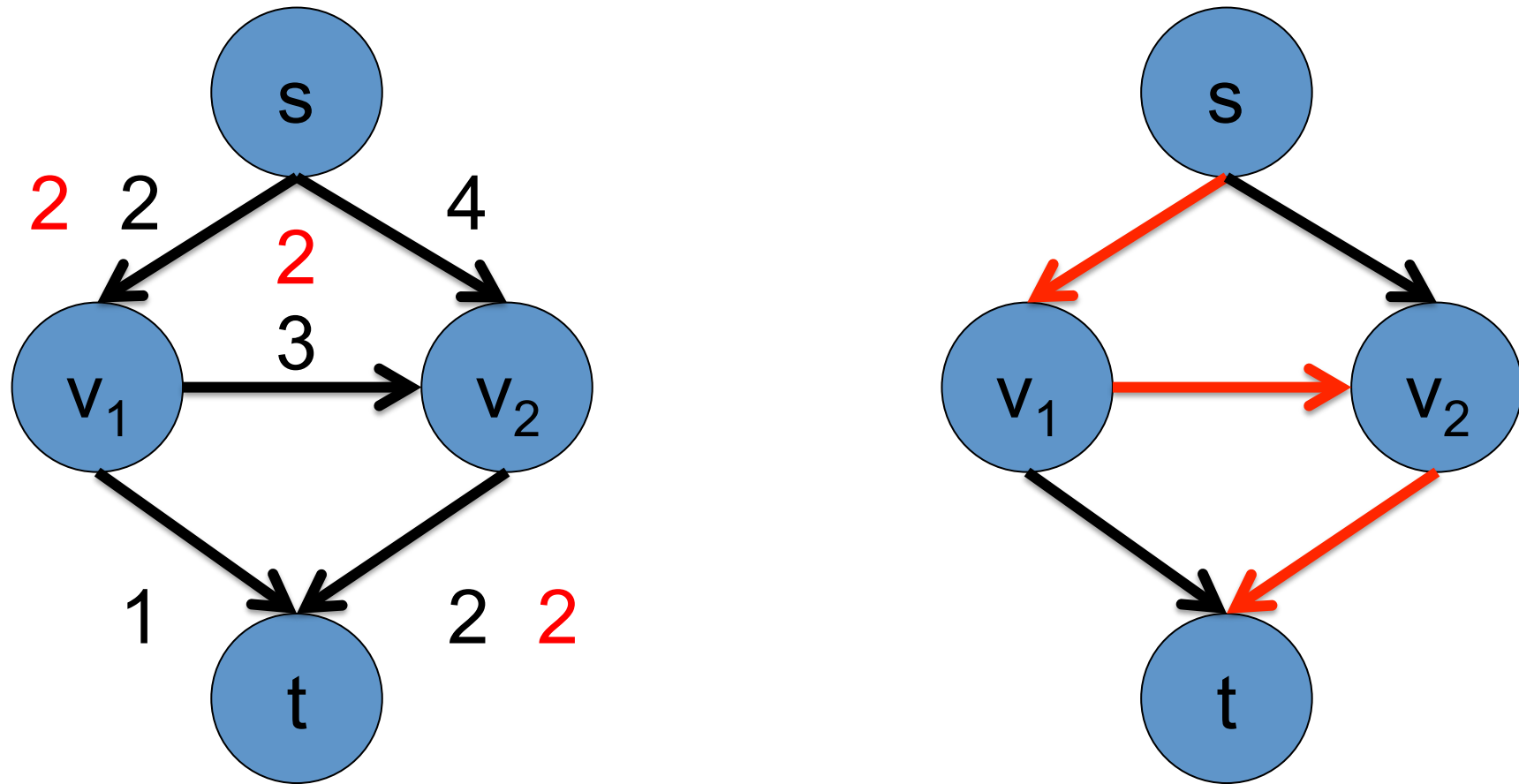
Find an s - t path in the residual graph.

Maximum Flow using Residual Graphs



For inverse arcs in path, subtract flow K .

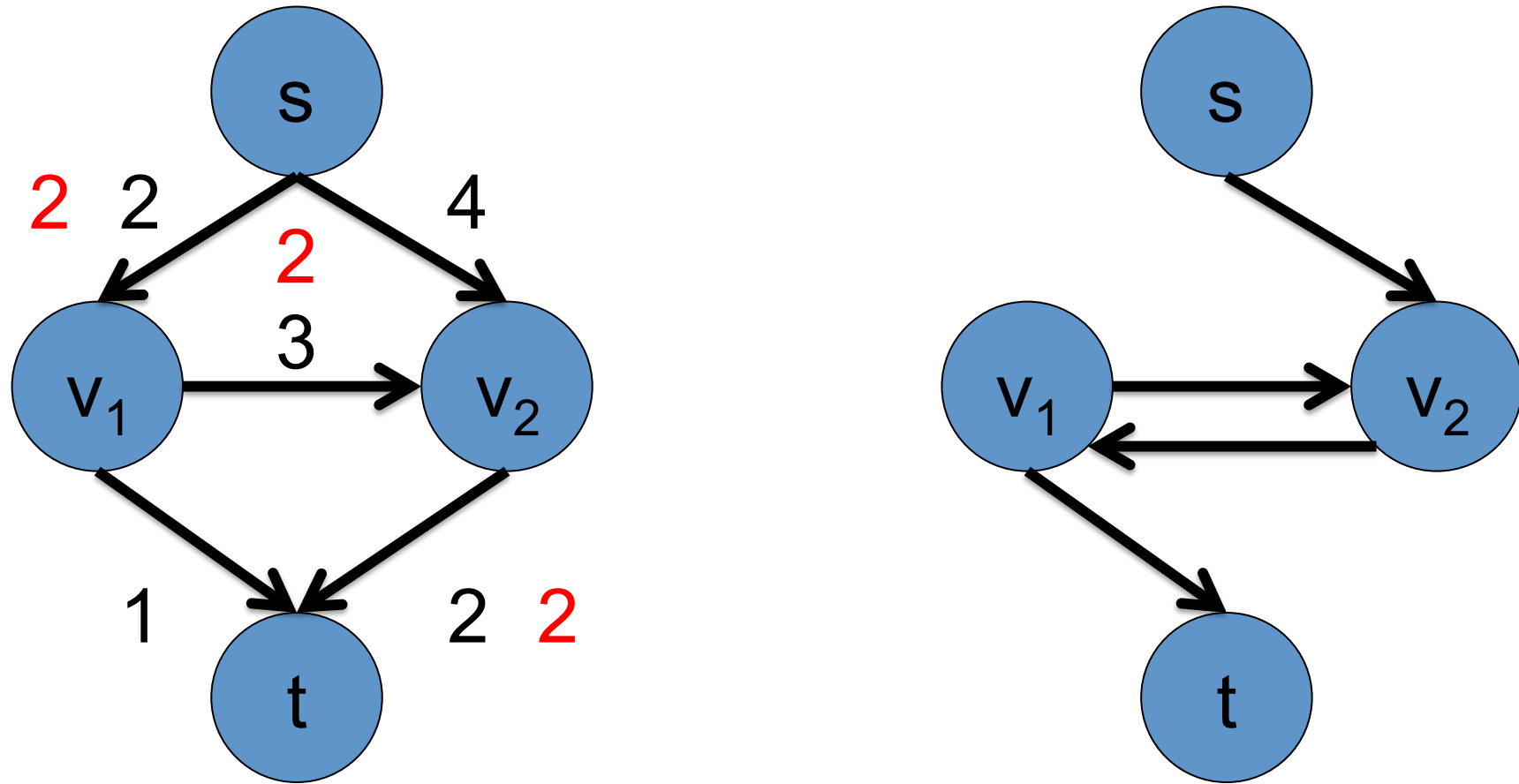
Maximum Flow using Residual Graphs



Choose maximum allowable value of K .

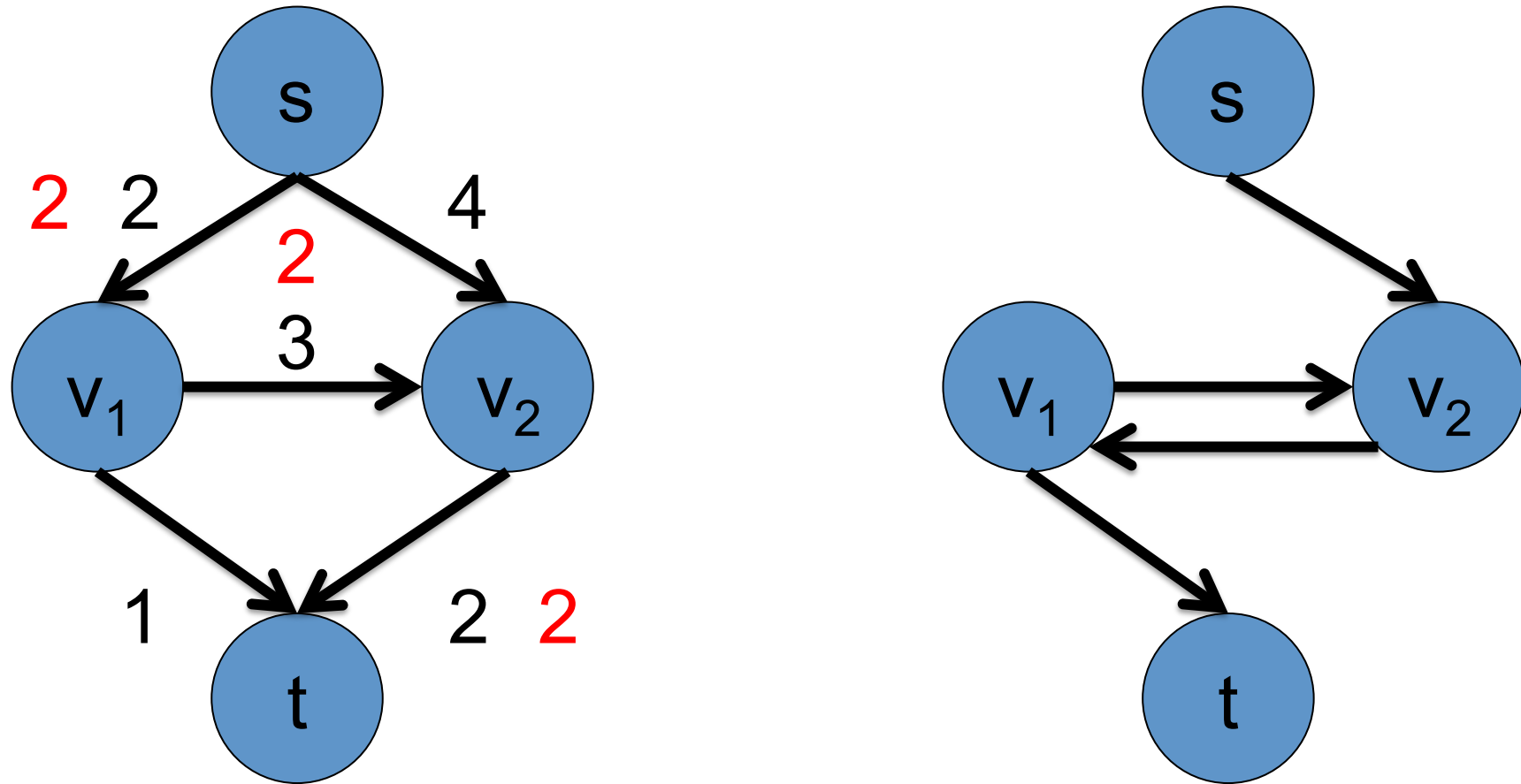
For forward arcs in path, add flow K .

Maximum Flow using Residual Graphs



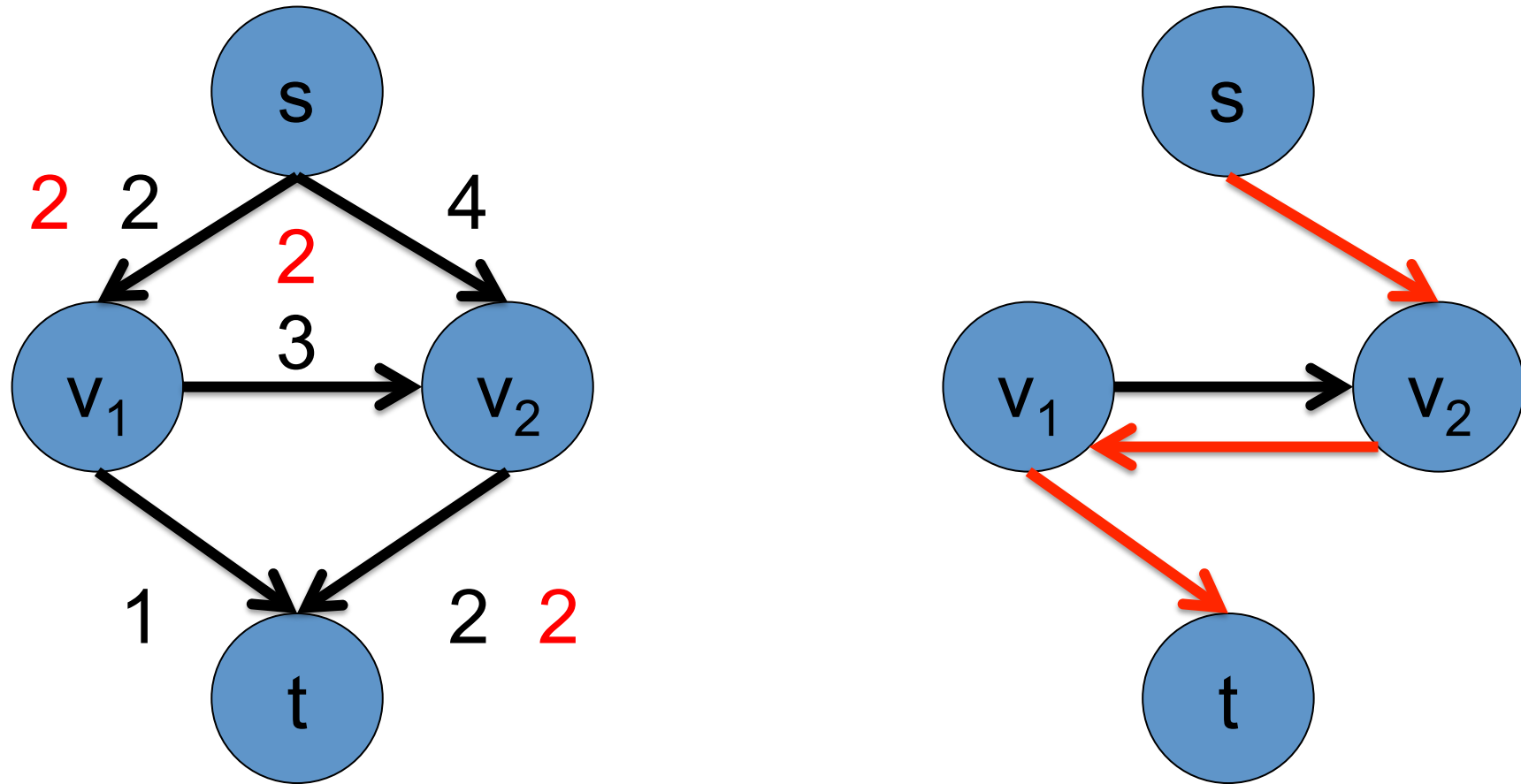
Update the residual graph.

Maximum Flow using Residual Graphs



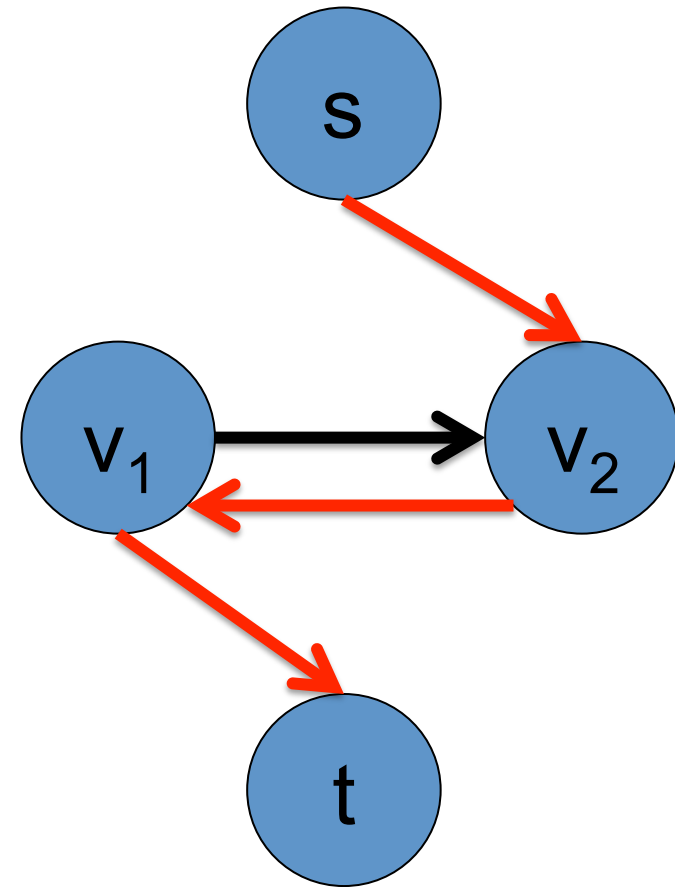
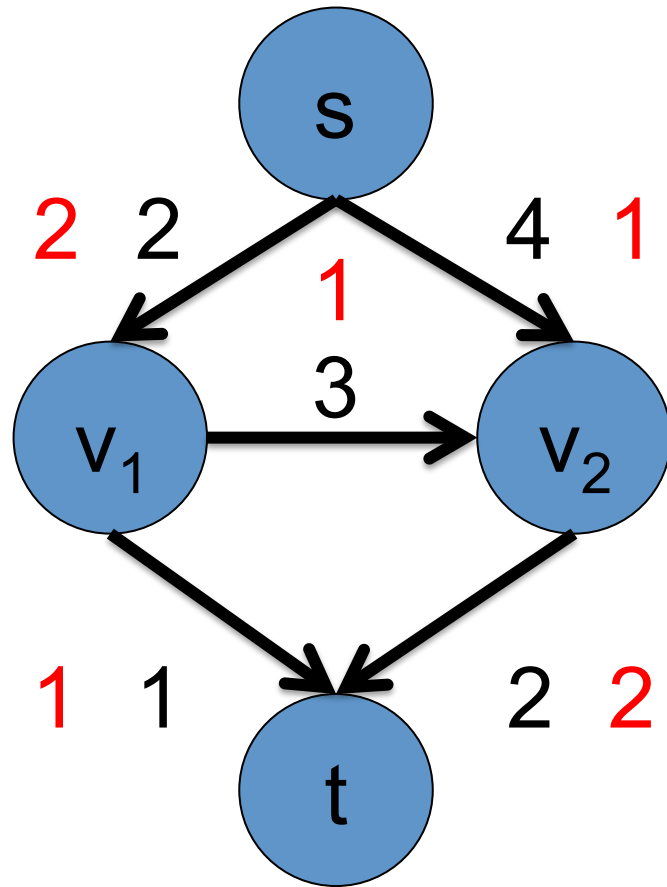
Find an s-t path in the residual graph.

Maximum Flow using Residual Graphs



Find an s - t path in the residual graph.

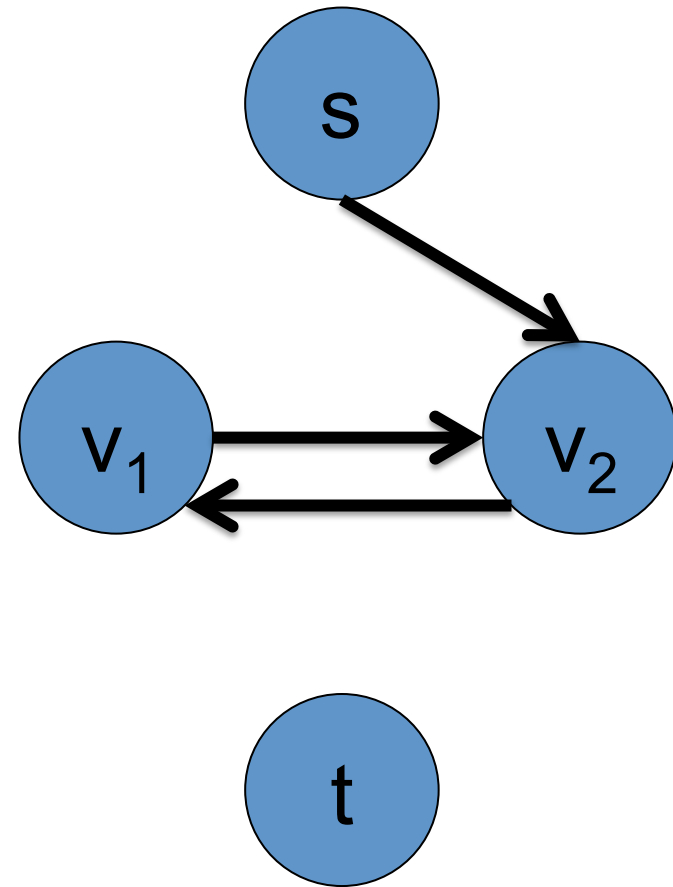
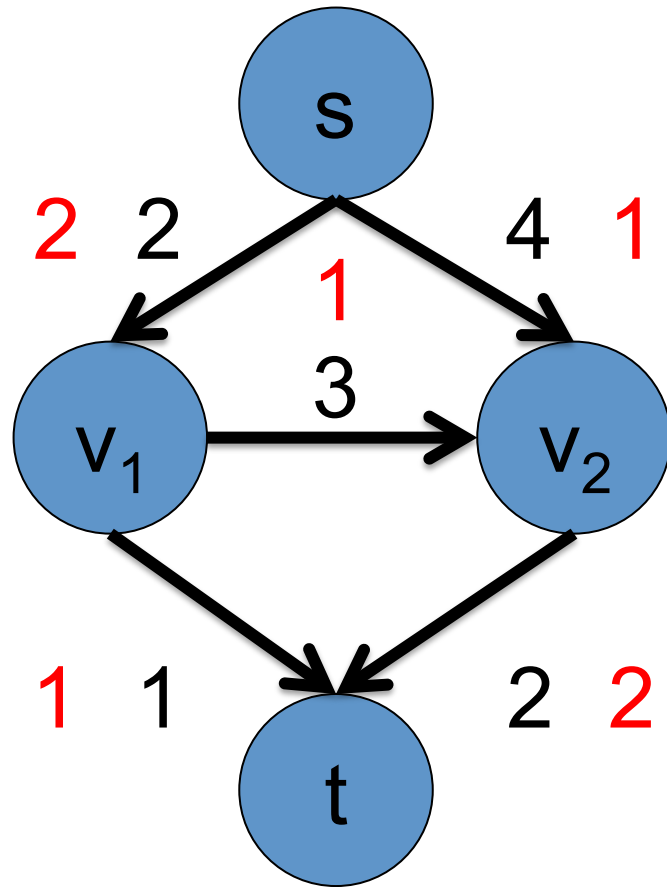
Maximum Flow using Residual Graphs



Choose maximum allowable value of K .

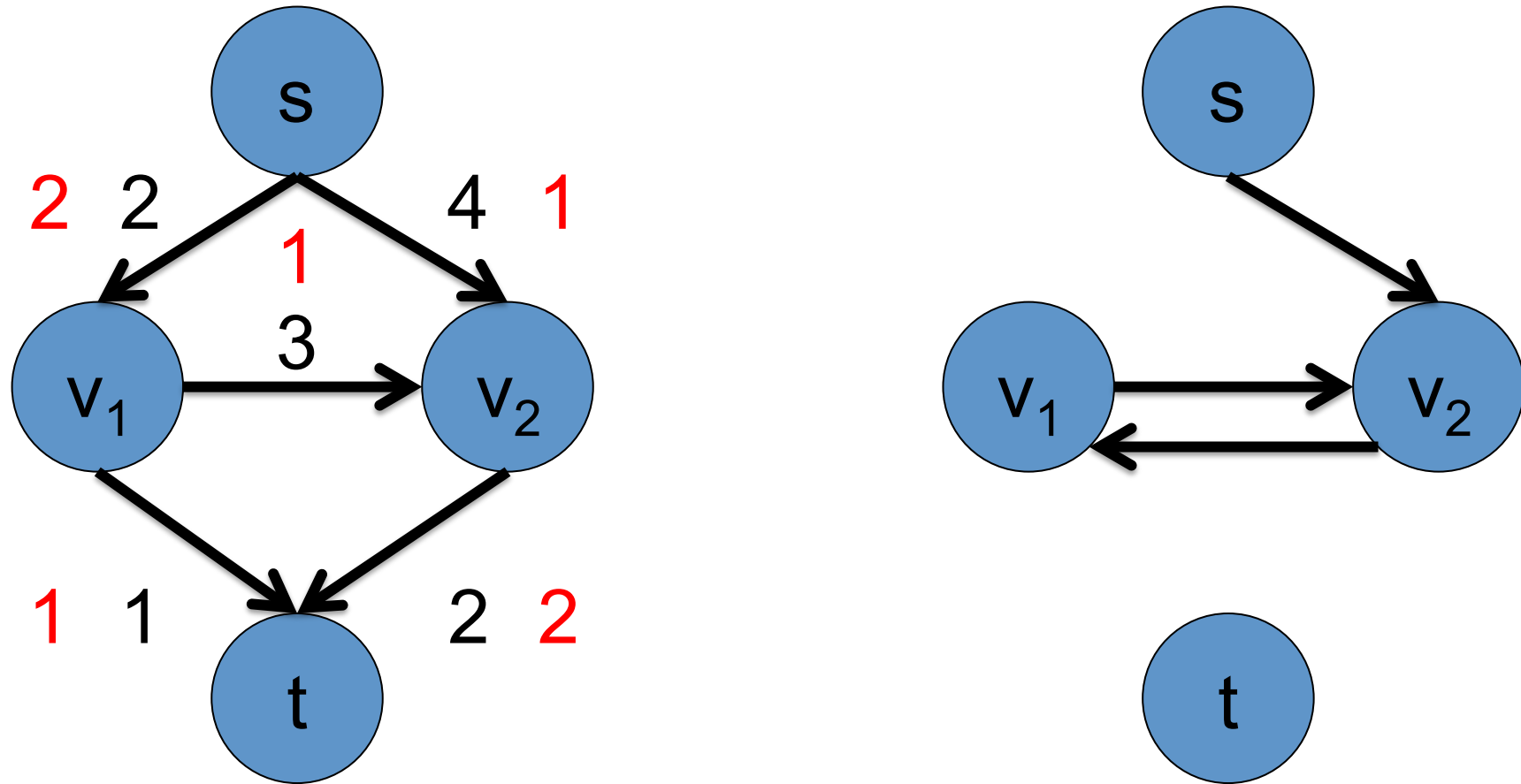
Add K to (s, v_2) and (v_1, t) . Subtract K from (v_1, v_2) .

Maximum Flow using Residual Graphs



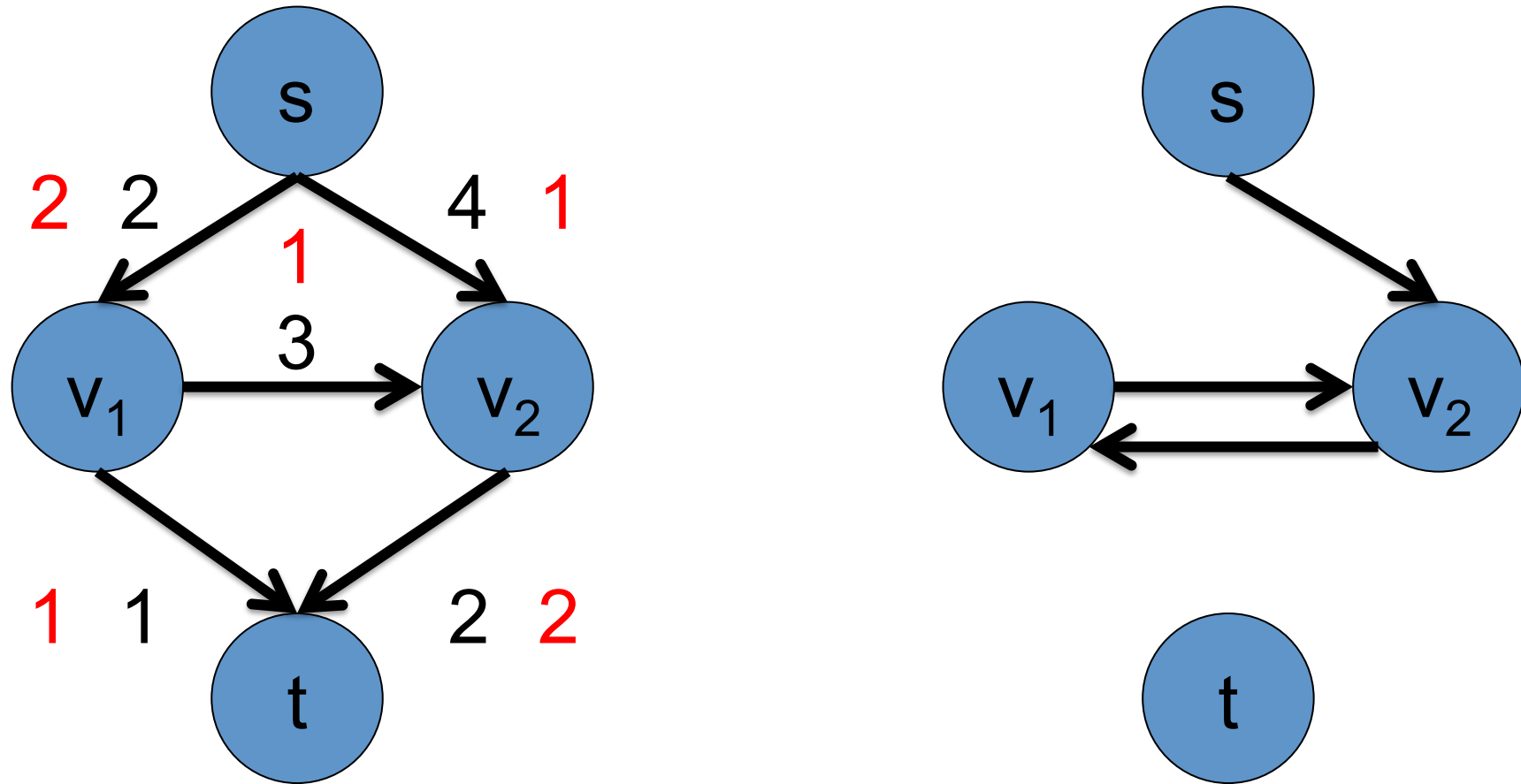
Update the residual graph.

Maximum Flow using Residual Graphs



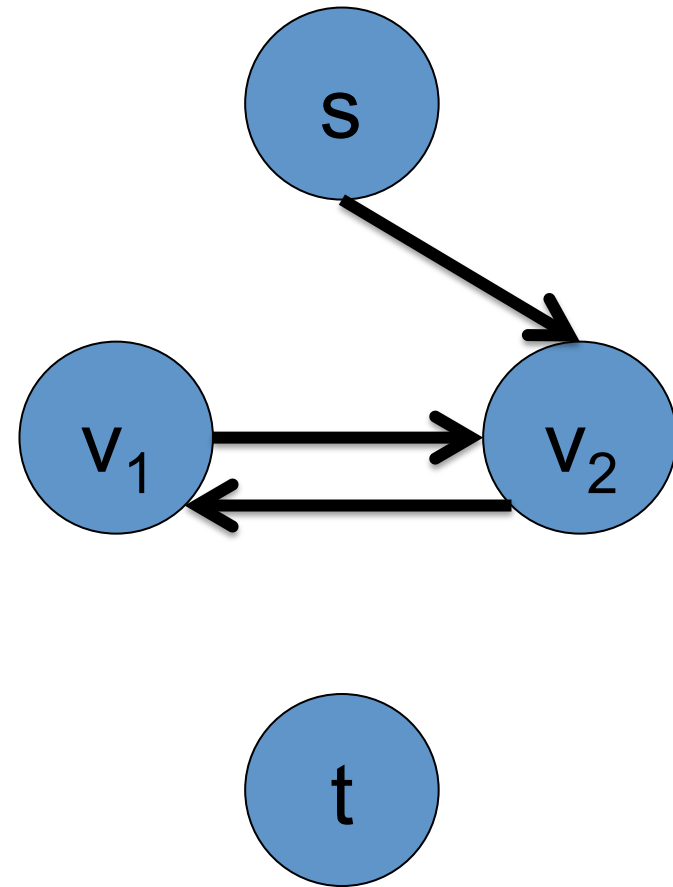
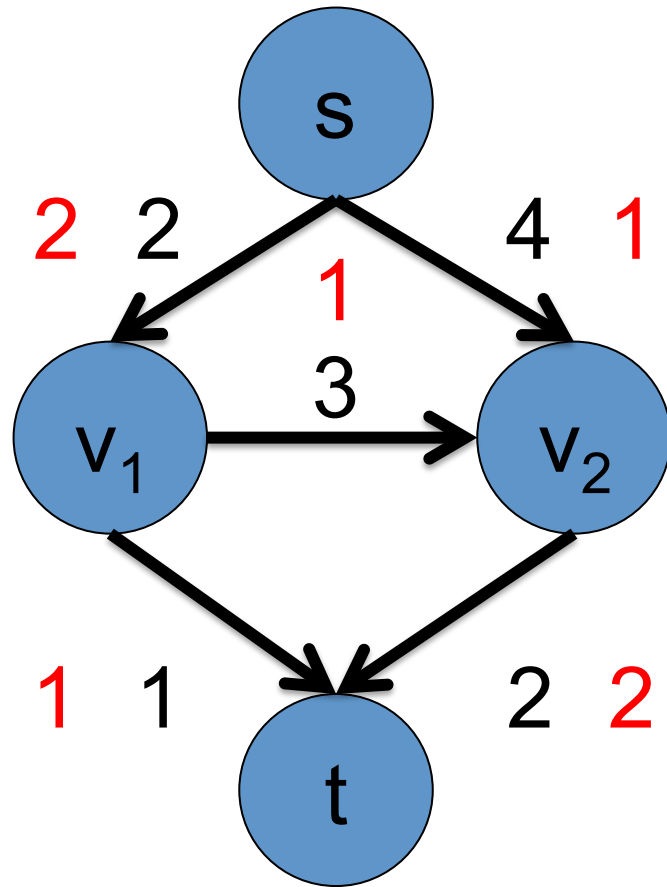
Find an s-t path in the residual graph.

Maximum Flow using Residual Graphs



No more s-t paths. Stop.

Maximum Flow using Residual Graphs



Correct Answer.