

# Discrete Inference and Learning

## Lecture 1

MVA

2017 – 2018

<http://thoth.inrialpes.fr/~alahari/disinflern>

Slides based on material from Stephen Gould, Pushmeet Kohli, Nikos Komodakis, M. Pawan Kumar, Carsten Rother

# Optimization problems

- Can be written as

$\min_x f(x)$  (optimize an objective function)

s.t.  $x \in \mathcal{C}$  (subject to some constraints)

**feasible set**, containing all  $x$   
satisfying the constraints

discrete variables

# Optimization problems

- Can be written as

$$\min_x f(x) \quad (\text{optimize an objective function})$$

$$\text{s.t. } x \in \mathcal{C} \quad (\text{subject to some constraints})$$

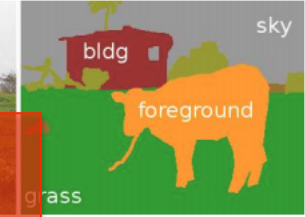
- Two main problems in this context
  - Optimize the objective (**inference**)
  - Learn the parameters of  $f$  (**learning**)

# Optimization problems

- Several applications, e.g., computer vision



Interactive figure-ground segmentation [Boykov and Jolly, 2001; Boykov and Funka-Lea, 2006]



Surface context [Hoiem et al., 2005]

Semantic labeling [He et al., 2004; Shotton et al., 2006; Gould et al., 2009]

## Low-level vision problems



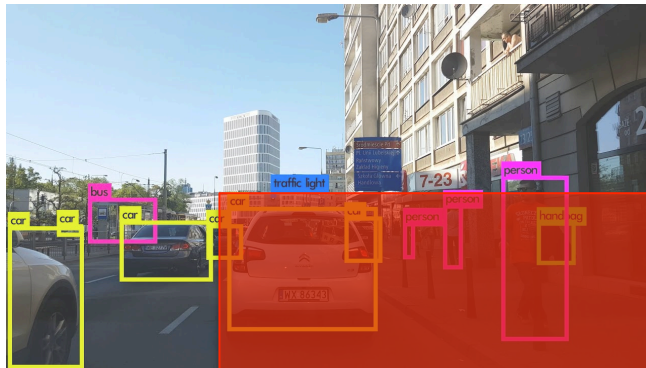
Stereo matching [Kolmogorov and Zabih, 2001; Scharstein and Szeliski, 2002]



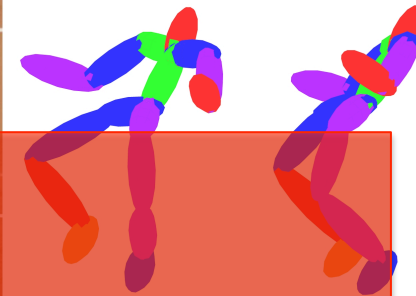
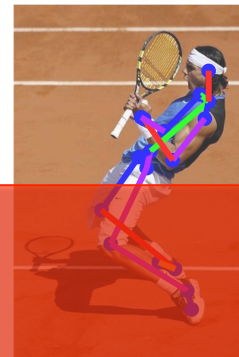
Image denoising [Felzenszwalb and Huttenlocher 2004]

# Optimization problems

- Several applications, e.g., computer vision

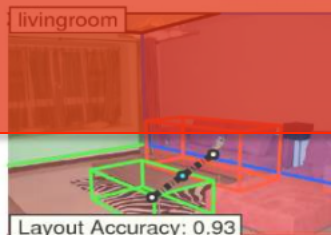
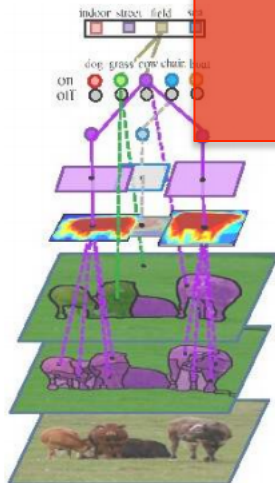


Object detection [Felzenszwalb et al., 2008]



Pose estimation [Alber and Black, 2015; Ramakrishna et al., 2012]

## High-level vision problems

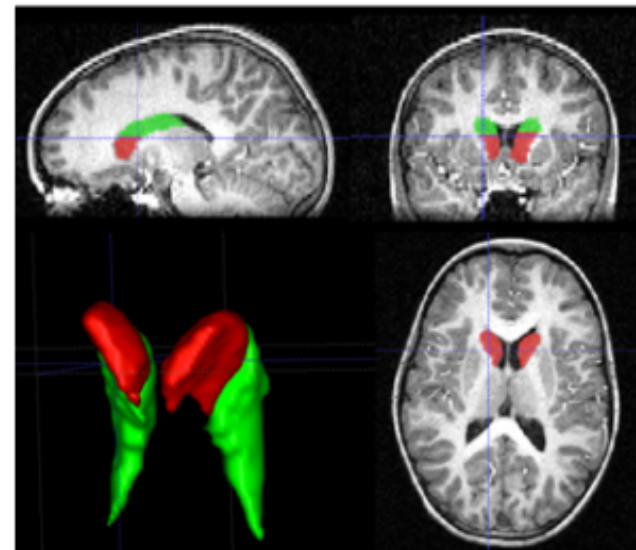
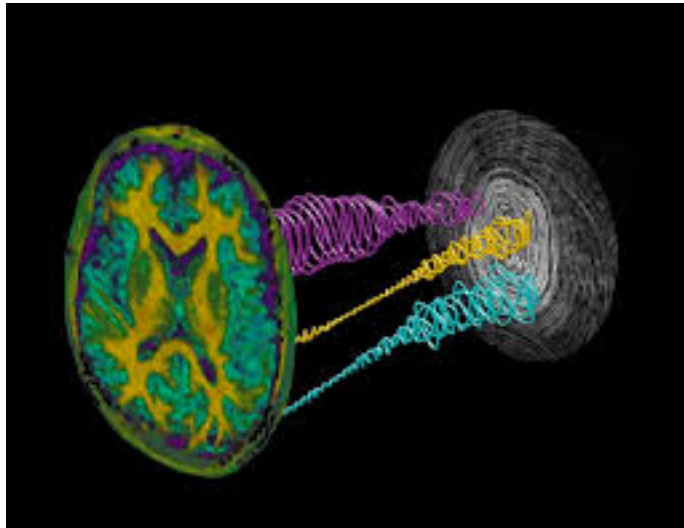


Scene understanding

[Fouhey et al., 2014; Ladicky et al., 2010; Xiao et al., 2013; Yao et al., 2012]

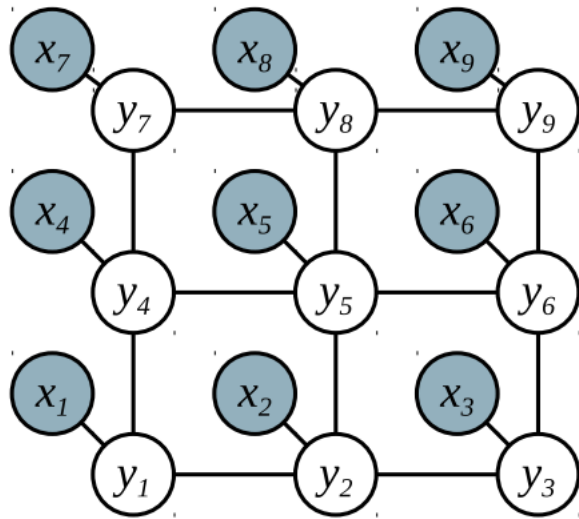
# Optimization problems

- Several applications, e.g., medical imaging

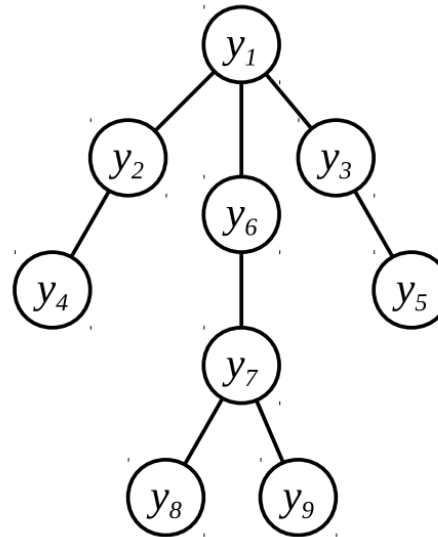


# Optimization problems

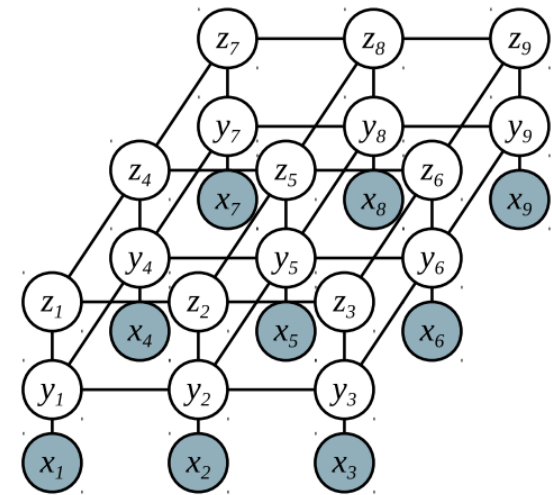
- Inherent in all these problems are graphical models



Pixel labeling



Object detection  
Pose estimation



Scene understanding

# Conditional Markov Random Fields

- Also known as: Markov networks, undirected graphical models, MRFs
- Note: Not making a distinction between CRFs and MRFs
- $\mathbf{X} \in \mathcal{X}$  : observed random variables
- $\mathbf{Y} = (Y_1, \dots, Y_n) \in \mathcal{Y}$ : output random variables
- $\mathbf{Y}_c$  are subset of variables for clique  $c \subseteq \{1, \dots, n\}$
- Define a factored probability distribution

$$P(\mathbf{Y} | \mathbf{X}) = \frac{1}{Z(\mathbf{X})} \prod_c \psi_c(\mathbf{Y}_c; \mathbf{X})$$

Partition function =  $\sum_{\mathbf{Y} \in \mathcal{Y}} \prod_c \psi_c(\mathbf{Y}_c; \mathbf{X})$  **Exponential number of configurations !**



# Maximum a posteriori (MAP) inference

$$\begin{aligned} \mathbf{y}^* &= \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} P(\mathbf{y} \mid \mathbf{x}) \\ &= \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} \frac{1}{Z(\mathbf{x})} \prod_c \psi_c(\mathbf{Y}_c; \mathbf{X}) \\ &= \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} \log \left( \frac{1}{Z(\mathbf{x})} \prod_c \psi_c(\mathbf{Y}_c; \mathbf{X}) \right) \\ &= \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} \sum_c \log \psi_c(\mathbf{Y}_c; \mathbf{X}) - \log Z(\mathbf{X}) \\ &= \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} \sum_c \log \psi_c(\mathbf{Y}_c; \mathbf{X}) \rightarrow -E(\mathbf{Y}; \mathbf{X}) \end{aligned}$$

# Maximum a posteriori (MAP) inference

$$\begin{aligned}\mathbf{y}^* &= \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} P(\mathbf{y} \mid \mathbf{x}) = \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} \sum_c \log \Psi_c(\mathbf{Y}_c; \mathbf{X}) \\ &= \operatorname{argmin}_{\mathbf{y} \in \mathcal{Y}} E(\mathbf{y}; \mathbf{x})\end{aligned}$$

MAP inference  $\Leftrightarrow$  Energy minimization

The energy function is  $E(\mathbf{Y}; \mathbf{X}) = \sum_c \psi_c(\mathbf{Y}_c; \mathbf{X})$

where  $\psi_c(\cdot) = -\log \Psi_c(\cdot)$

 Clique potential

# Clique potentials

- Defines a mapping from an assignment of random variables to a real number

$$\psi_c : \mathcal{Y}_c \times \mathcal{X} \rightarrow \mathbb{R}$$

- Encodes a preference for assignments to the random variables (lower is better)

- Parameterized as  $\psi_c(\mathbf{y}_c; \mathbf{x}) = \mathbf{w}_c^T \phi_c(\mathbf{y}_c; \mathbf{x})$

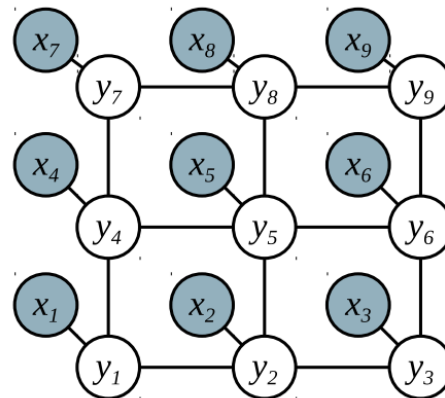


Parameters

# Clique potentials

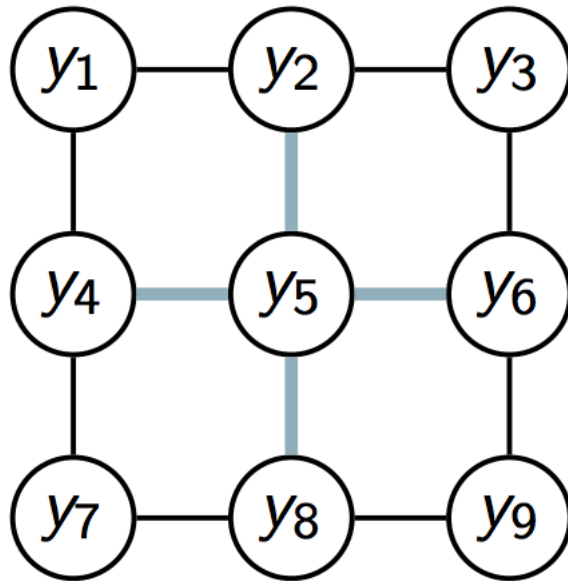
- Arity

$$\begin{aligned} E(\mathbf{y}; \mathbf{x}) &= \sum_c \psi_c(\mathbf{y}_c; \mathbf{x}) \\ &= \underbrace{\sum_{i \in \mathcal{V}} \psi_i^U(y_i; \mathbf{x})}_{\text{unary}} + \underbrace{\sum_{ij \in \mathcal{E}} \psi_{ij}^P(y_i, y_j; \mathbf{x})}_{\text{pairwise}} + \underbrace{\sum_{c \in \mathcal{C}} \psi_c^H(\mathbf{y}_c; \mathbf{x})}_{\text{higher-order}}. \end{aligned}$$

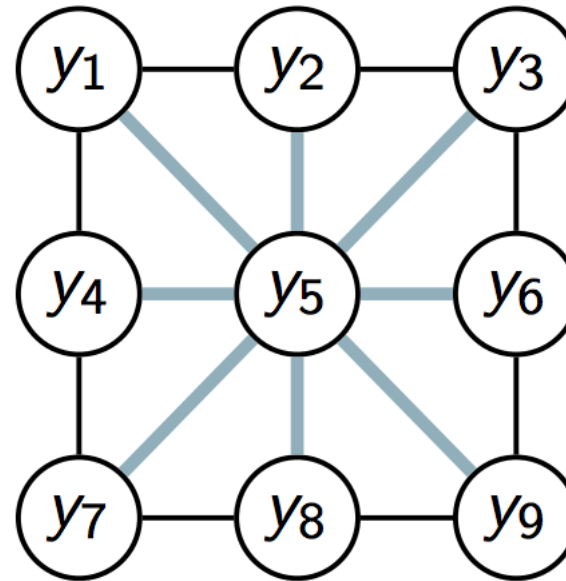


# Clique potentials

- Arity

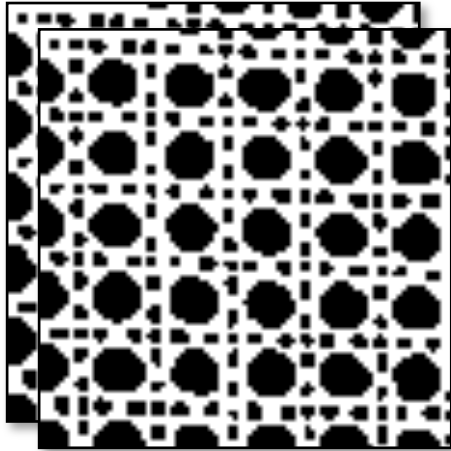


4-connected,  $\mathcal{N}_4$

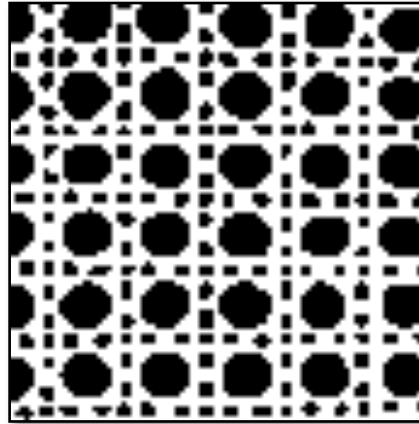


8-connected,  $\mathcal{N}_8$

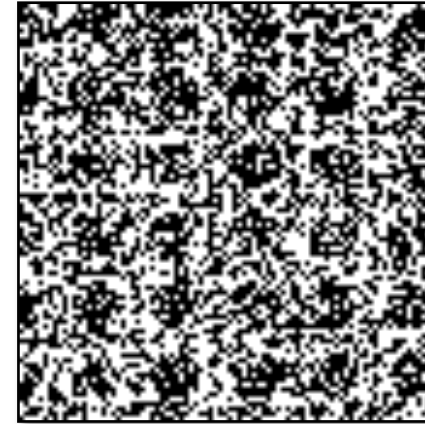
# Reason 1: Texture modelling



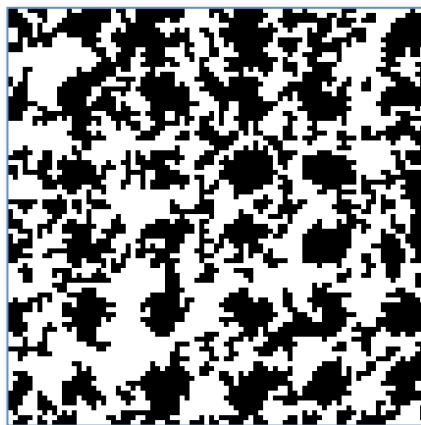
Training images



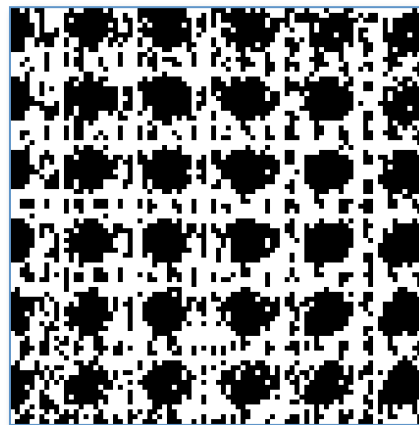
Test image



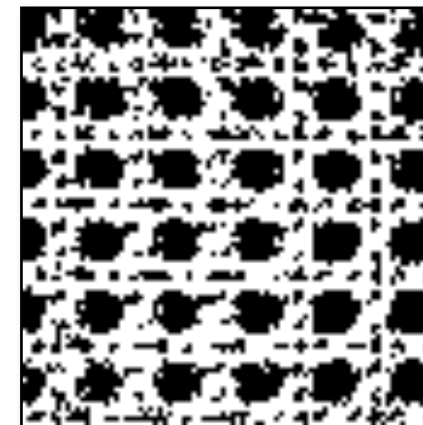
Test image (60% Noise)



Result MRF  
4-connected  
(neighbours)

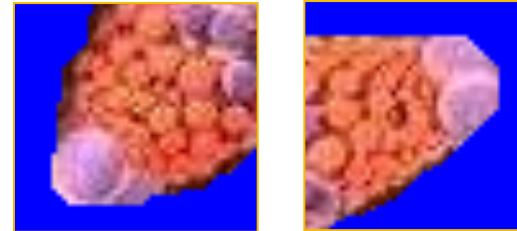
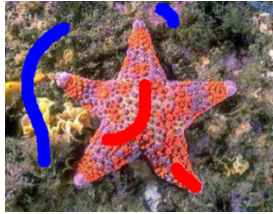


Result MRF  
4-connected

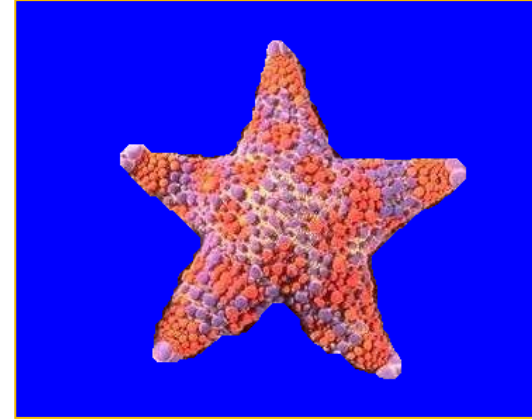


Result MRF  
9-connected  
(7 attractive; 2 repulsive)

# Reason2: Discretization artefacts



4-connected  
Euclidean

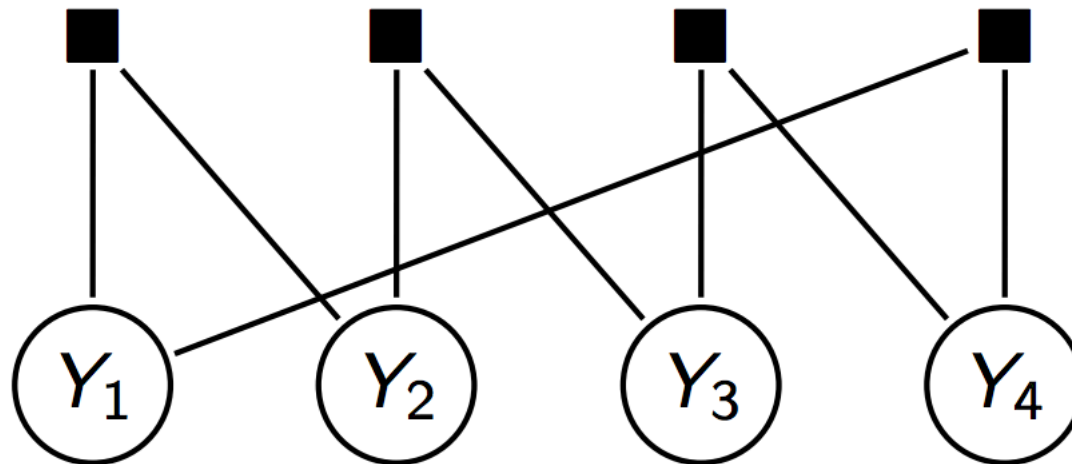


8-connected  
Euclidean

# Graphical representation

- Example

$$E(\mathbf{y}) = \psi(y_1, y_2) + \psi(y_2, y_3) + \psi(y_3, y_4) + \psi(y_4, y_1)$$



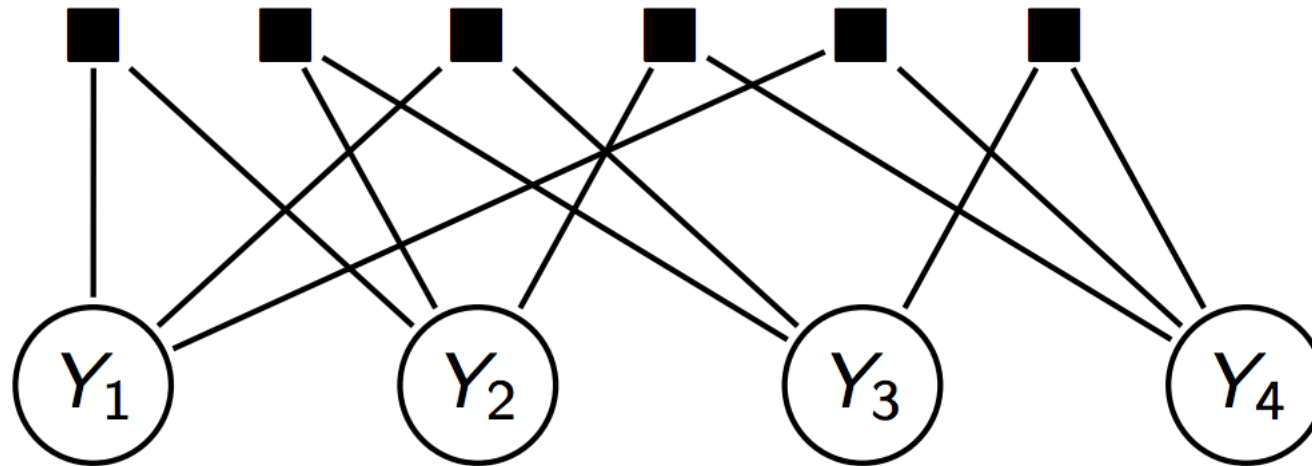
factor graph



# Graphical representation

- Example

$$E(\mathbf{y}) = \sum_{i,j} \psi(y_i, y_j)$$

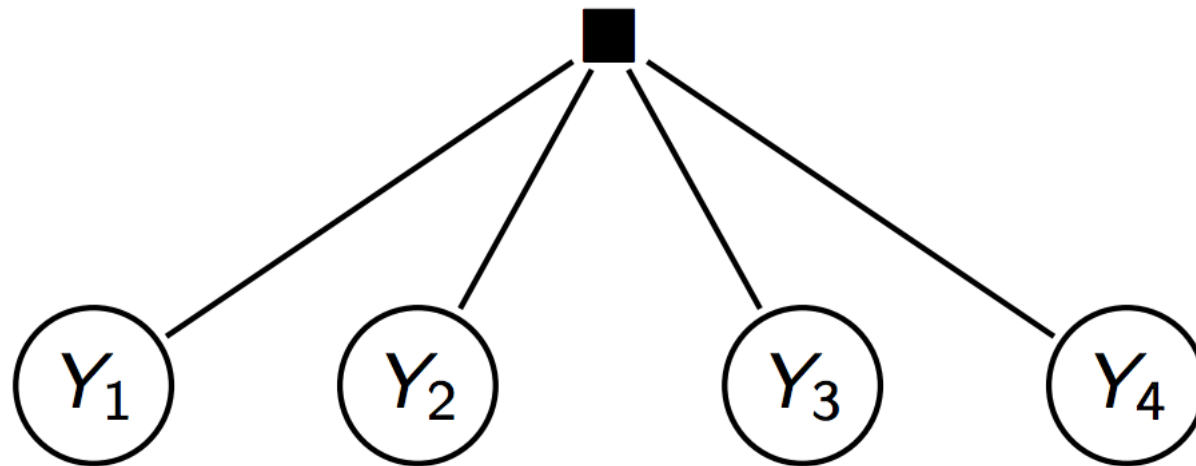


factor graph

# Graphical representation

- Example

$$E(\mathbf{y}) = \psi(y_1, y_2, y_3, y_4)$$



factor graph

# A Computer Vision Application

## Binary Image Segmentation



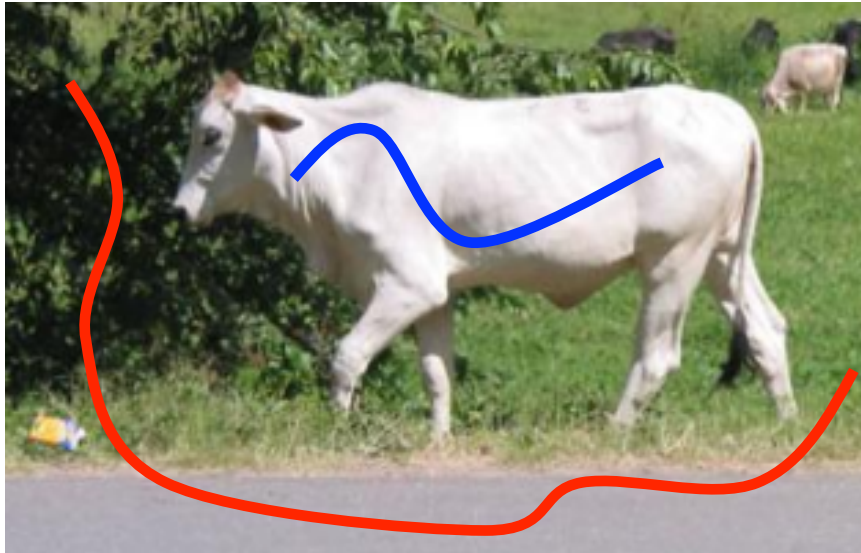
**How ?**

Cost function    Models *our* knowledge about natural images

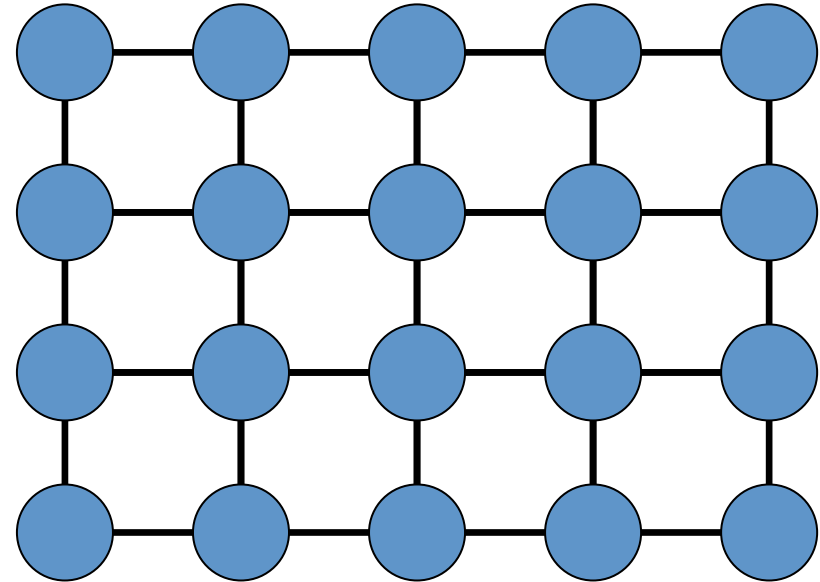
Optimize cost function to obtain the segmentation

# A Computer Vision Application

## Binary Image Segmentation



Object - white, Background - green/grey



Graph  $G = (V, E)$

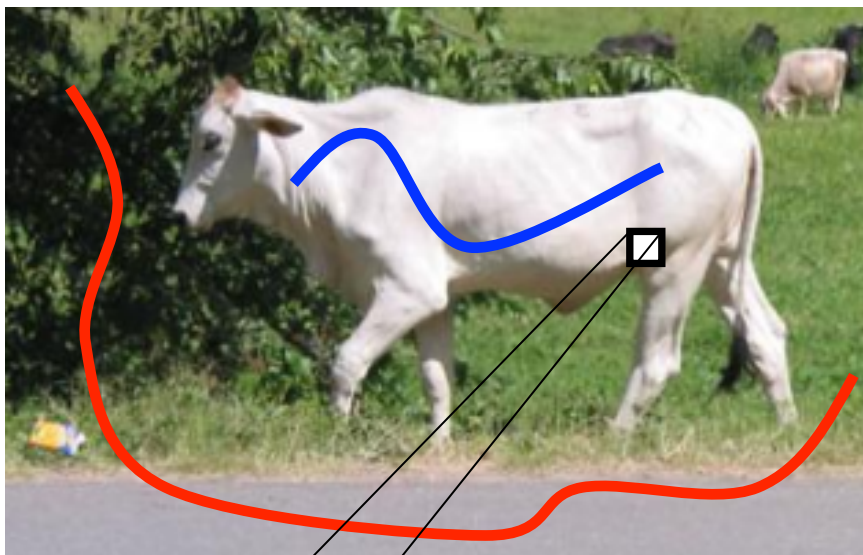
Each vertex corresponds to a pixel

Edges define a 4-neighbourhood *grid* graph

Assign a label to each vertex from  $L = \{\text{obj}, \text{bkg}\}$

# A Computer Vision Application

## Binary Image Segmentation

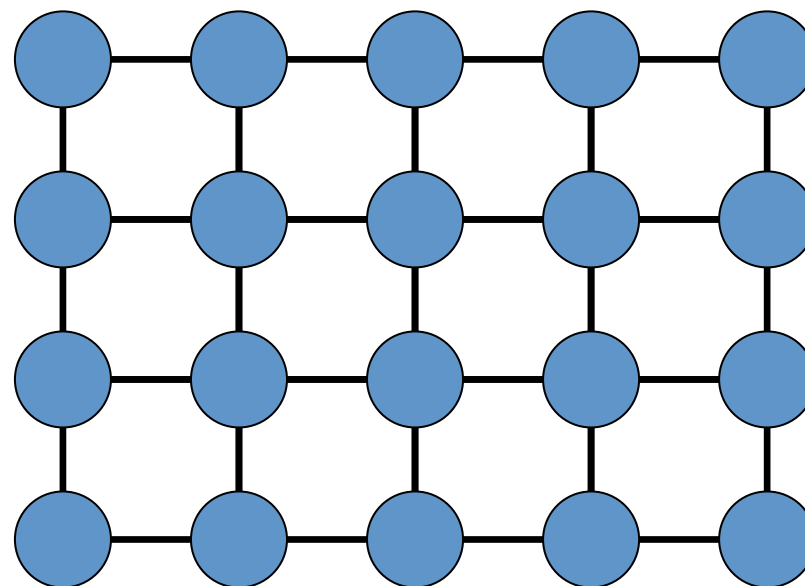


Object - white, Background - green/grey

Cost of a labelling  $f : V \rightarrow L$



Cost of label 'obj' low Cost of label 'bkg' high

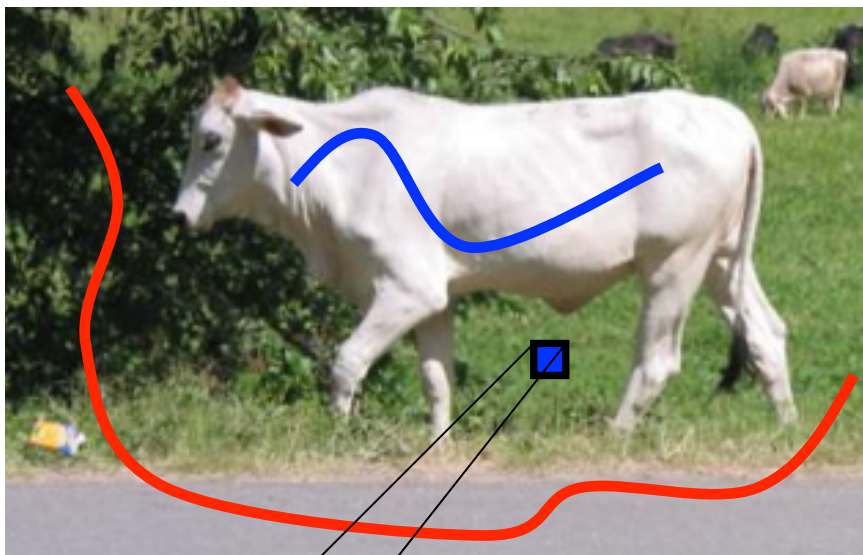


Graph  $G = (V, E)$

Per Vertex Cost

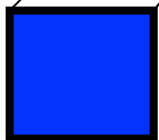
# A Computer Vision Application

## Binary Image Segmentation

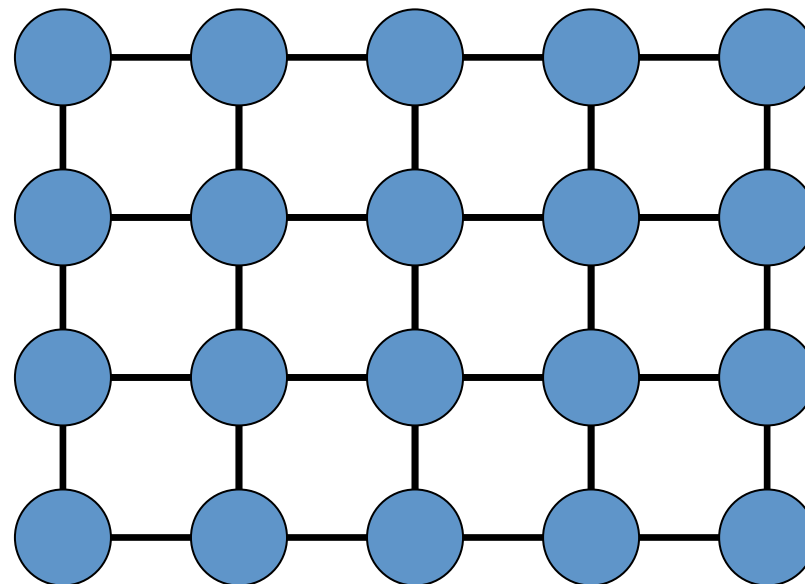


Object - white, Background - green/grey

Cost of a labelling  $f : V \rightarrow L$



Cost of label 'obj' high Cost of label 'bkg' low



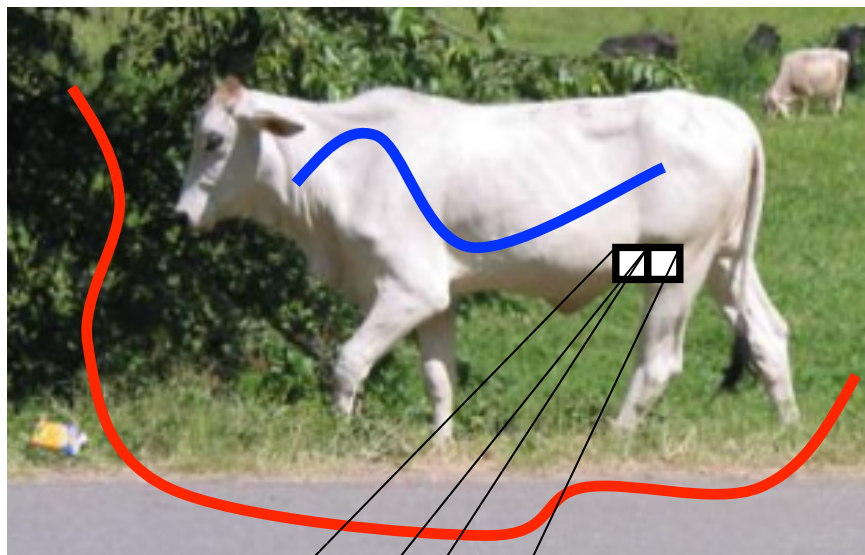
Graph  $G = (V, E)$

Per Vertex Cost

UNARY COST

# A Computer Vision Application

## Binary Image Segmentation



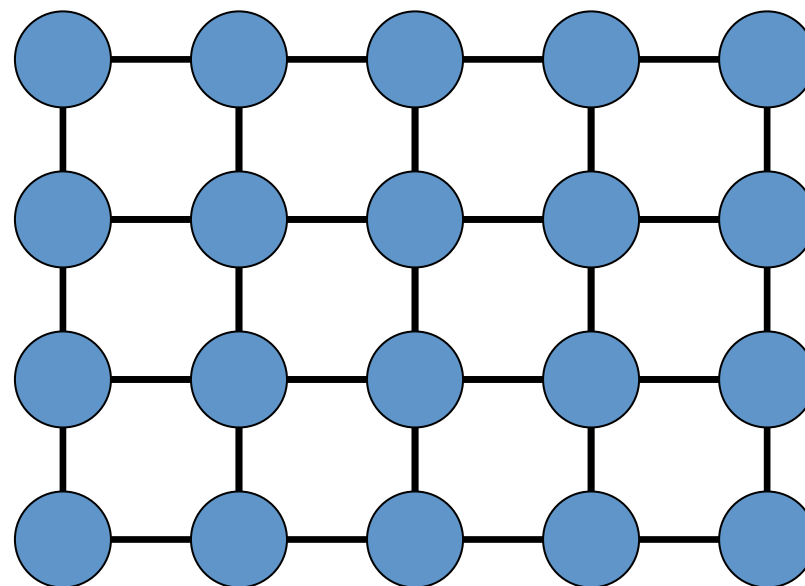
Object - white, Background - green/grey

Cost of a labelling  $f : V \rightarrow L$



Cost of same label low

Cost of different labels high

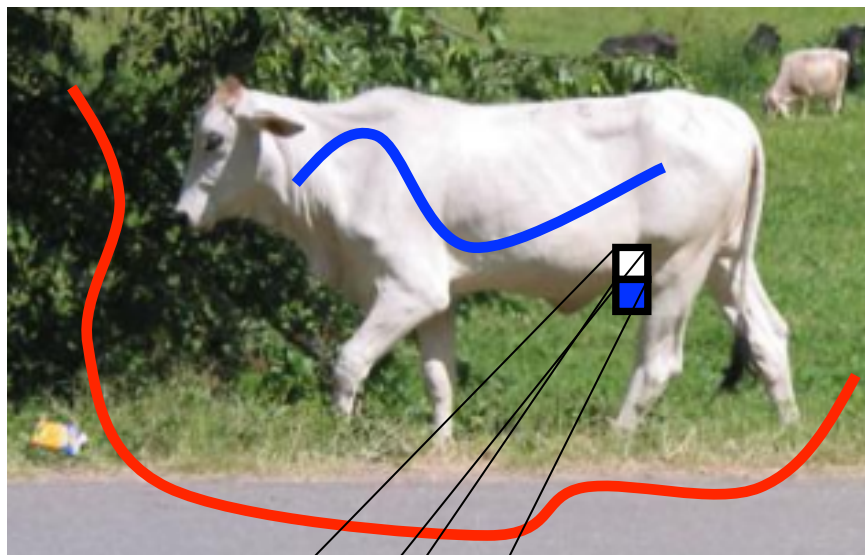


Graph  $G = (V, E)$

Per Edge Cost

# A Computer Vision Application

## Binary Image Segmentation



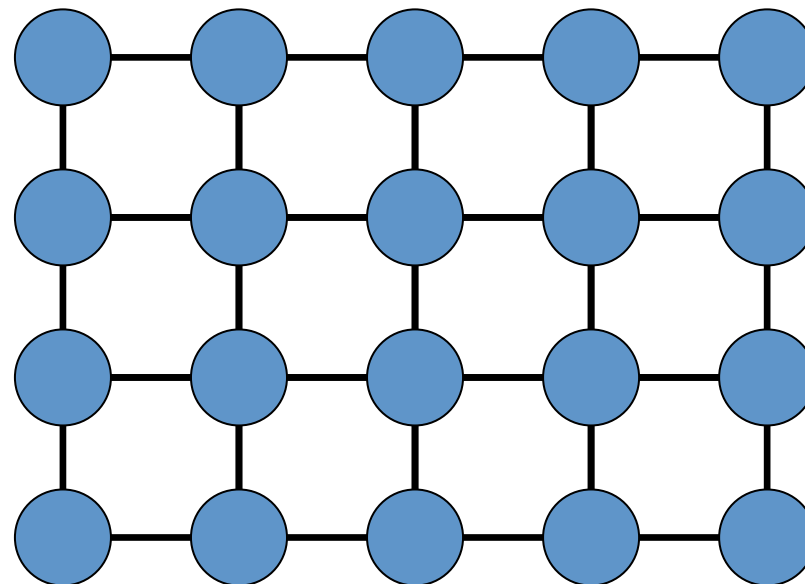
Object - white, Background - green/grey

Cost of a labelling  $f : V \rightarrow L$



Cost of same label high

Cost of different labels low



Graph  $G = (V, E)$

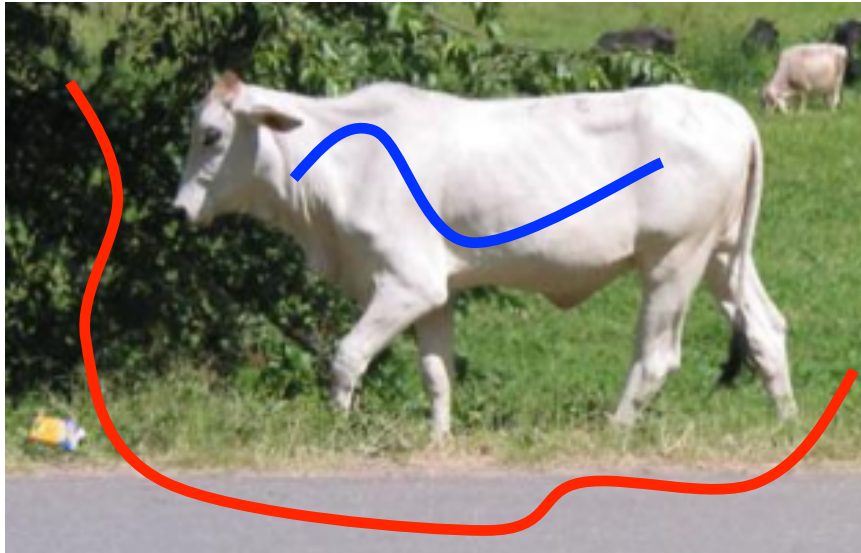
Per Edge Cost

**PAIRWISE  
COST**

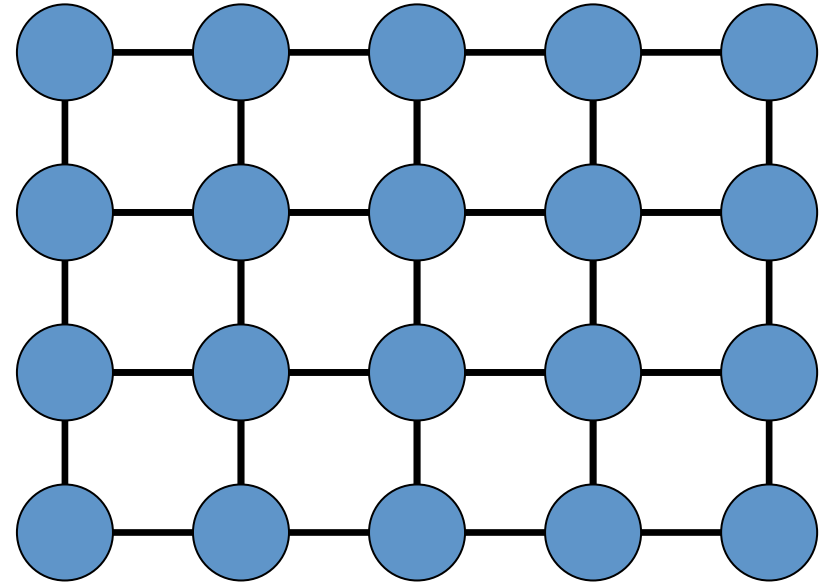


# A Computer Vision Application

## Binary Image Segmentation



Object - white, Background - green/grey

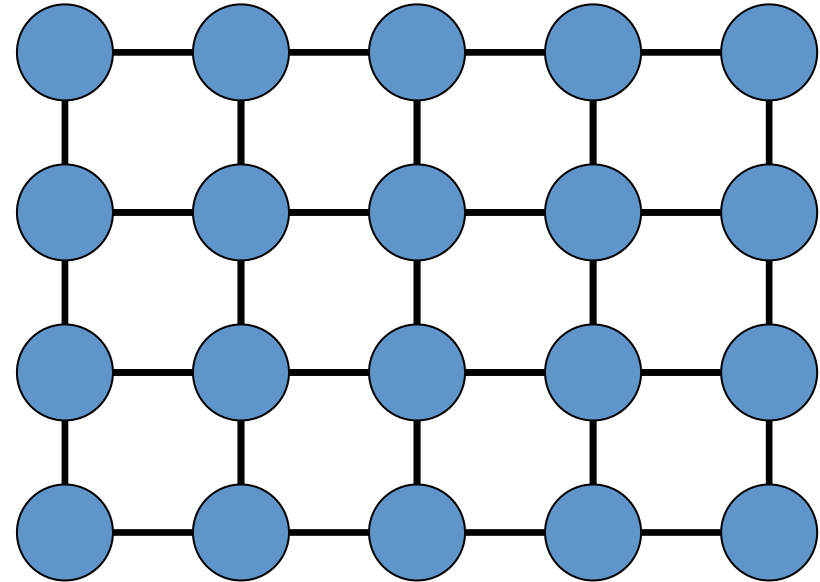


Graph  $G = (V, E)$

Problem: Find the labelling with minimum cost  $f^*$

# A Computer Vision Application

## Binary Image Segmentation

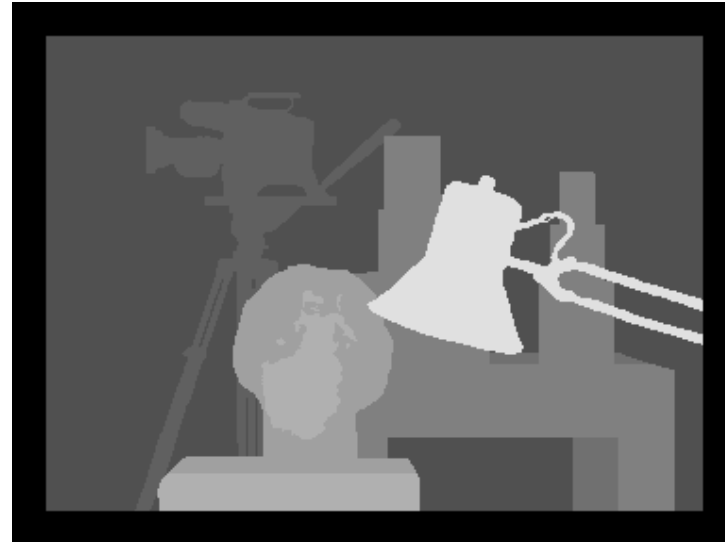
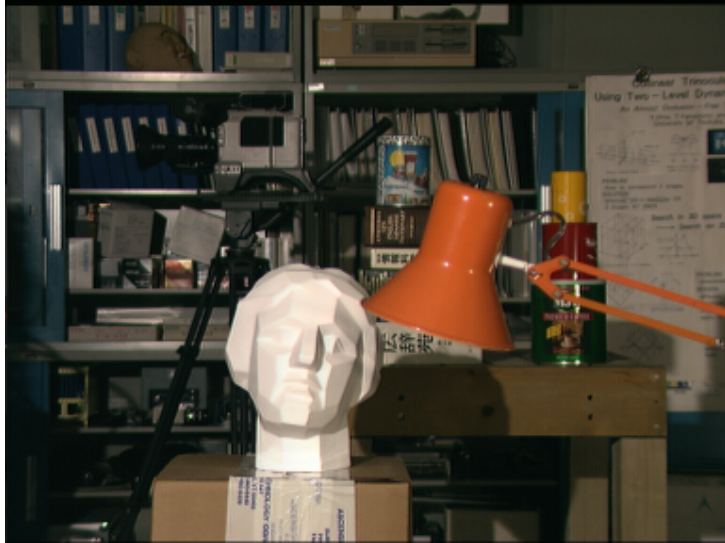


Graph  $G = (V, E)$

Problem: Find the labelling with minimum cost  $f^*$

# Another Computer Vision Application

## Stereo Correspondence



Disparity Map

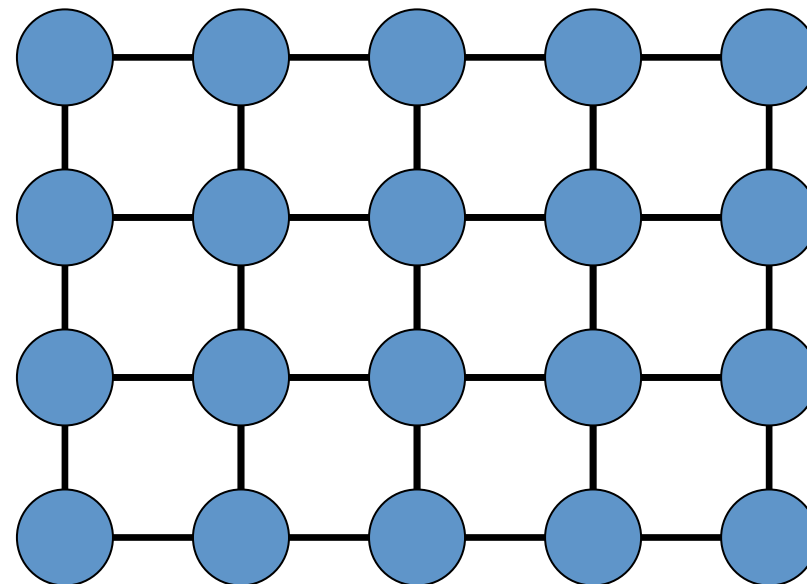
**How ?**

Minimizing a cost function



# Another Computer Vision Application

## Stereo Correspondence



Graph  $G = (V, E)$

Vertex corresponds to a pixel

Edges define grid graph

$L = \{\text{disparities}\}$

# Another Computer Vision Application

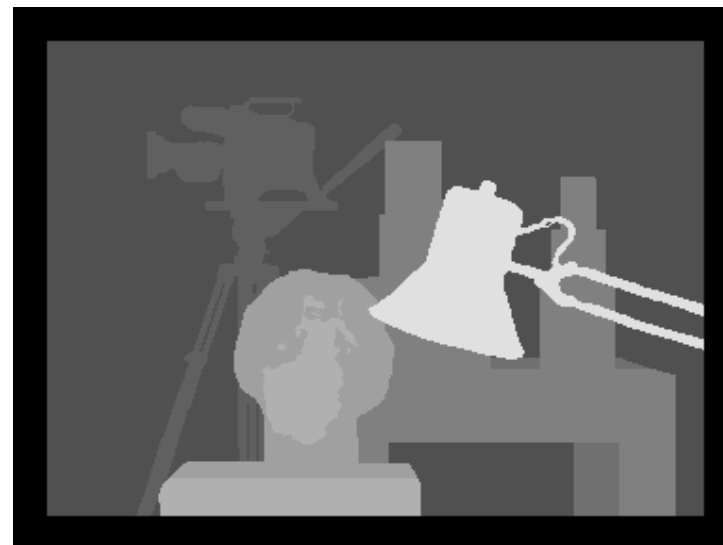
## Stereo Correspondence



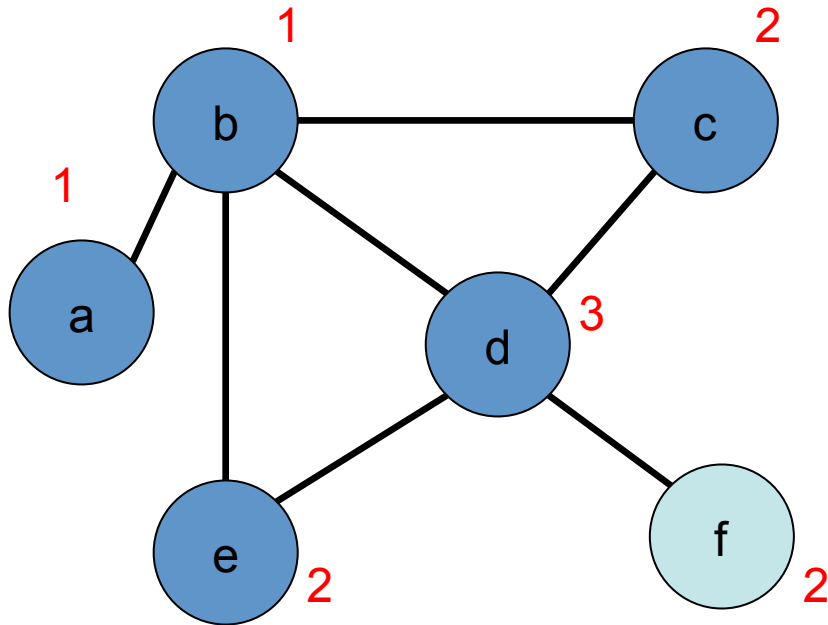
Cost of labelling  $f$ :

Unary cost + Pairwise Cost

Find minimum cost  $f^*$



# The General Problem



Graph  $G = (V, E)$

Discrete label set  $L = \{1, 2, \dots, h\}$

Assign a label to each vertex

$f: V \rightarrow L$

Cost of a labelling  $Q(f)$

Unary Cost

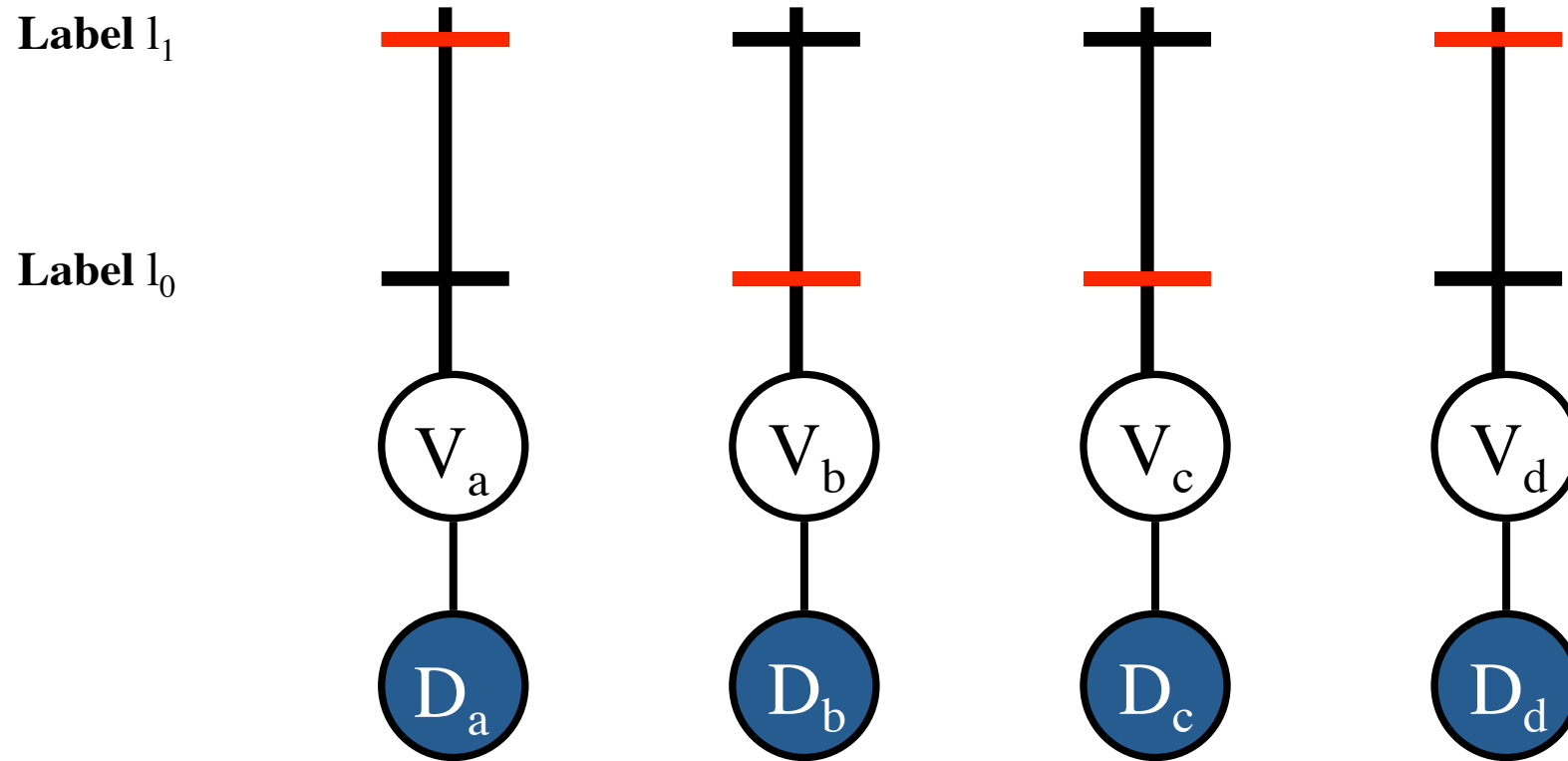
Pairwise Cost

Find  $f^* = \arg \min Q(f)$

# Overview

- Basics: problem formulation
  - Energy Function
  - MAP Estimation
  - Computing min-marginals
  - Reparameterization
- Solutions
  - Belief Propagation and related methods [Lecture 1]
  - Graph cuts [Lecture 2]

# Energy Function



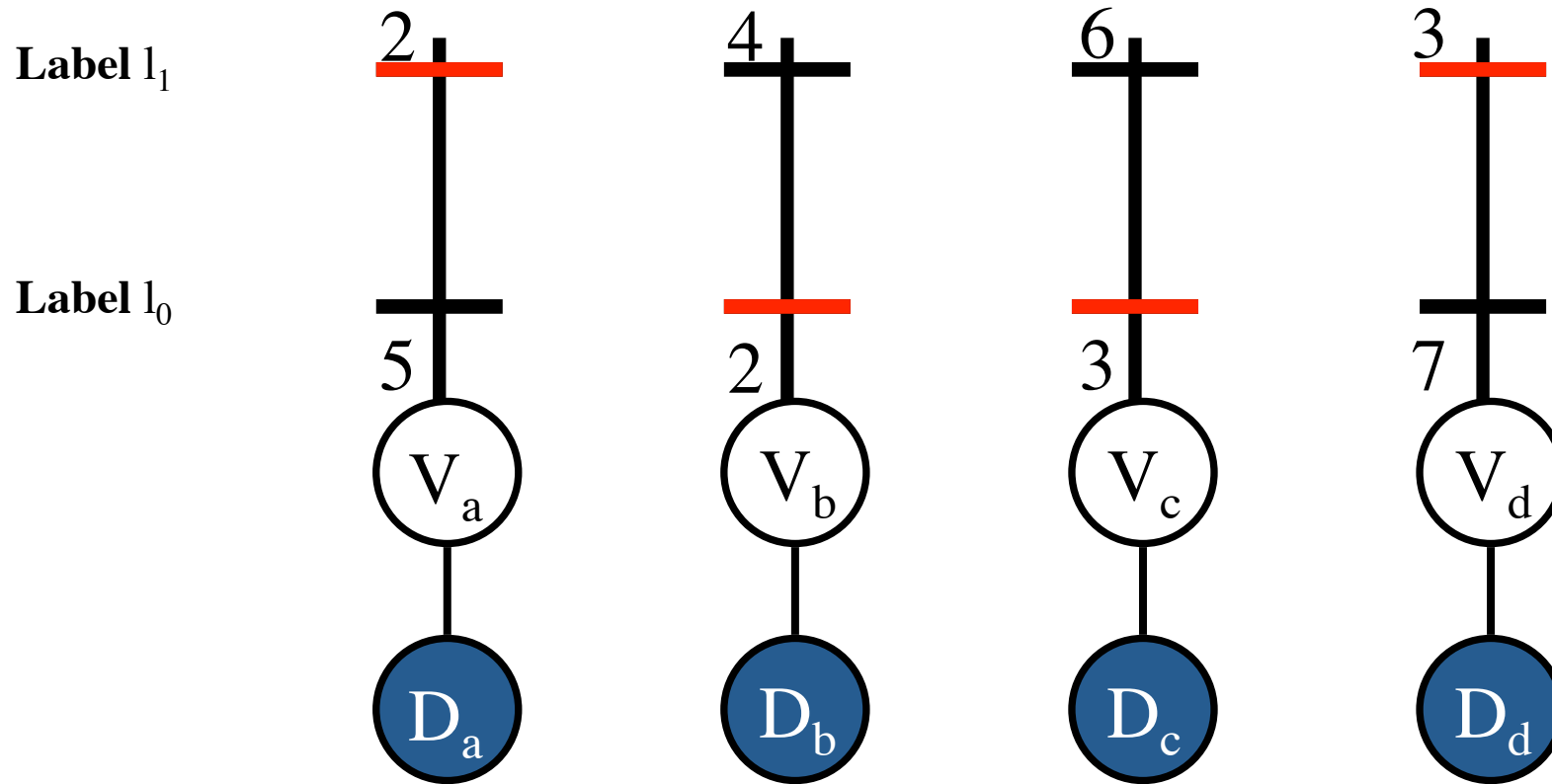
Random Variables  $V = \{V_a, V_b, \dots\}$

Labels  $L = \{l_0, l_1, \dots\}$  Data  $D$

Labelling  $f: \{a, b, \dots\} \rightarrow \{0, 1, \dots\}$



# Energy Function



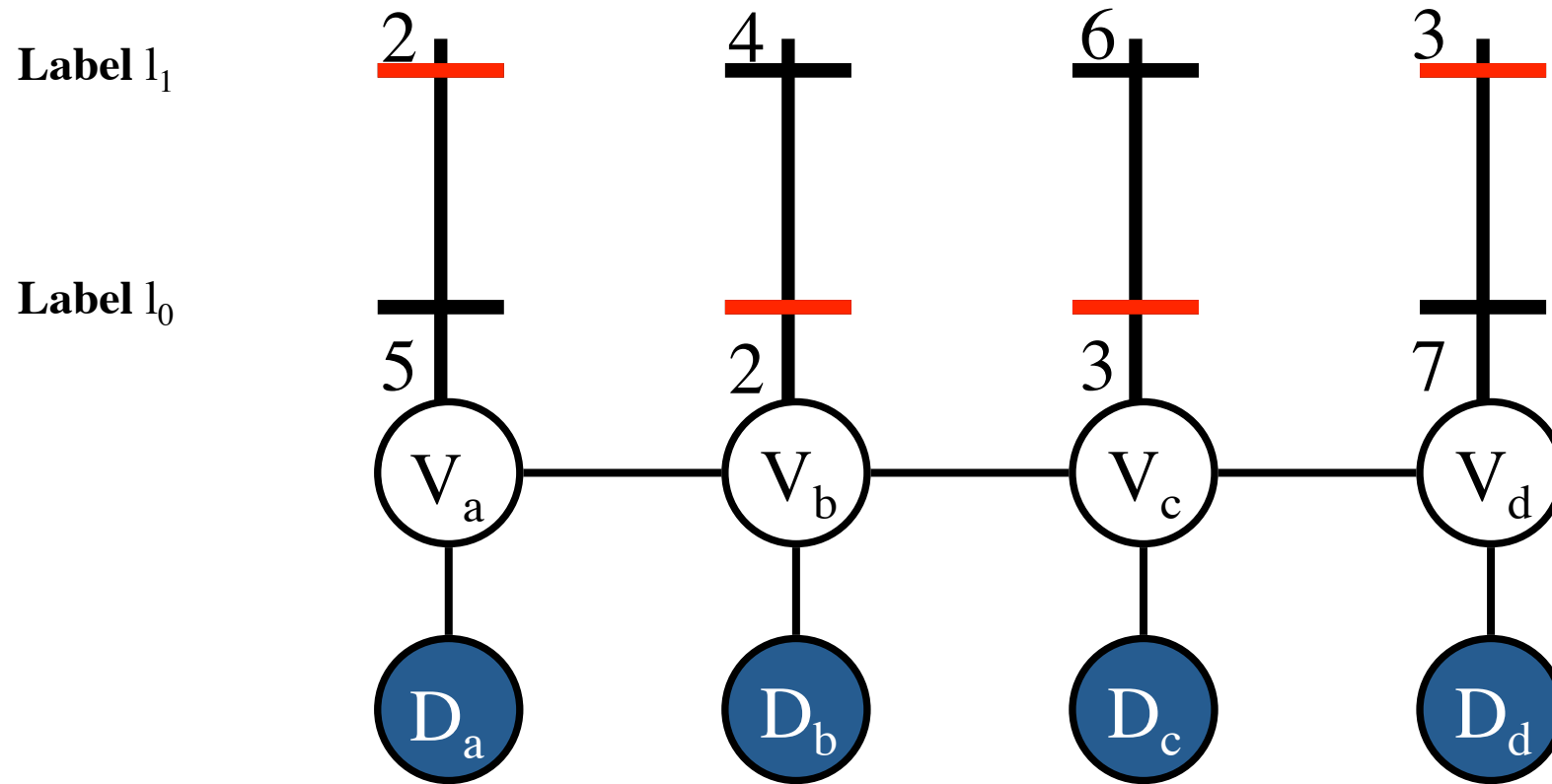
$$Q(f) = \sum_a \theta_{a;f(a)}$$

Unary Potential

Easy to minimize

Neighbourhood

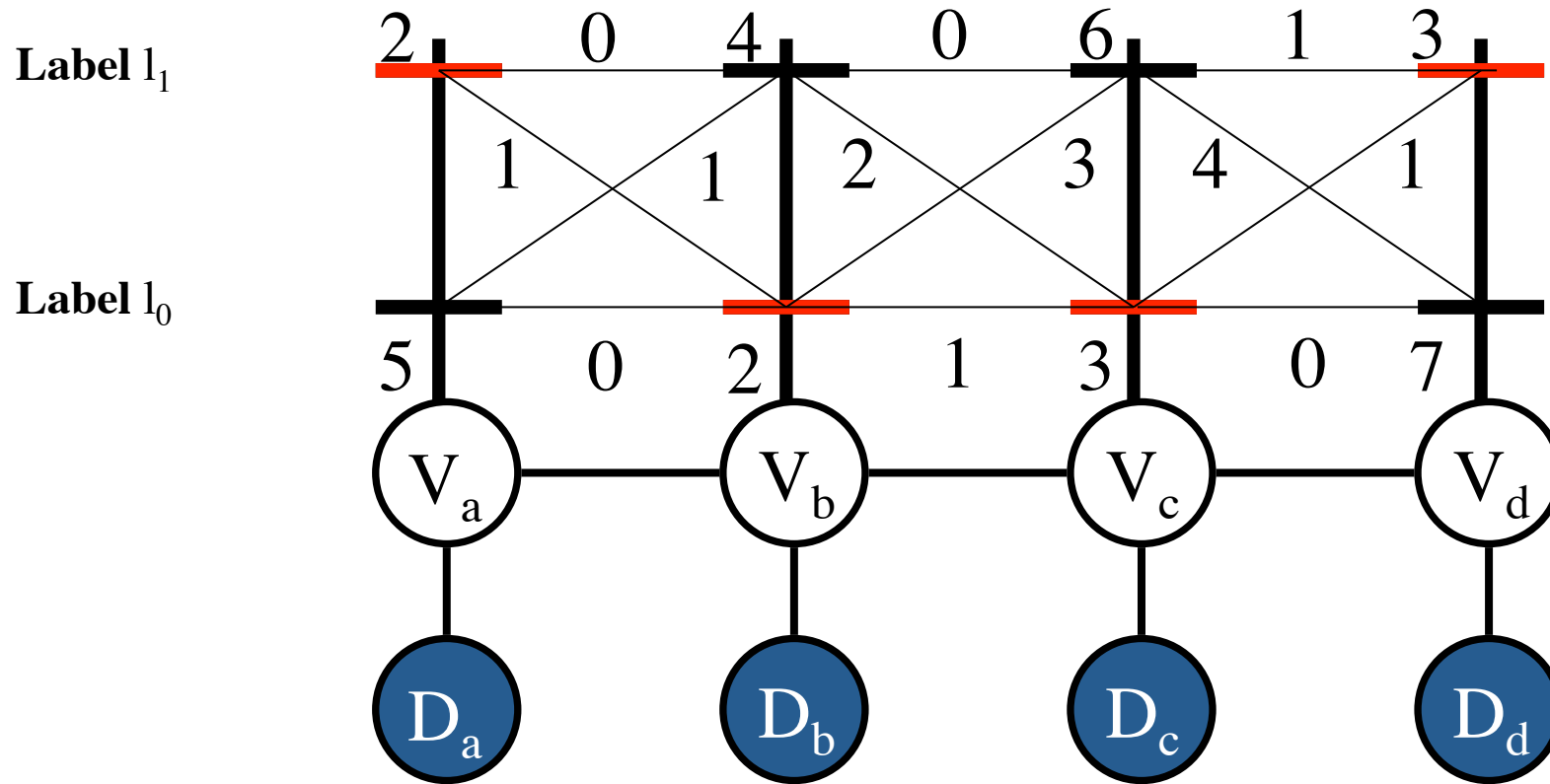
# Energy Function



$E : (a,b) \in E$  iff  $V_a$  and  $V_b$  are neighbours

$$E = \{ (a,b) , (b,c) , (c,d) \}$$

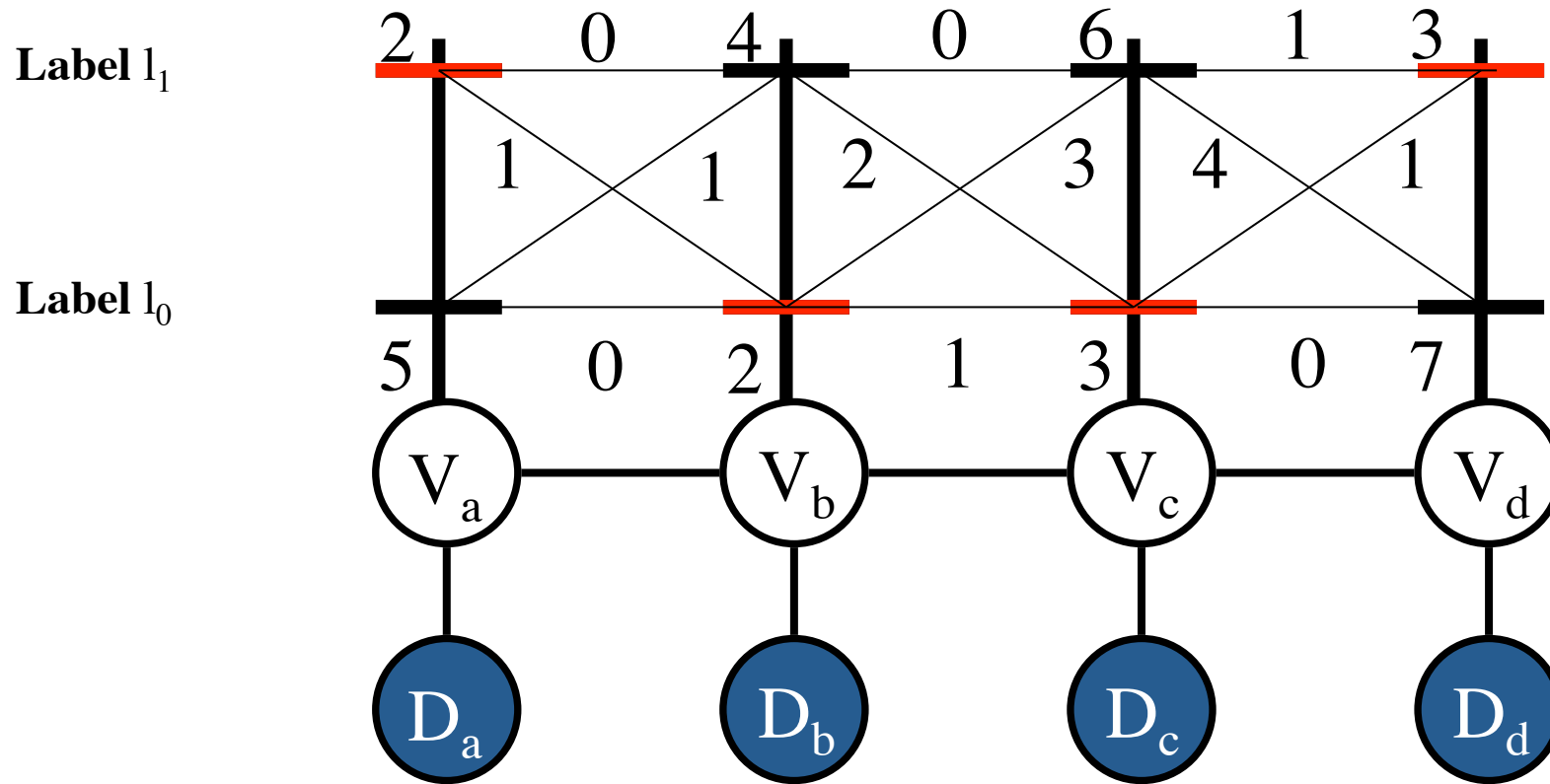
# Energy Function



Pairwise Potential

$$Q(f) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

# Energy Function



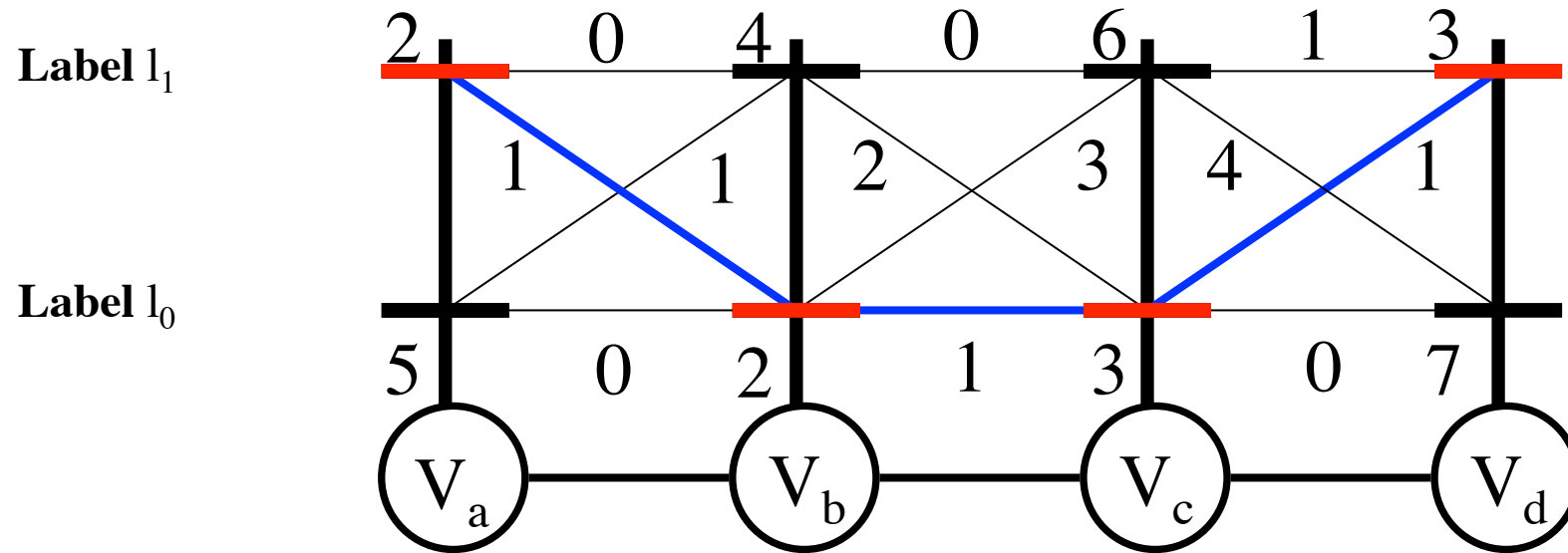
$$Q(f; \theta) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

Parameter

# Overview

- Basics: problem formulation
  - Energy Function
  - MAP Estimation
  - Computing min-marginals
  - Reparameterization
- Solutions
  - Belief Propagation and related methods [Lecture 1]
  - Graph cuts [Lecture 2]

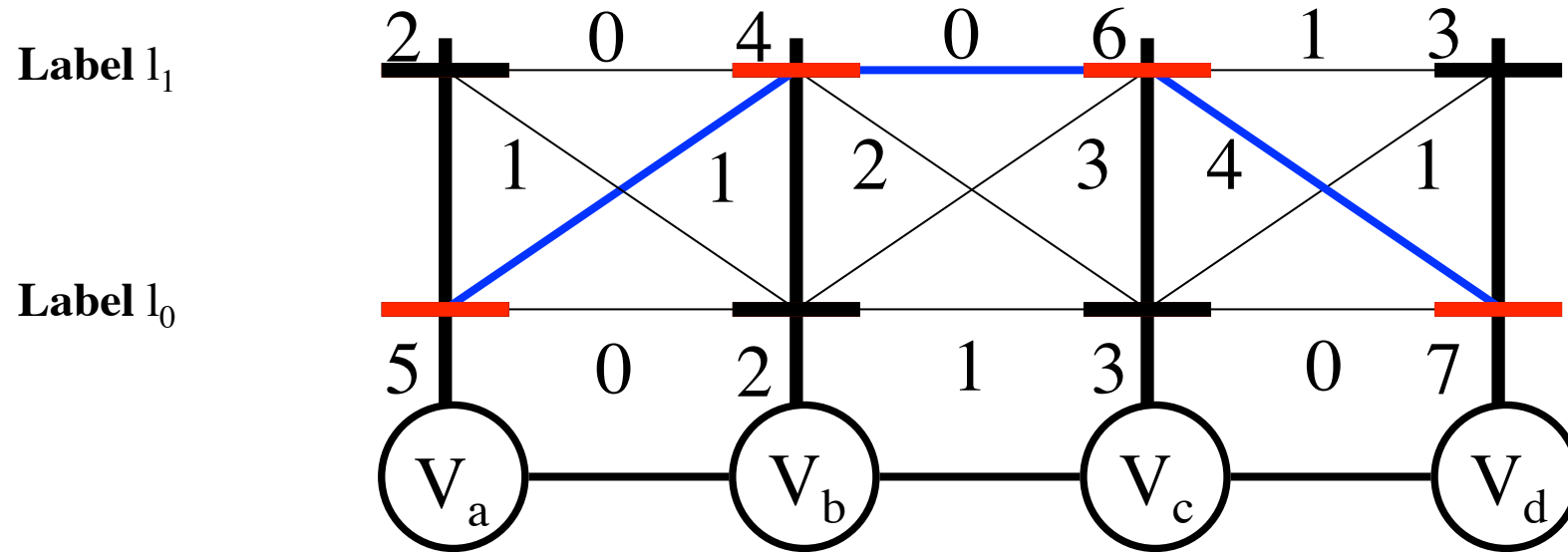
# MAP Estimation



$$Q(f; \theta) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

$$2 + 1 + 2 + 1 + 3 + 1 + 3 = 13$$

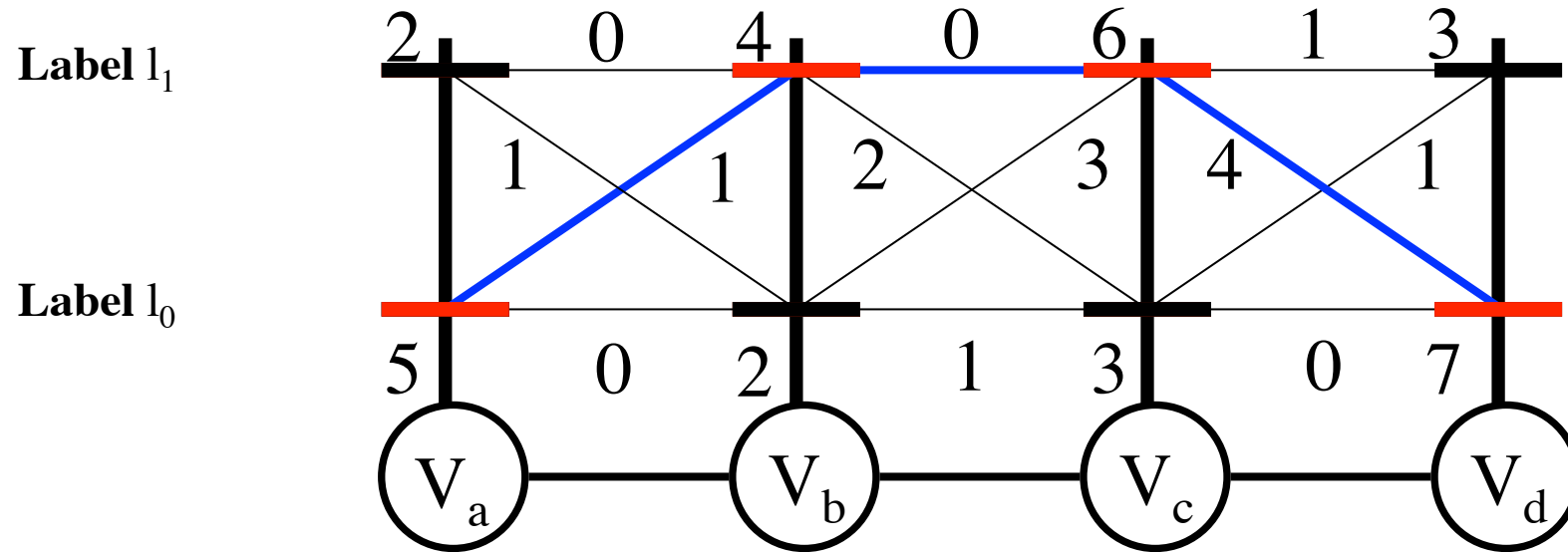
# MAP Estimation



$$Q(f; \theta) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

$$5 + 1 + 4 + 0 + 6 + 4 + 7 = 27$$

# MAP Estimation



$$q^* = \min Q(\mathbf{f}; \theta) = Q(\mathbf{f}^*; \theta)$$

$$Q(\mathbf{f}; \theta) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

$$\mathbf{f}^* = \arg \min Q(\mathbf{f}; \theta)$$

Equivalent to maximizing the associated probability



# MAP Estimation

16 possible labellings

$$f^* = \{1, 0, 0, 1\}$$

$$q^* = 13$$

f(a)	f(b)	f(c)	f(d)	Q(f; $\theta$ )
0	0	0	0	18
0	0	0	1	15
0	0	1	0	27
0	0	1	1	20
0	1	0	0	22
0	1	0	1	19
0	1	1	0	27
0	1	1	1	20

f(a)	f(b)	f(c)	f(d)	Q(f; $\theta$ )
1	0	0	0	16
1	0	0	1	13
1	0	1	0	25
1	0	1	1	18
1	1	0	0	18
1	1	0	1	15
1	1	1	0	23
1	1	1	1	16

# Computational Complexity

Segmentation

$2^{|V|}$



$|V|$  = number of pixels  $\approx 153600$

Can we do better than brute-force?

MAP Estimation is NP-hard !!

# MAP Inference / Energy Minimization

- Computing the assignment minimizing the energy in NP-hard in general

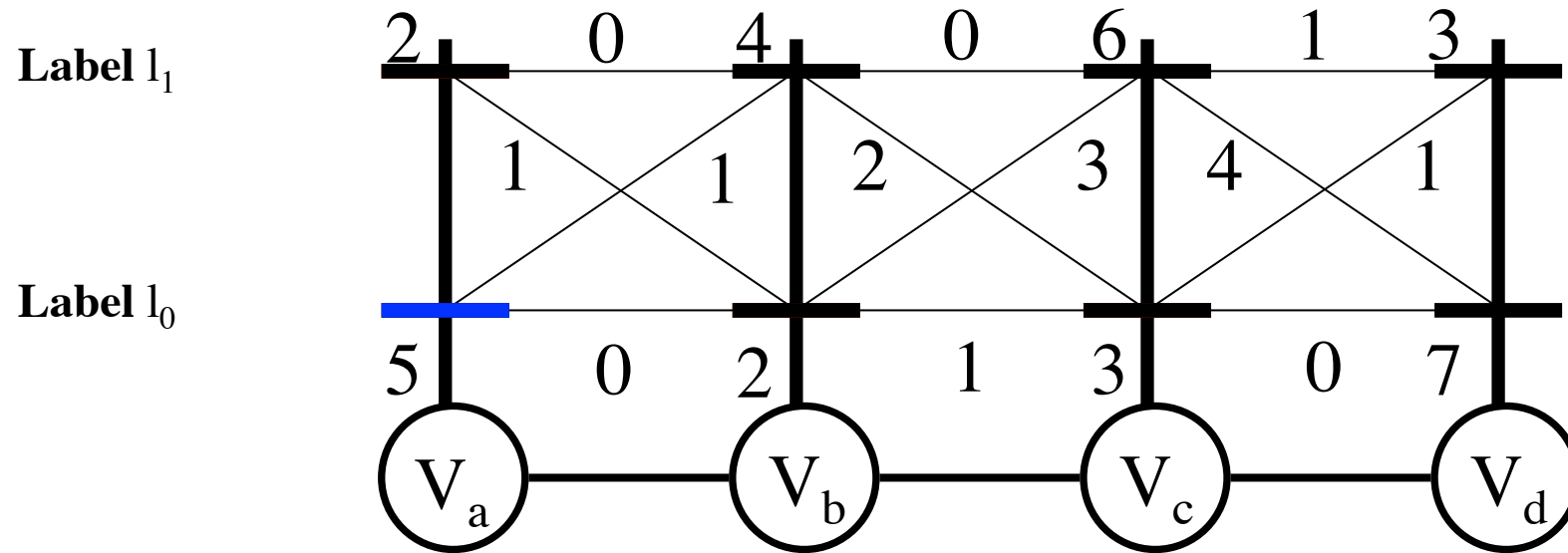
$$\operatorname{argmin}_{\mathbf{y} \in \mathcal{Y}} E(\mathbf{y}; \mathbf{x}) = \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} P(\mathbf{y} | \mathbf{x})$$

- Exact inference is possible in some cases, e.g.,
  - Low treewidth graphs  $\rightarrow$  message-passing
  - Submodular potentials  $\rightarrow$  graph cuts
- Efficient approximate inference algorithms exist
  - Message passing on general graphs
  - Move-making algorithms
  - Relaxation algorithms

# Overview

- Basics: problem formulation
  - Energy Function
  - MAP Estimation
  - Computing min-marginals
  - Reparameterization
- Solutions
  - Belief Propagation and related methods [Lecture 1]
  - Graph cuts [Lecture 2]

# Min-Marginals



Not a marginal (no summation)

$$f^* = \arg \min Q(f; \theta) \quad \text{such that } f(a) = i$$

Min-marginal  $q_{a,i}$

# Min-Marginals

16 possible labellings

$$Q_{a;0} = 15$$

f(a)	f(b)	f(c)	f(d)	Q(f; $\theta$ )
0	0	0	0	18
0	0	0	1	15
0	0	1	0	27
0	0	1	1	20
0	1	0	0	22
0	1	0	1	19
0	1	1	0	27
0	1	1	1	20

f(a)	f(b)	f(c)	f(d)	Q(f; $\theta$ )
1	0	0	0	16
1	0	0	1	13
1	0	1	0	25
1	0	1	1	18
1	1	0	0	18
1	1	0	1	15
1	1	1	0	23
1	1	1	1	16

# Min-Marginals

16 possible labellings

$$Q_{a;1} = 13$$

f(a)	f(b)	f(c)	f(d)	Q(f; $\theta$ )
0	0	0	0	18
0	0	0	1	15
0	0	1	0	27
0	0	1	1	20
0	1	0	0	22
0	1	0	1	19
0	1	1	0	27
0	1	1	1	20

f(a)	f(b)	f(c)	f(d)	Q(f; $\theta$ )
1	0	0	0	16
1	0	0	1	13
1	0	1	0	25
1	0	1	1	18
1	1	0	0	18
1	1	0	1	15
1	1	1	0	23
1	1	1	1	16

# Min-Marginals and MAP

- Minimum min-marginal of any variable = energy of MAP labelling

$$\min_i q_{a;i}$$

$$\min_i ( \min_f Q(f; \theta) \text{ such that } f(a) = i )$$

$V_a$  has to take one label

$$\min_f Q(f; \theta)$$



# Summary

## Energy Function

$$Q(\mathbf{f}; \theta) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

## MAP Estimation

$$\mathbf{f}^* = \arg \min Q(\mathbf{f}; \theta)$$

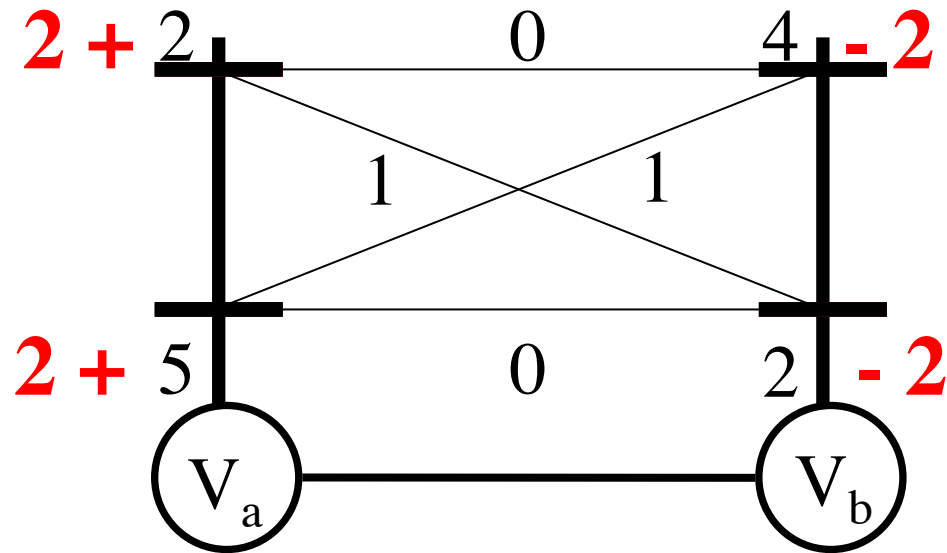
## Min-marginals

$$q_{a;i} = \min Q(\mathbf{f}; \theta) \quad \text{s.t. } f(a) = i$$

# Overview

- Basics: problem formulation
  - Energy Function
  - MAP Estimation
  - Computing min-marginals
  - Reparameterization
- Solutions
  - Belief Propagation and related methods [Lecture 1]
  - Graph cuts [Lecture 2]

# Reparameterization



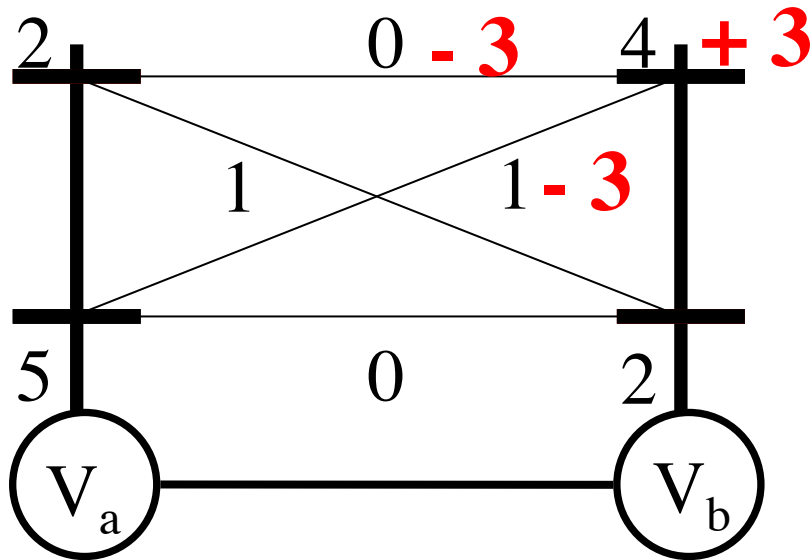
f(a)	f(b)	Q(f; $\theta$ )
0	0	7 + 2 - 2
0	1	10 + 2 - 2
1	0	5 + 2 - 2
1	1	6 + 2 - 2

Add a constant to all  $\theta_{a;i}$

Subtract that constant from all  $\theta_{b;k}$

$$Q(f; \theta') = Q(f; \theta)$$

# Reparameterization



f(a)	f(b)	Q(f; $\theta$ )
0	0	7
0	1	10 - 3 + 3
1	0	5
1	1	6 - 3 + 3

Add a constant to one  $\theta_{b;k}$

Subtract that constant from  $\theta_{ab;ik}$  for all 'i'

$$Q(f; \theta') = Q(f; \theta)$$

# Reparameterization

$\theta'$  is a reparameterization of  $\theta$ , iff

$$Q(f; \theta') = Q(f; \theta), \text{ for all } f \quad \theta' \equiv \theta$$

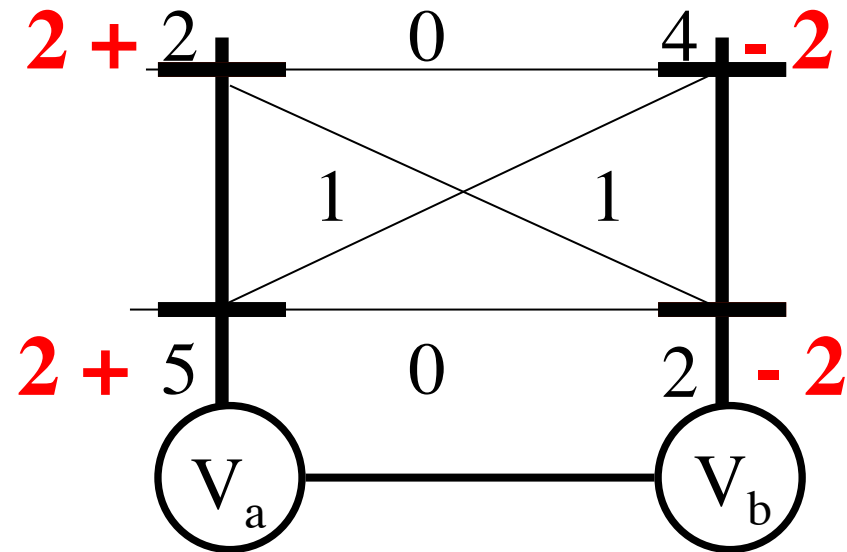
Equivalently

$$\theta'_{a;i} = \theta_{a;i} + M_{ba;i}$$

$$\theta'_{b;k} = \theta_{b;k} + M_{ab;k}$$

$$\theta'_{ab;ik} = \theta_{ab;ik} - M_{ab;k} - M_{ba;i}$$

Kolmogorov, PAMI, 2006



# Recap

## MAP Estimation

$$\mathbf{f}^* = \arg \min \mathbf{Q}(\mathbf{f}; \boldsymbol{\theta})$$

$$\mathbf{Q}(\mathbf{f}; \boldsymbol{\theta}) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

## Min-marginals

$$q_{a;i} = \min \mathbf{Q}(\mathbf{f}; \boldsymbol{\theta}) \quad \text{s.t. } f(a) = i$$

## Reparameterization

$$\mathbf{Q}(\mathbf{f}; \boldsymbol{\theta}') = \mathbf{Q}(\mathbf{f}; \boldsymbol{\theta}), \text{ for all } \mathbf{f} \quad \boldsymbol{\theta}' \equiv \boldsymbol{\theta}$$

# Overview

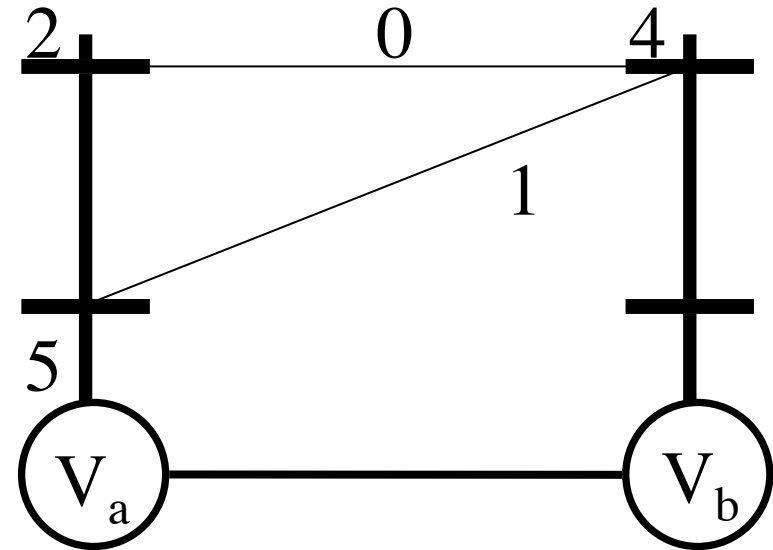
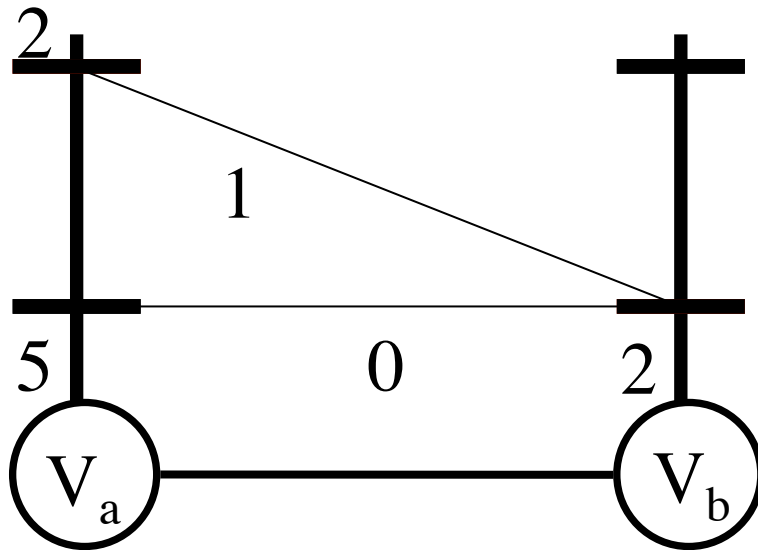
- Basics: problem formulation
  - Energy Function
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  - Reparameterization
- Solutions
  - Belief Propagation and related methods [Lecture 1]
  - Graph cuts [Lecture 2]

# Belief Propagation

- Remember, some MAP problems are easy
- Belief Propagation gives exact MAP for chains
- Exact MAP for trees
- Clever Reparameterization



# Two Variables

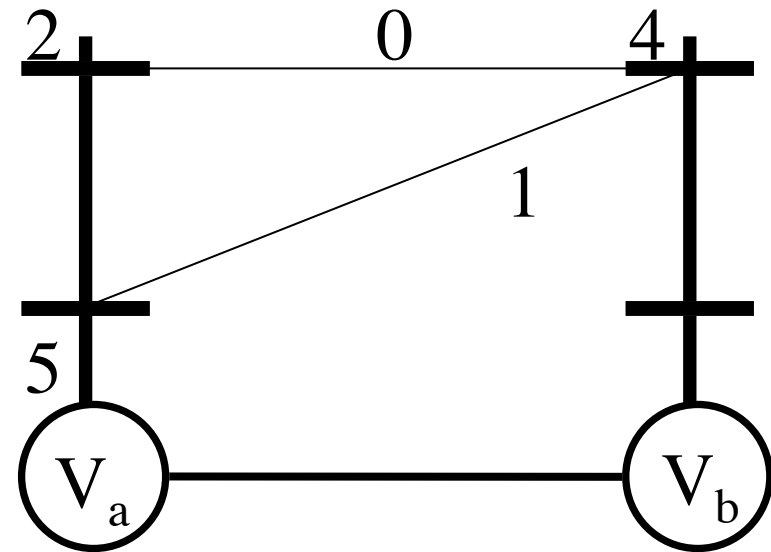
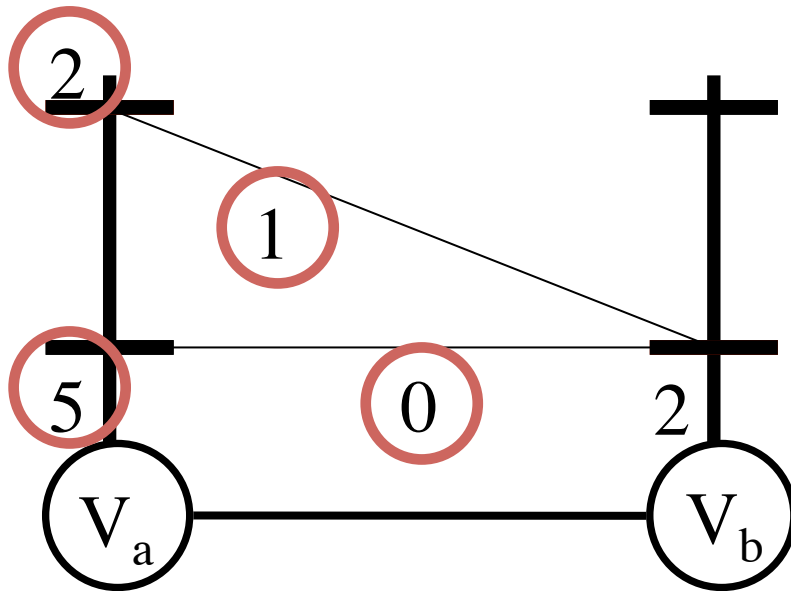


Add a constant to one  $\theta_{b;k}$

Subtract that constant from  $\theta_{ab;ik}$  for all 'i'

Choose the *right* constant  $\theta'_{b;k} = q_{b;k}$

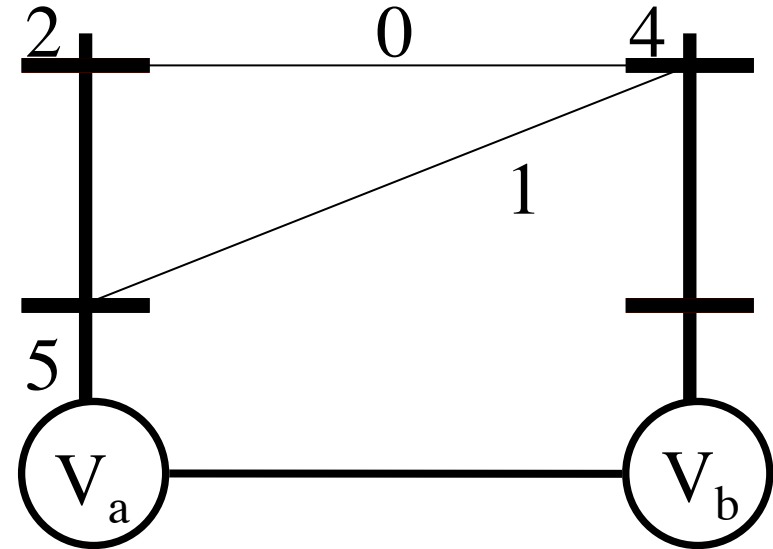
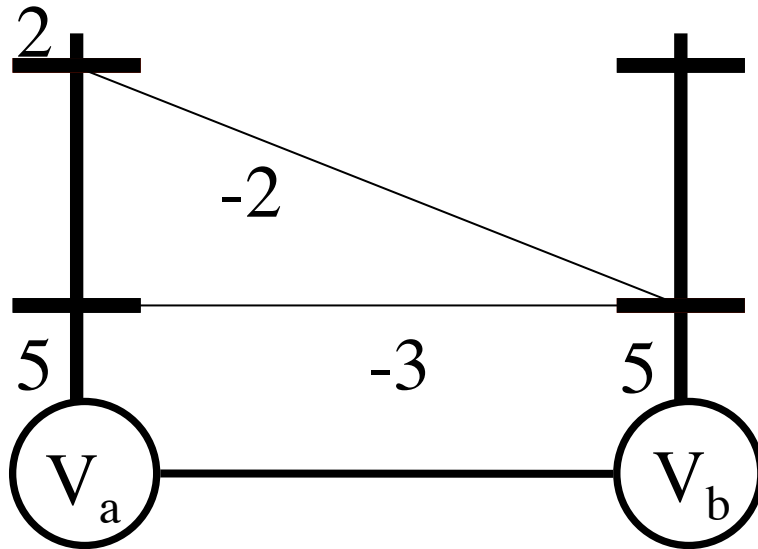
# Two Variables



$$M_{ab;0} = \min \begin{cases} \theta_{a;0} + \theta_{ab;00} = 5 + 0 \\ \theta_{a;1} + \theta_{ab;10} = 2 + 1 \end{cases}$$

Choose the *right* constant  $\theta'_{b;k} = q_{b;k}$

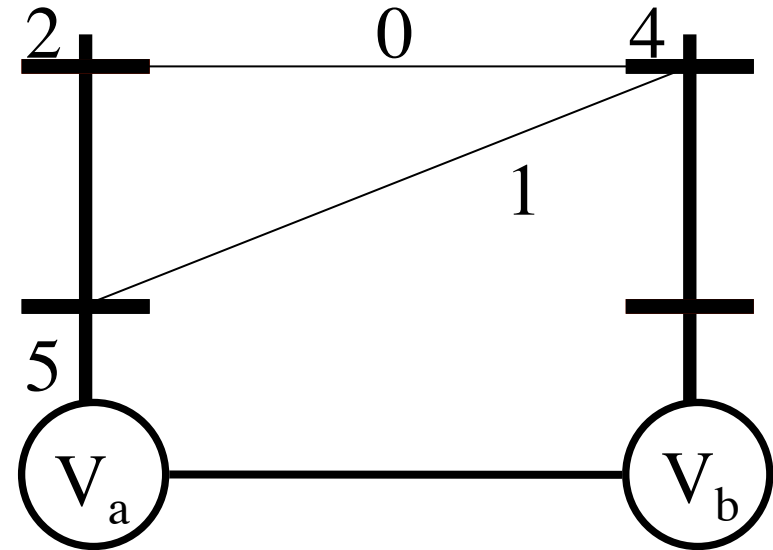
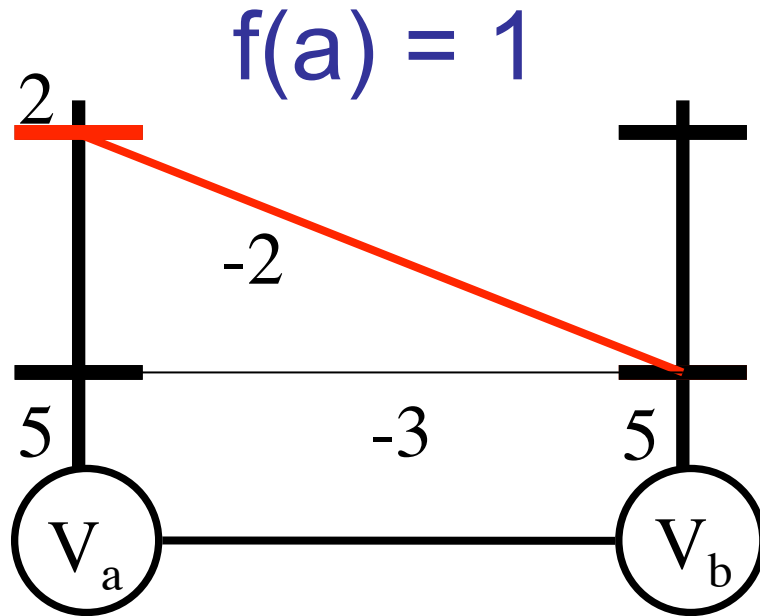
# Two Variables



Choose the *right* constant

$$\theta'_{b;k} = q_{b;k}$$

# Two Variables

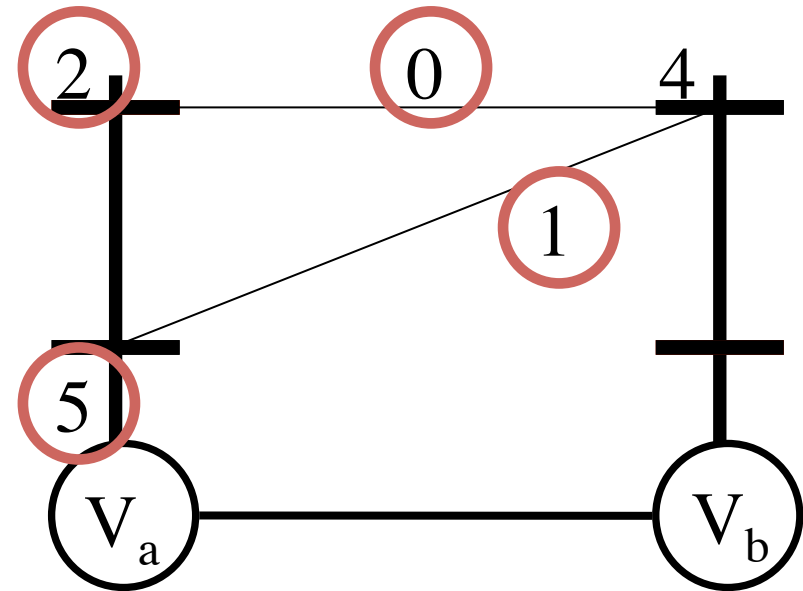
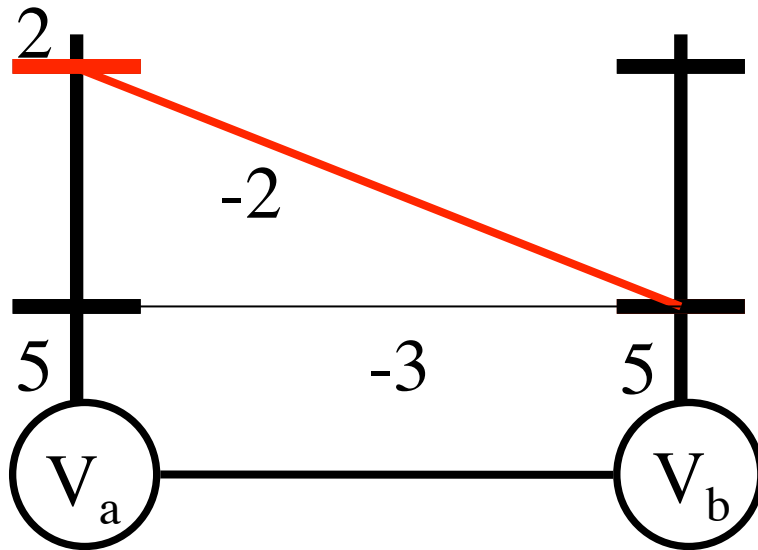


$$\theta'_{b;0} = q_{b;0}$$

Potentials along the red path add up to 0

Choose the **right** constant  $\theta'_{b;k} = q_{b;k}$

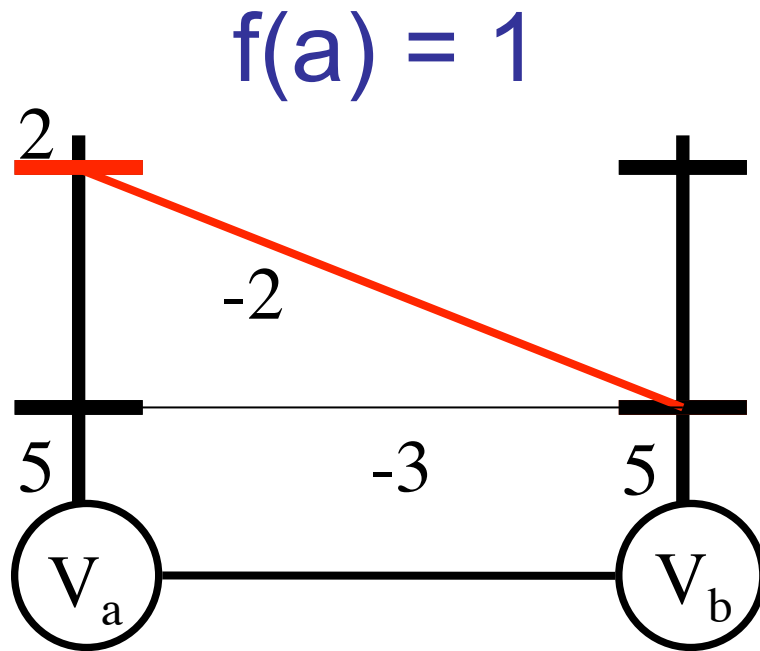
# Two Variables



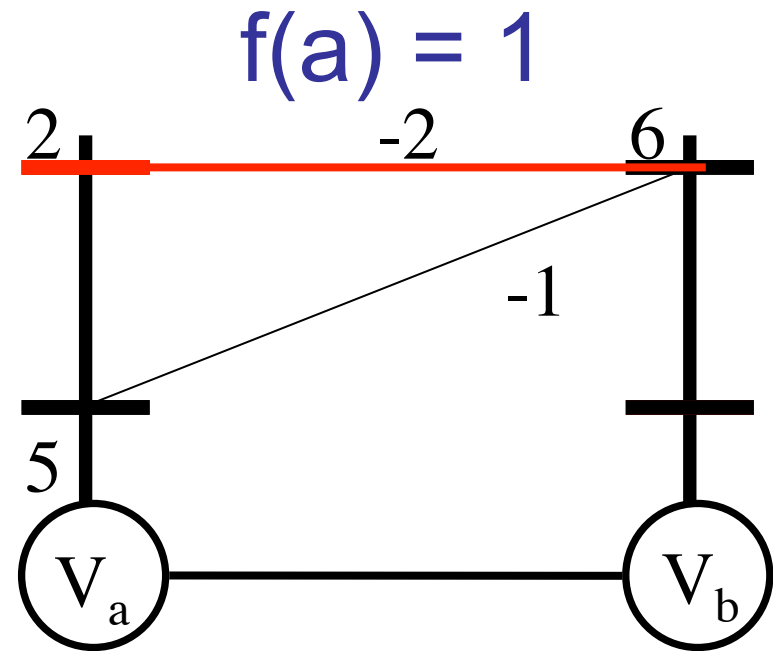
$$M_{ab;1} = \min \begin{cases} \theta_{a;0} + \theta_{ab;01} = 5 + 1 \\ \theta_{a;1} + \theta_{ab;11} = 2 + 0 \end{cases}$$

Choose the *right* constant  $\theta'_{b;k} = q_{b;k}$

# Two Variables



$$\theta'_{b;0} = q_{b;0}$$

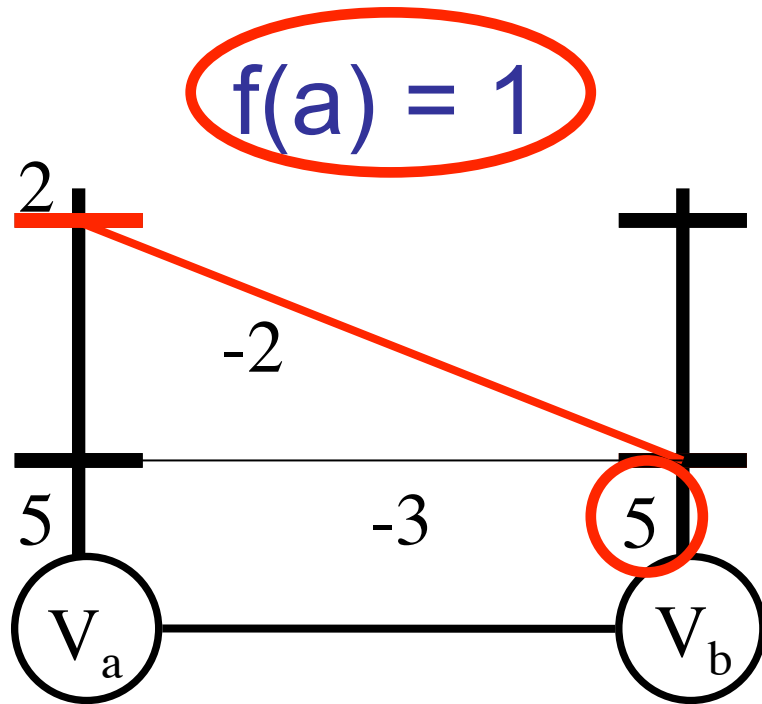


$$\theta'_{b;1} = q_{b;1}$$

Minimum of min-marginals = MAP estimate

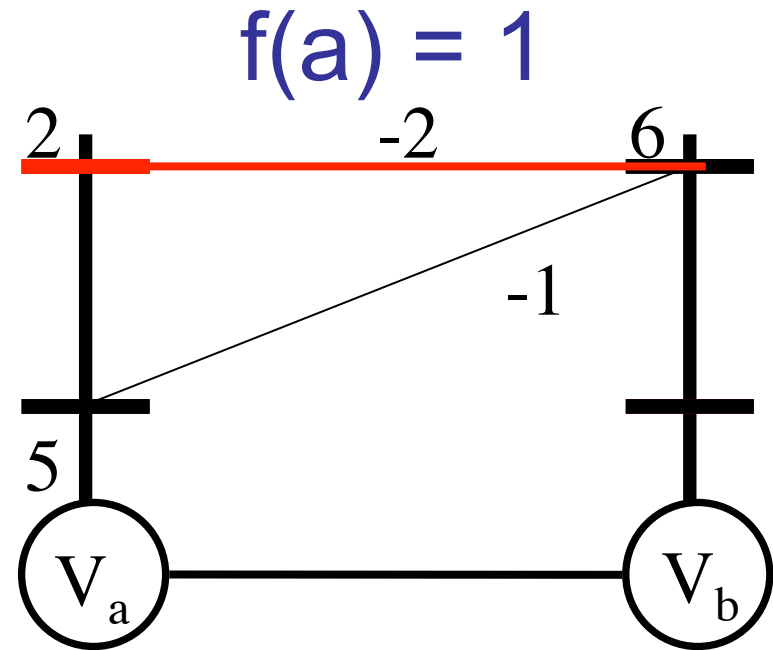
Choose the *right* constant  $\theta'_{b;k} = q_{b;k}$

# Two Variables



$$\theta'_{b;0} = q_{b;0}$$

$$f^*(b) = 0 \quad f^*(a) = 1$$

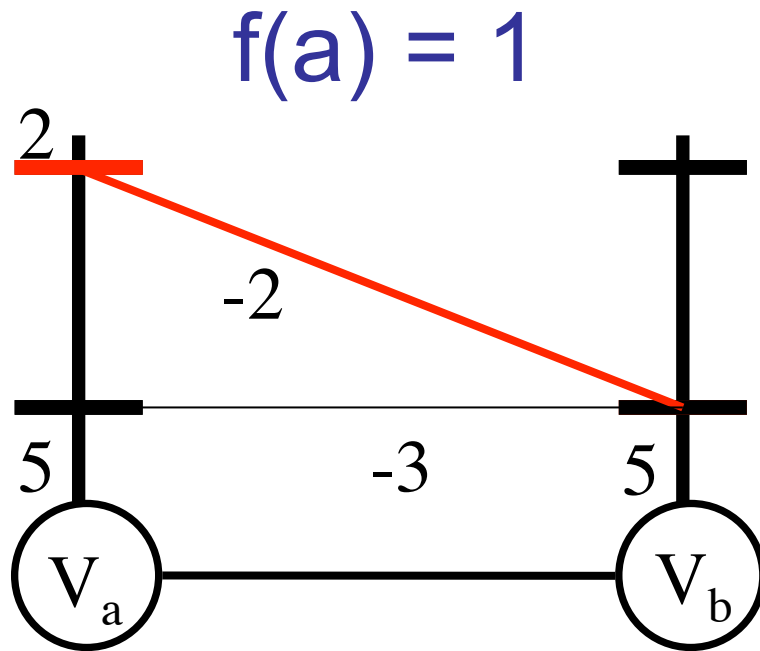


$$\theta'_{b;1} = q_{b;1}$$

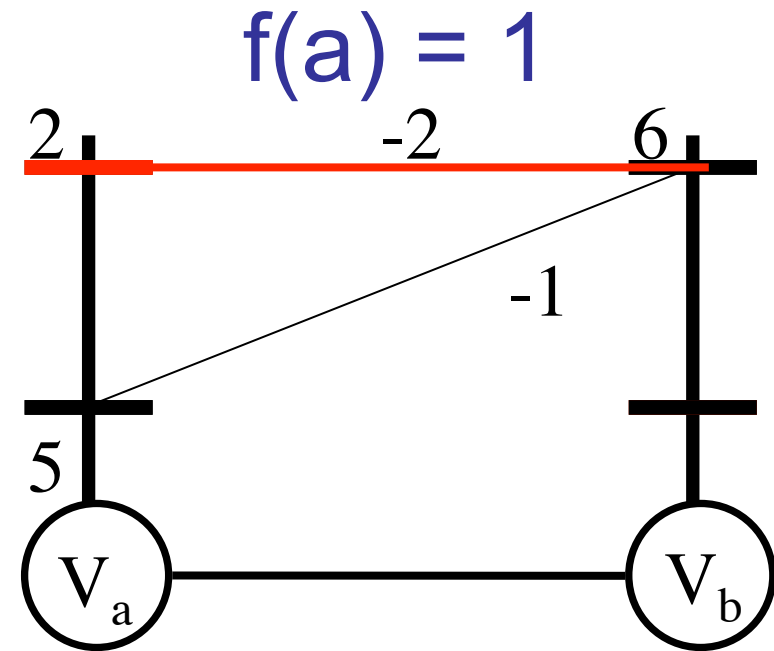
Choose the **right** constant

$$\theta'_{b;k} = q_{b;k}$$

# Two Variables



$$\theta'_{b;0} = q_{b;0}$$



$$\theta'_{b;1} = q_{b;1}$$

We get all the min-marginals of  $V_b$

Choose the *right* constant  $\theta'_{b;k} = q_{b;k}$



# Recap

We only need to know two sets of equations

General form of Reparameterization

$$\theta'_{a;i} = \theta_{a;i} + M_{ba;i} \quad \theta'_{b;k} = \theta_{b;k} + M_{ab;k}$$

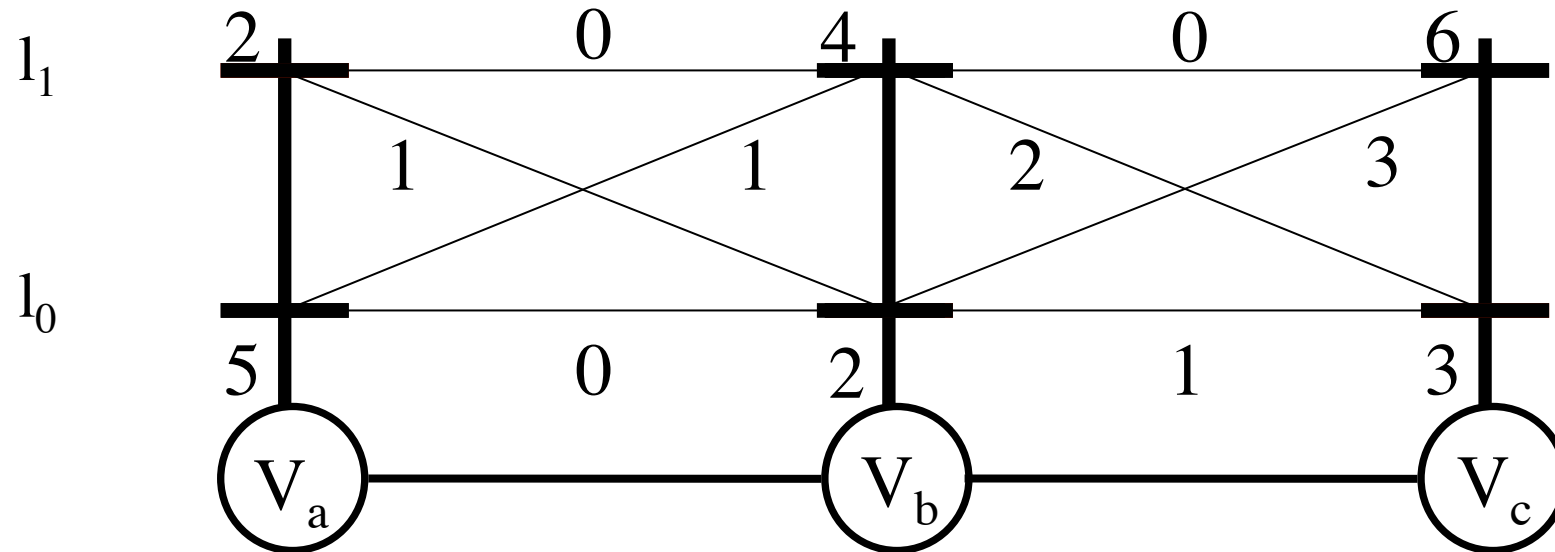
$$\theta'_{ab;ik} = \theta_{ab;ik} - M_{ab;k} - M_{ba;i}$$

Reparameterization of (a,b) in Belief Propagation

$$M_{ab;k} = \min_i \{ \theta_{a;i} + \theta_{ab;ik} \}$$

$$M_{ba;i} = 0$$

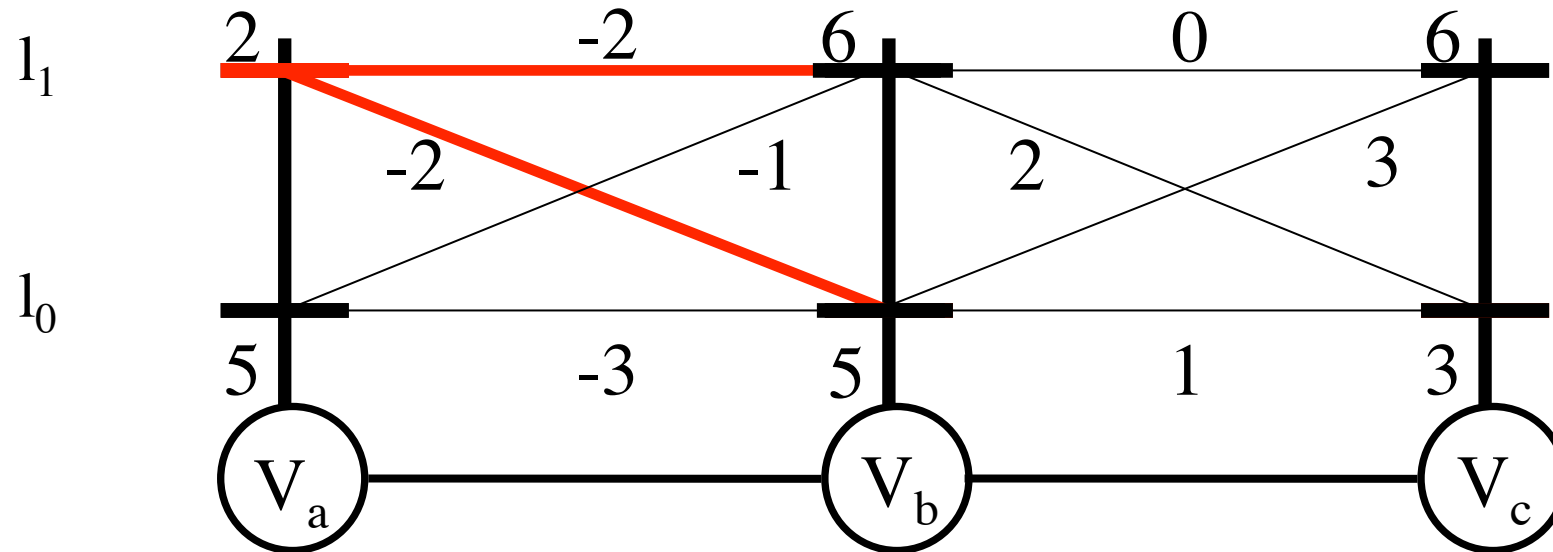
# Three Variables



Reparameterize the edge (a,b) as before

# Three Variables

$$f(a) = 1$$

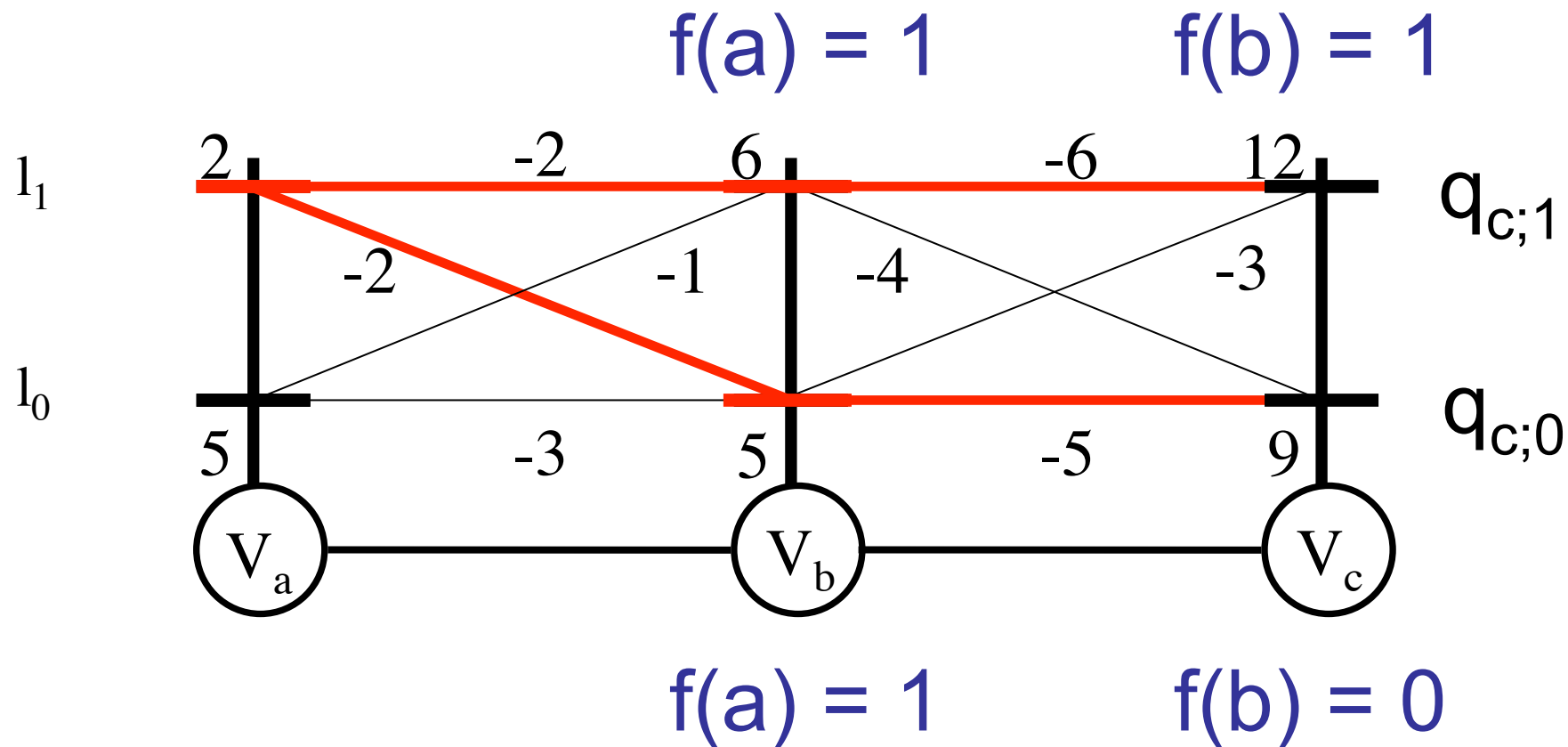


$$f(a) = 1$$

Reparameterize the edge (a,b) as before

Potentials along the red path add up to 0

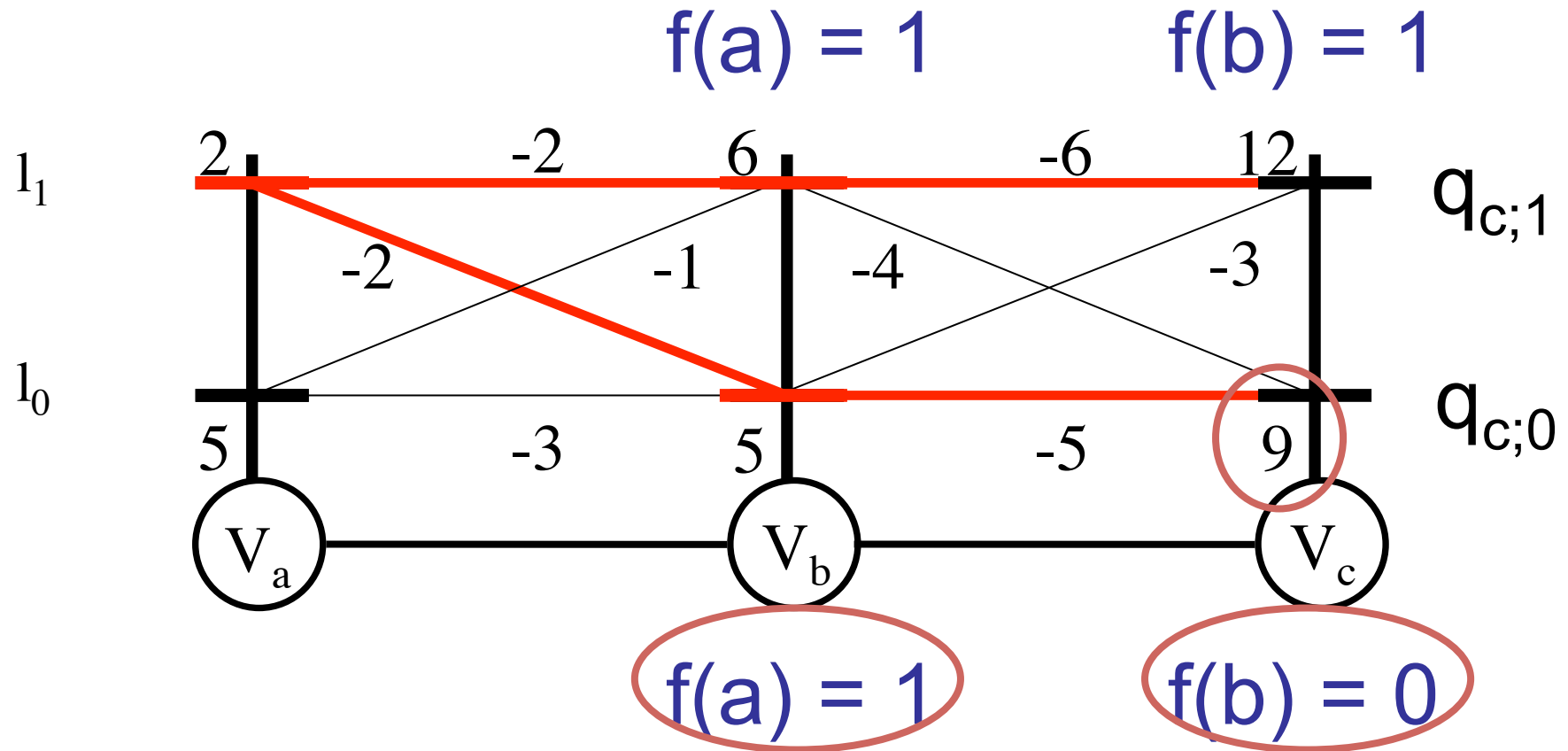
# Three Variables



Reparameterize the edge (b,c) as before

Potentials along the red path add up to 0

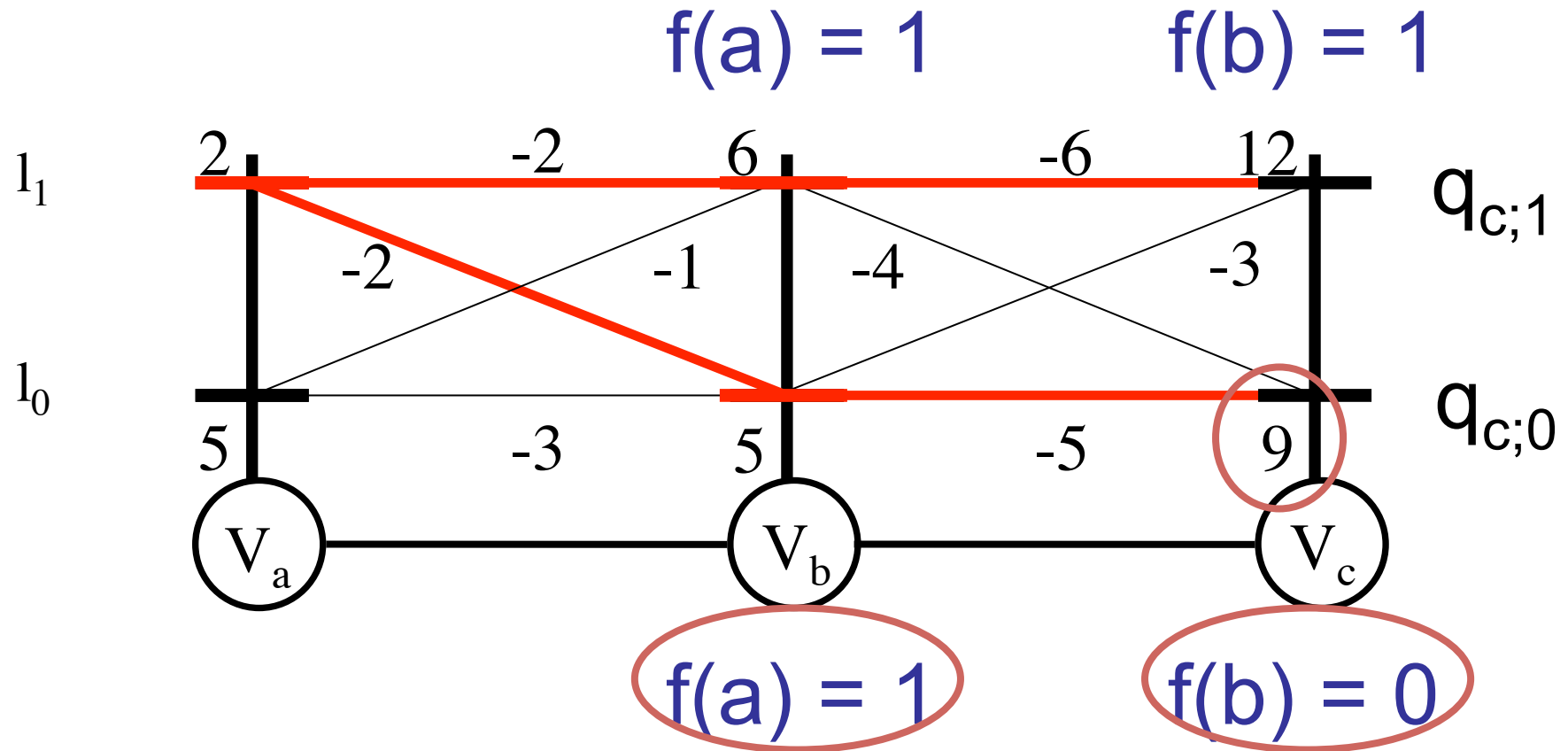
# Three Variables



$$f^*(c) = 0 \quad f^*(b) = 0 \quad f^*(a) = 1$$

**Generalizes to any length chain**

# Three Variables



$f^*(c) = 0 \quad f^*(b) = 0 \quad f^*(a) = 1$

**Only Dynamic Programming**

# Why Dynamic Programming?

3 variables  $\equiv$  2 variables + book-keeping

n variables  $\equiv$  (n-1) variables + book-keeping

Start from left, go to right

Reparameterize current edge (a,b)

$$M_{ab;k} = \min_i \{ \theta_{a;i} + \theta_{ab;ik} \}$$

$$\theta'_{b;k} = \theta_{b;k} + M_{ab;k} \quad \theta'_{ab;ik} = \theta_{ab;ik} - M_{ab;k}$$

Repeat

# Why Dynamic Programming?

Messages    Message Passing

Why stop at dynamic programming?

Start from left, go to right

Reparameterize current edge (a,b)

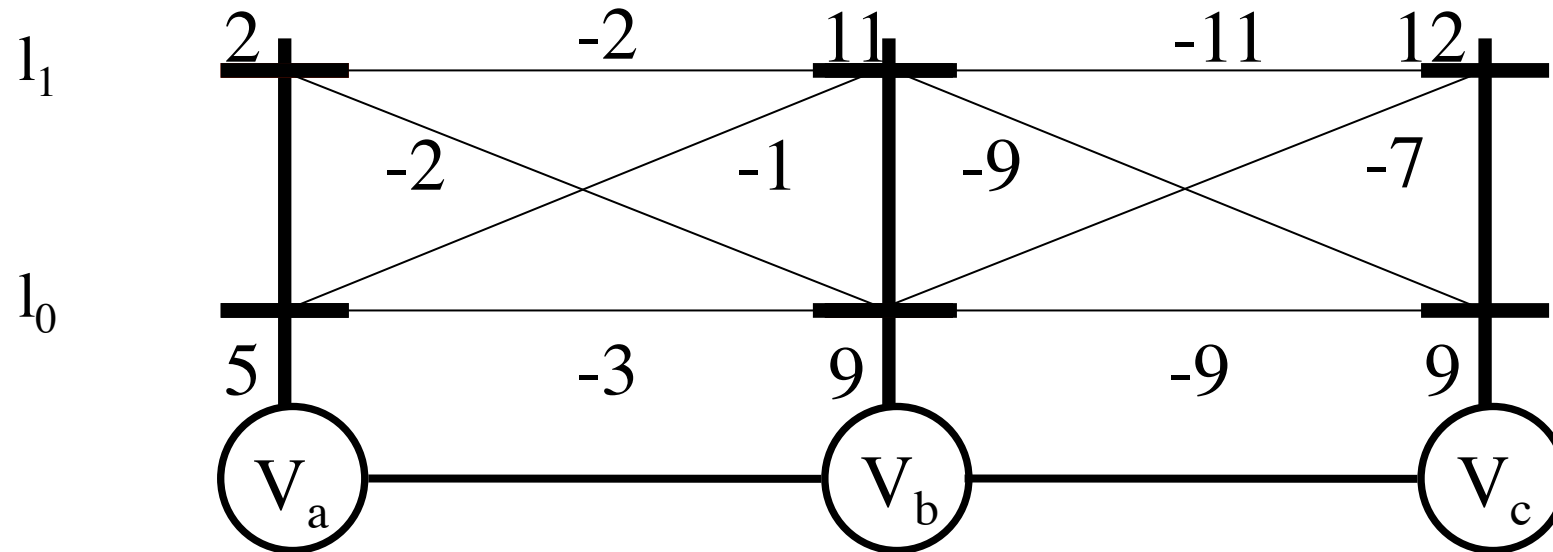
$$M_{ab;k} = \min_i \{ \theta_{a;i} + \theta_{ab;ik} \}$$

$$\theta'_{b;k} = \theta_{b;k} + M_{ab;k} \quad \theta'_{ab;ik} = \theta_{ab;ik} - M_{ab;k}$$

Repeat



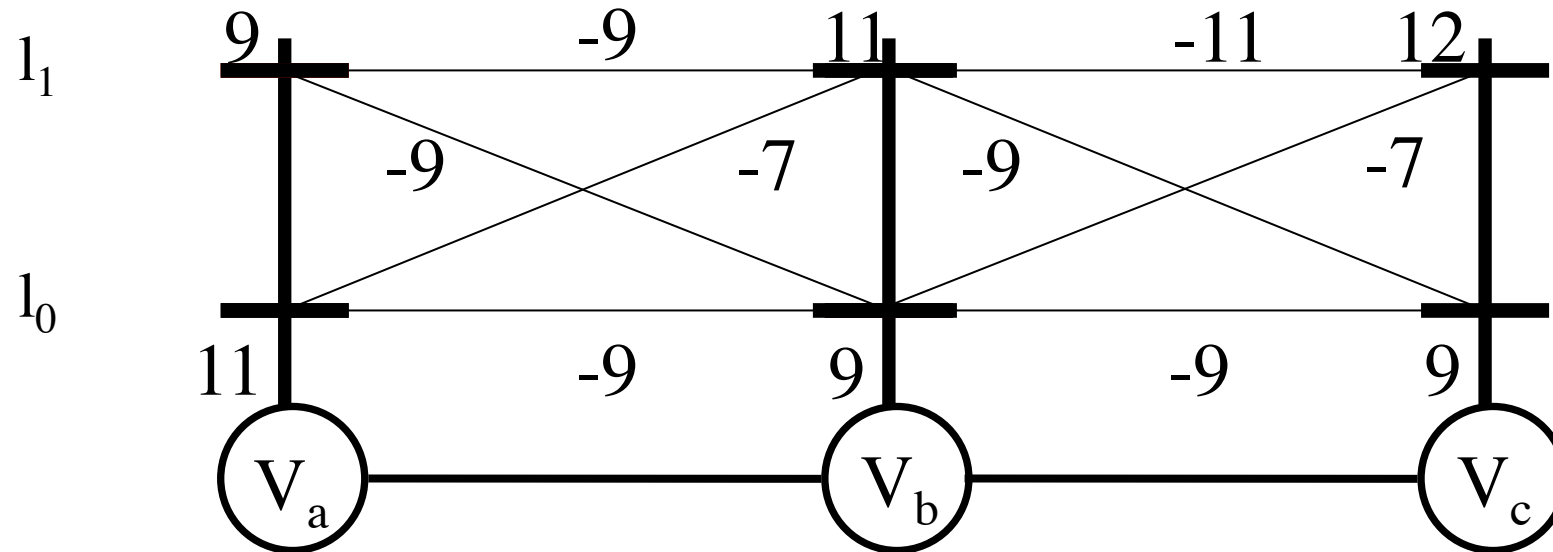
# Three Variables



Reparameterize the edge (c,b) as before

$$\theta'_{b;i} = q_{b;i}$$

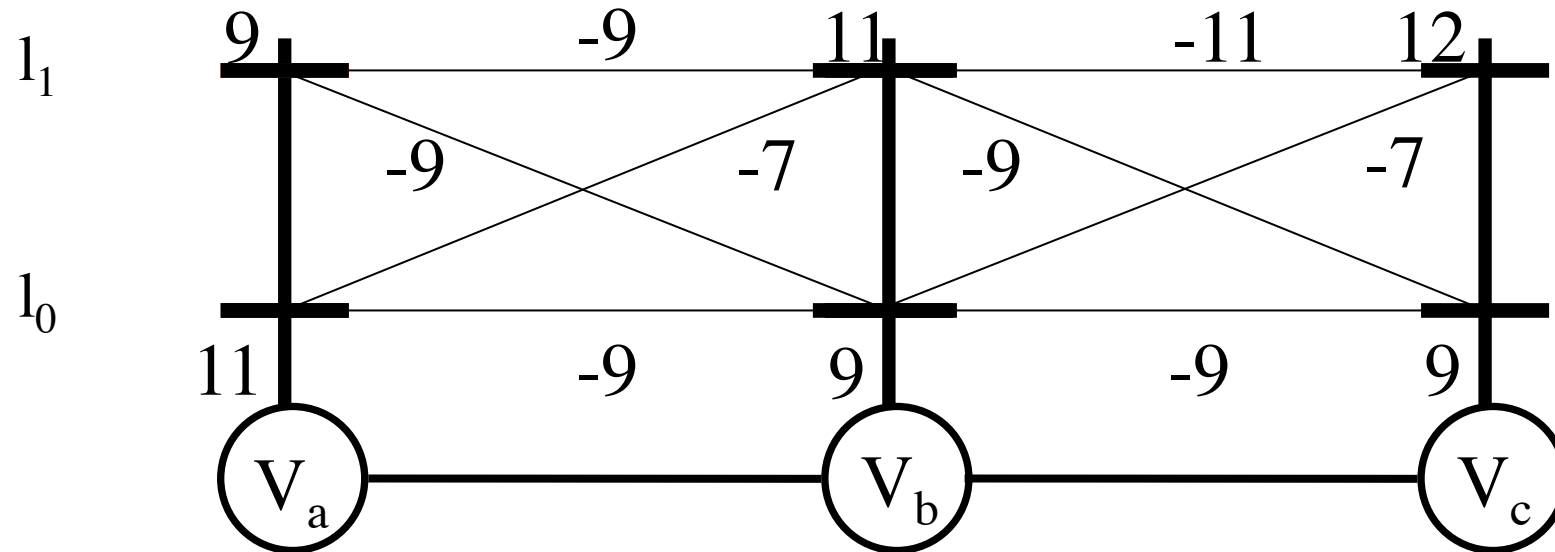
# Three Variables



Reparameterize the edge (b,a) as before

$$\theta'_{a;i} = q_{a;i}$$

# Three Variables



Forward Pass →

← Backward Pass

All min-marginals are computed

# Chains



Reparameterize the edge (1,2)

# Chains



Reparameterize the edge (1,2)

# Chains



Reparameterize the edge (2,3)

# Chains



Reparameterize the edge  $(n-1,n)$

Min-marginals  $e_n(i)$  for all labels

# Belief Propagation on Chains

Start from left, go to right

Reparameterize current edge (a,b)

$$M_{ab;k} = \min_i \{ \theta_{a;i} + \theta_{ab;ik} \}$$

$$\theta'_{b;k} = \theta_{b;k} + M_{ab;k} \quad \theta'_{ab;ik} = \theta_{ab;ik} - M_{ab;k}$$

Repeat till the end of the chain

Start from right, go to left

Repeat till the end of the chain



# Belief Propagation on Chains

- Generalizes to chains of any length
- A way of computing reparam constants
- Forward Pass - Start to End
  - MAP estimate
  - Min-marginals of final variable
- Backward Pass - End to start
  - All other min-marginals

# Computational Complexity

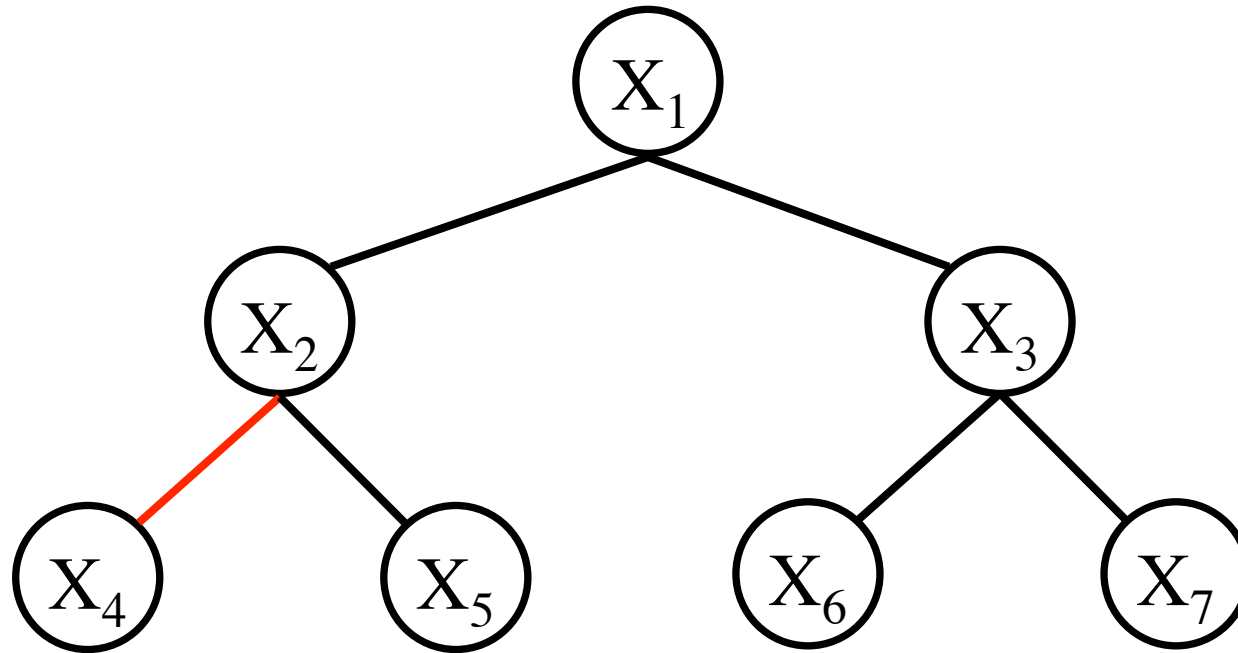
Number of reparameterization constants =  $(n-1)h$

Complexity for each constant =  $O(h)$

Total complexity =  $O(nh^2)$

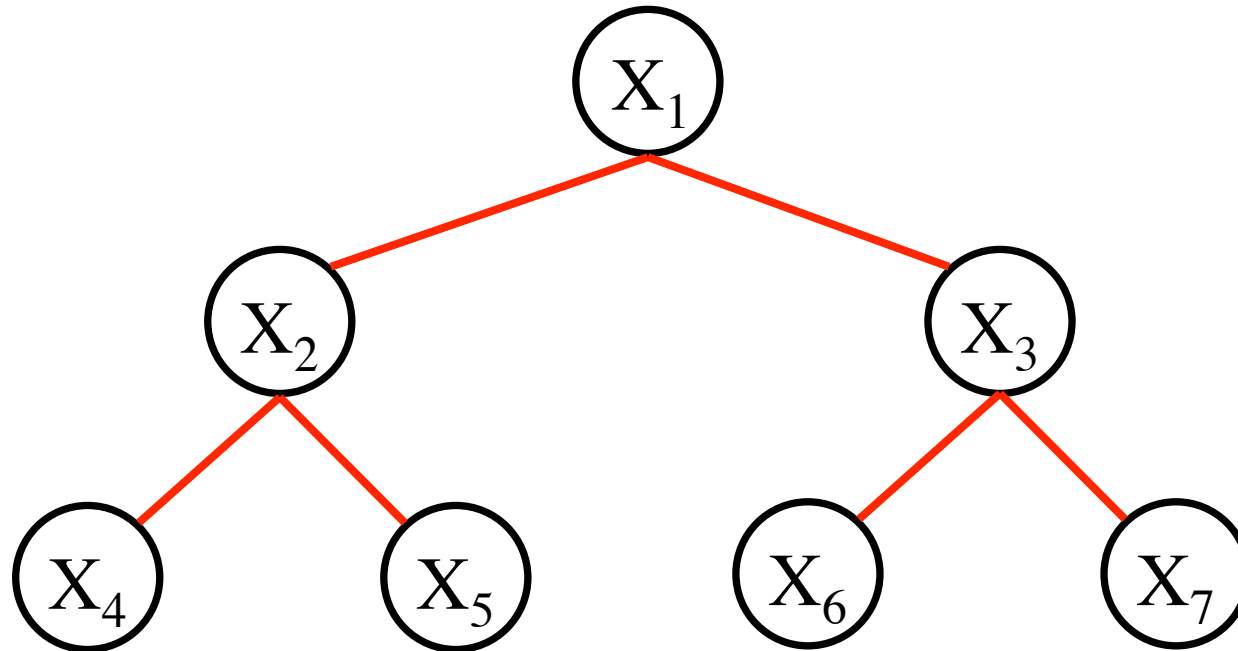
Better than brute-force  $O(h^n)$

# Trees



Reparameterize the edge (4,2)

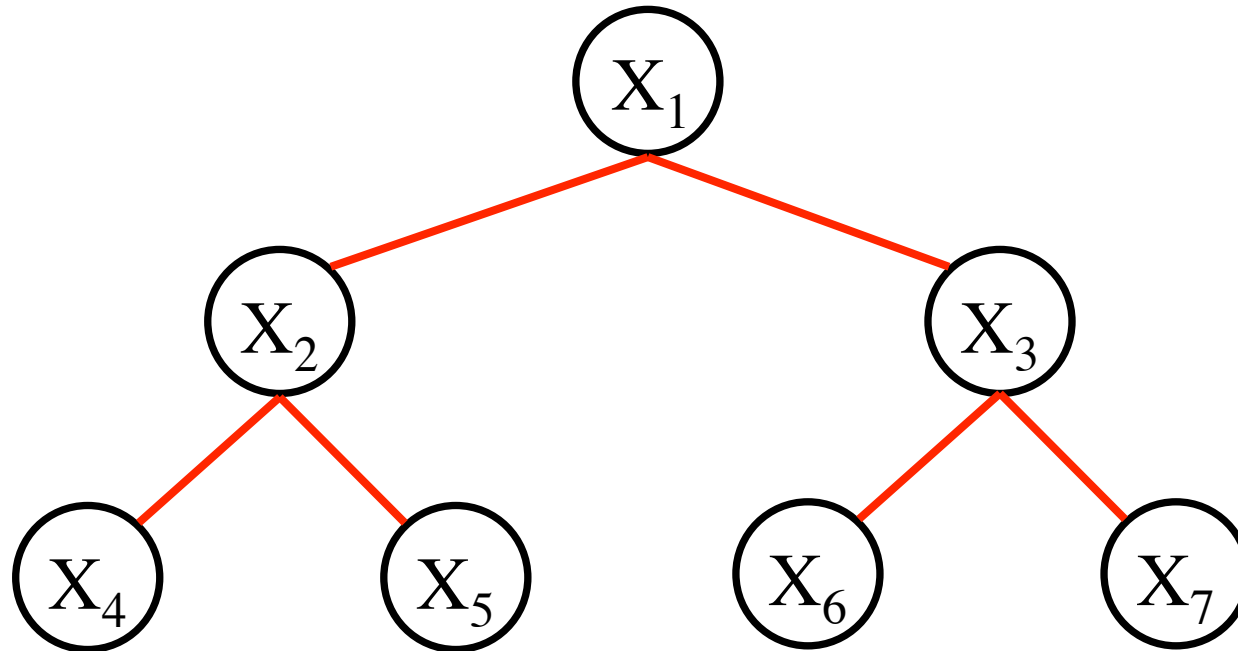
# Trees



Reparameterize the edge  $(3,1)$

Min-marginals  $e_1(i)$  for all labels

# Trees



Start from leaves and move towards root

Pick the minimum of min-marginals

Backtrack to find the best labeling  $\mathbf{x}$

# Computational Complexity

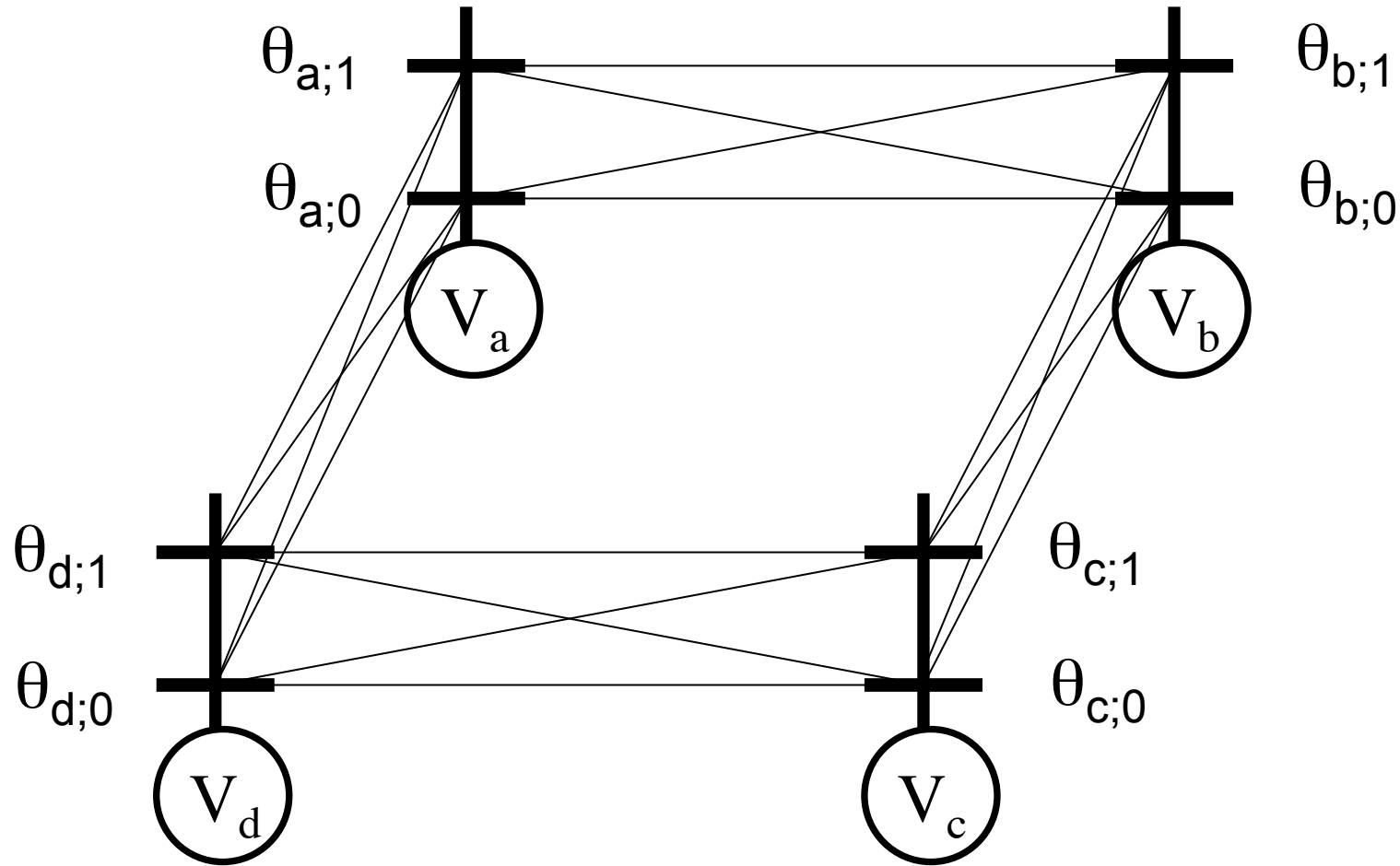
Number of reparameterization constants =  $(n-1)h$

Complexity for each constant =  $O(h)$

Total complexity =  $O(nh^2)$

Better than brute-force  $O(h^n)$

# Belief Propagation on Cycles

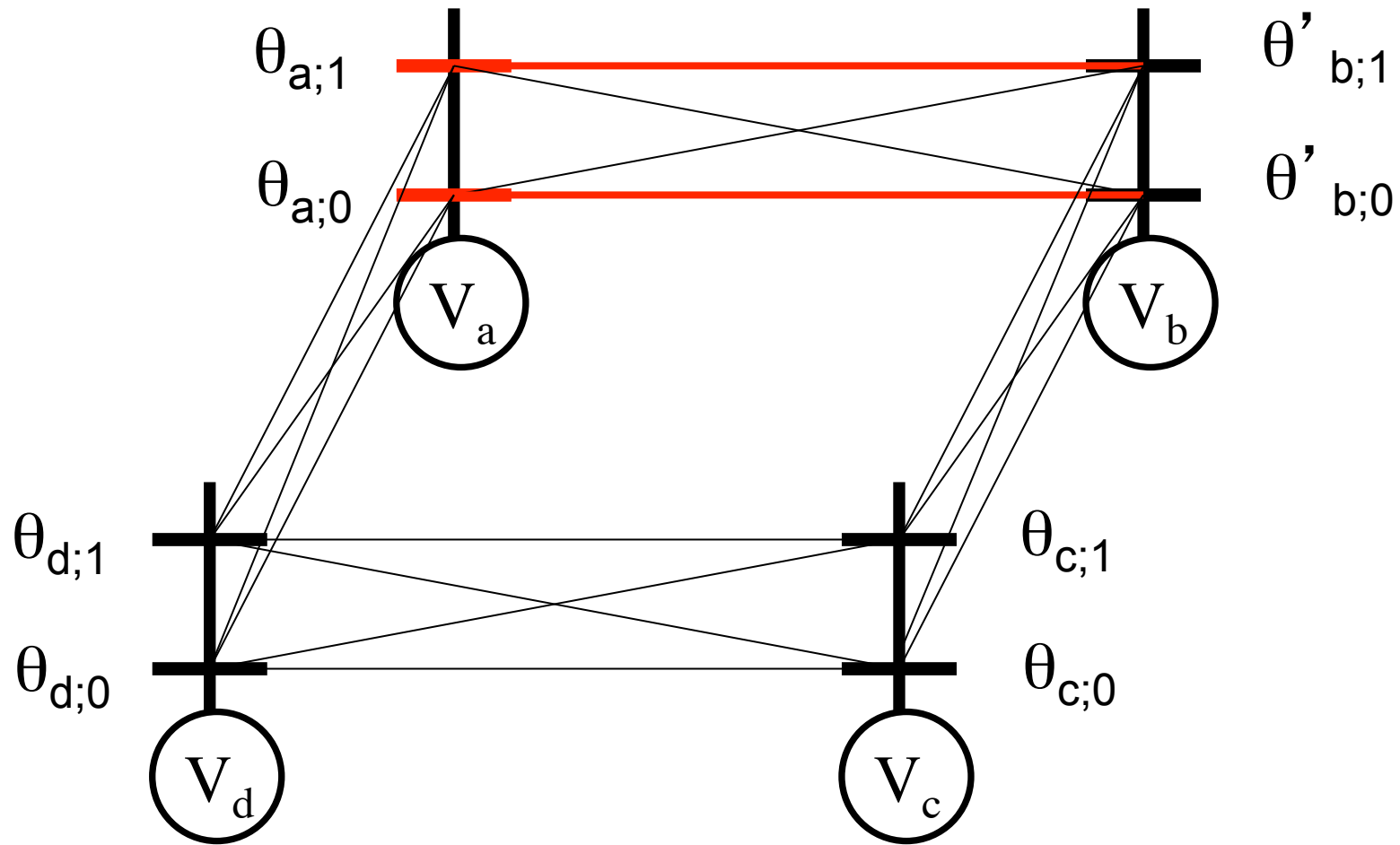


Where do we start?

Arbitrarily

Reparameterize (a,b)

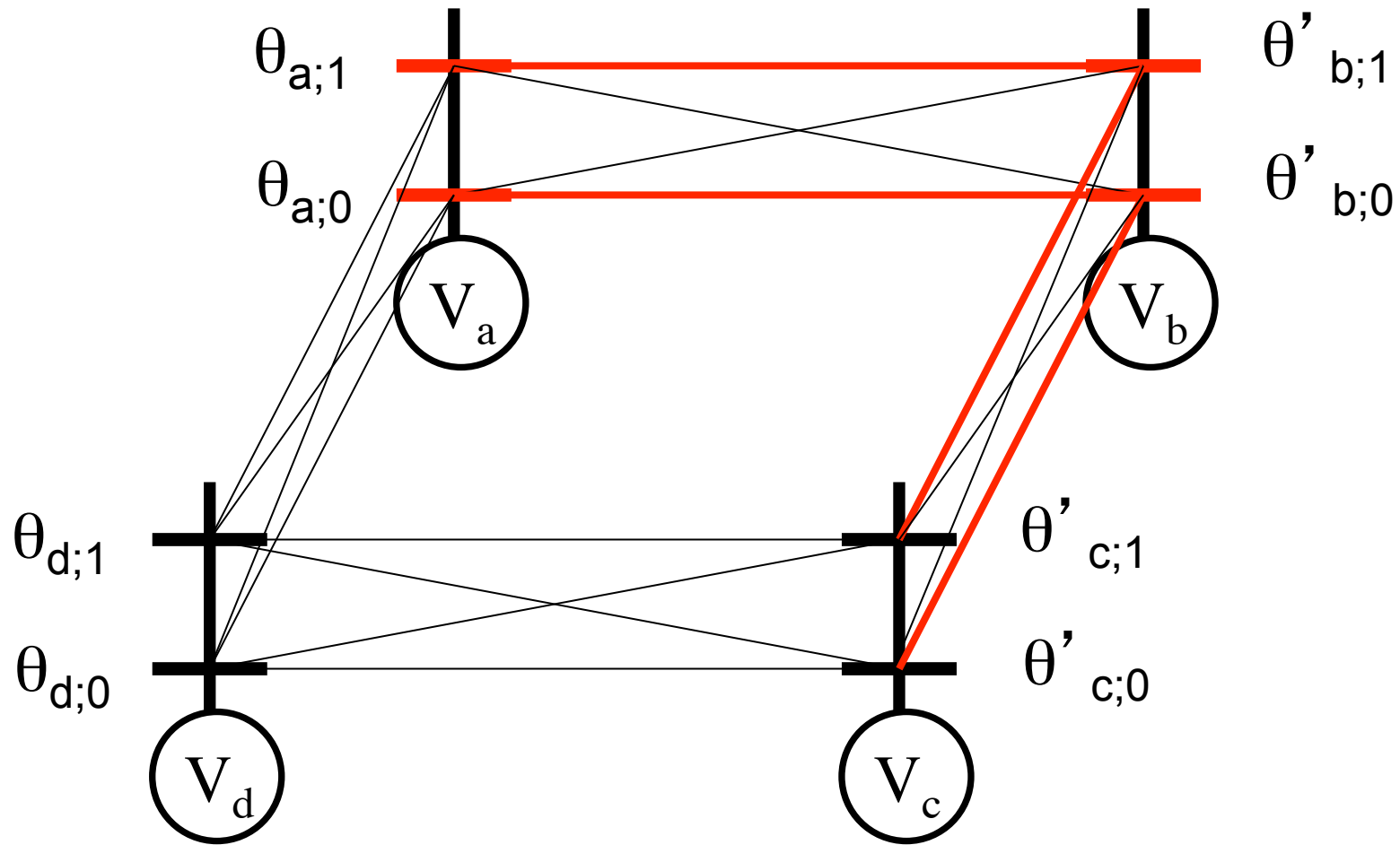
# Belief Propagation on Cycles



Potentials along the red path add up to 0

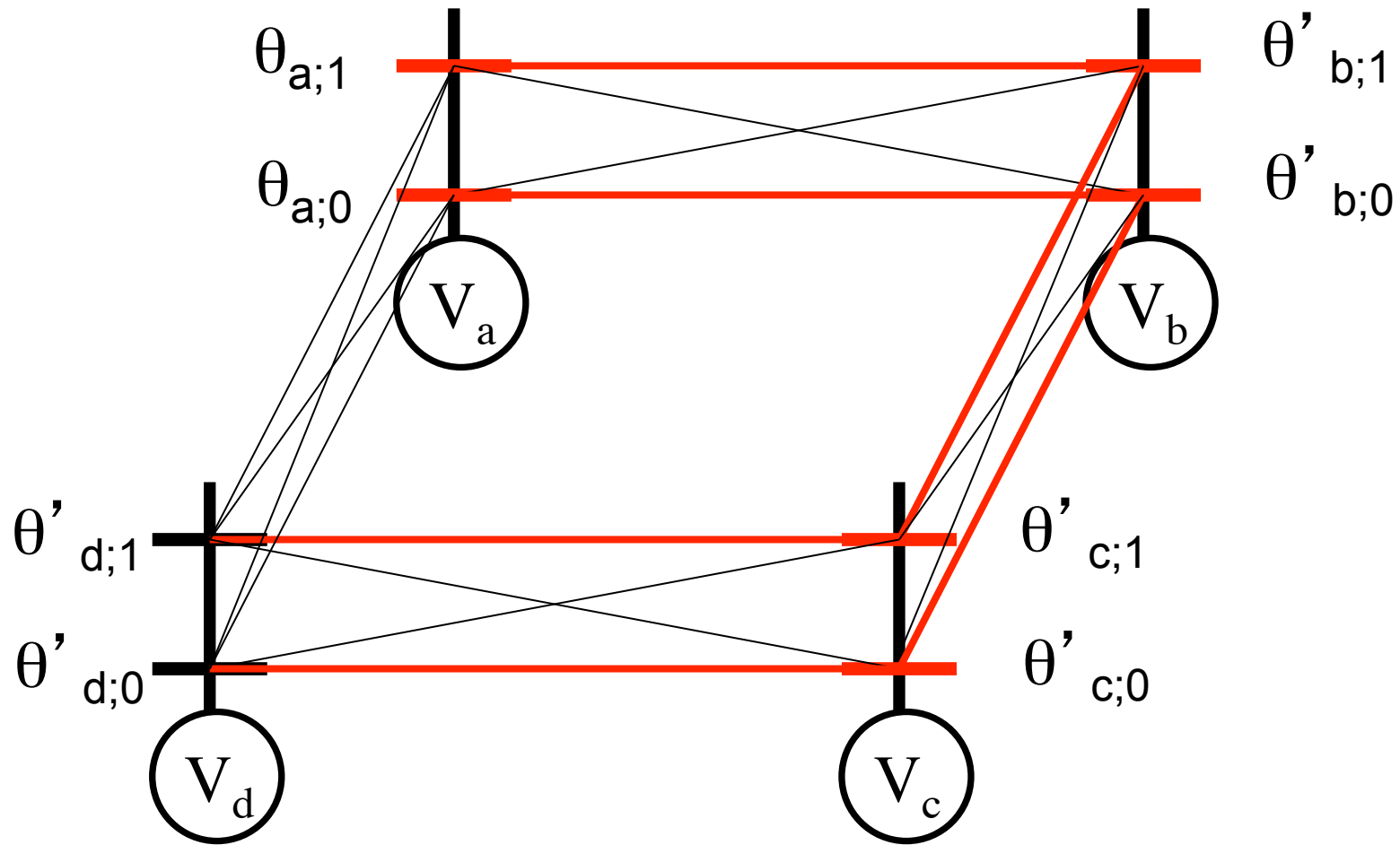


# Belief Propagation on Cycles



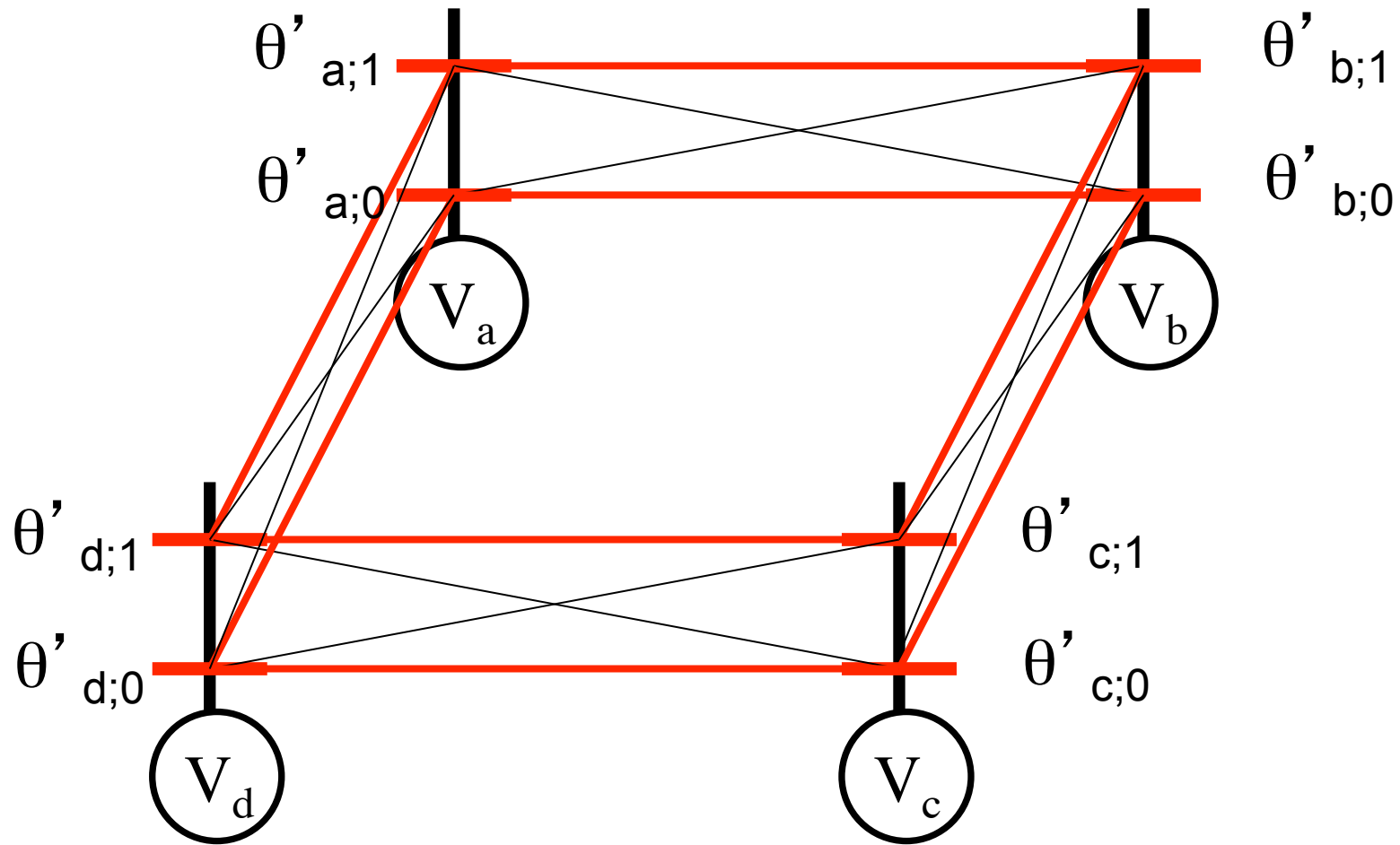
Potentials along the red path add up to 0

# Belief Propagation on Cycles



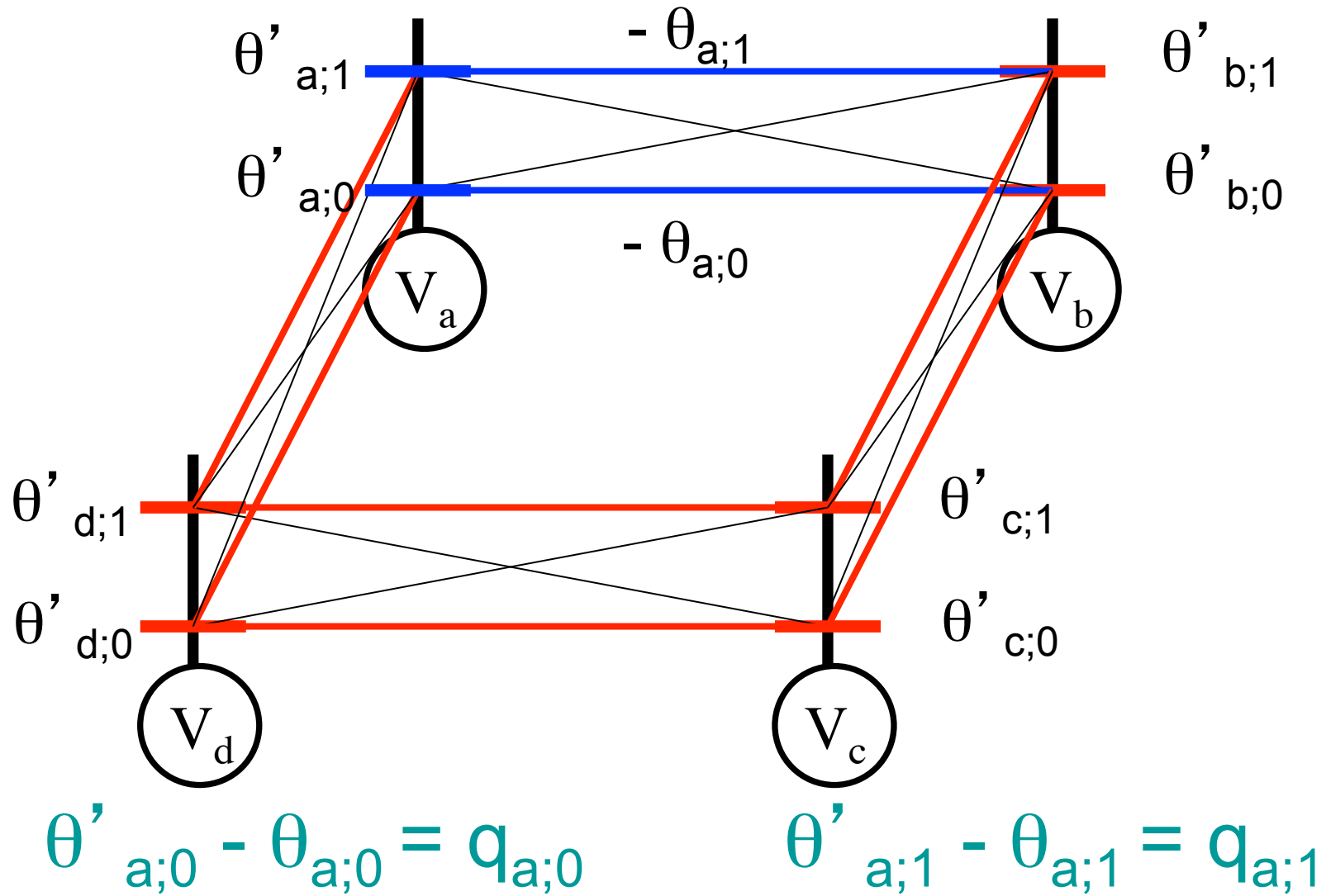
Potentials along the red path add up to 0

# Belief Propagation on Cycles



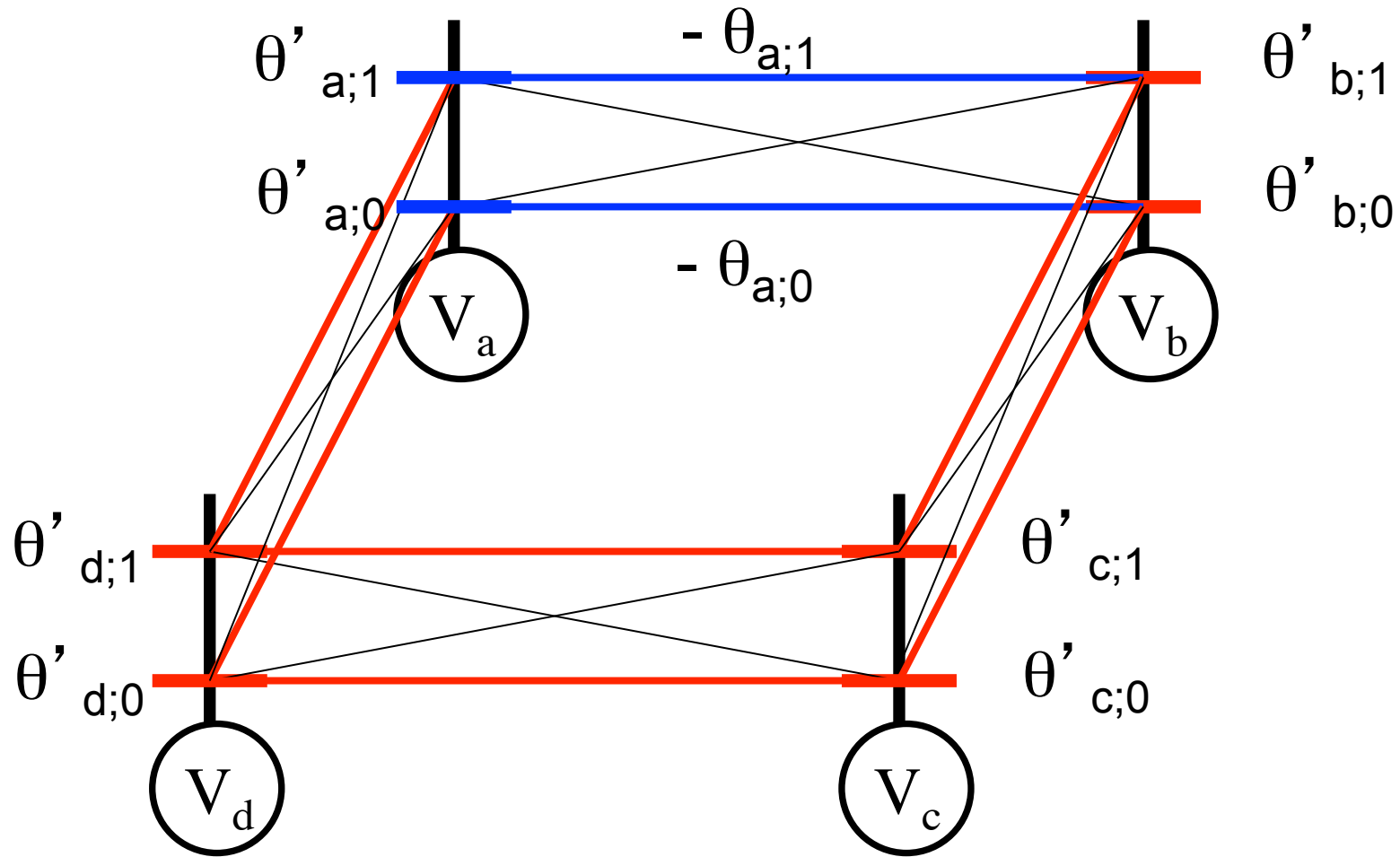
Potentials along the red path add up to 0

# Belief Propagation on Cycles



Potentials along the red path add up to 0

# Belief Propagation on Cycles

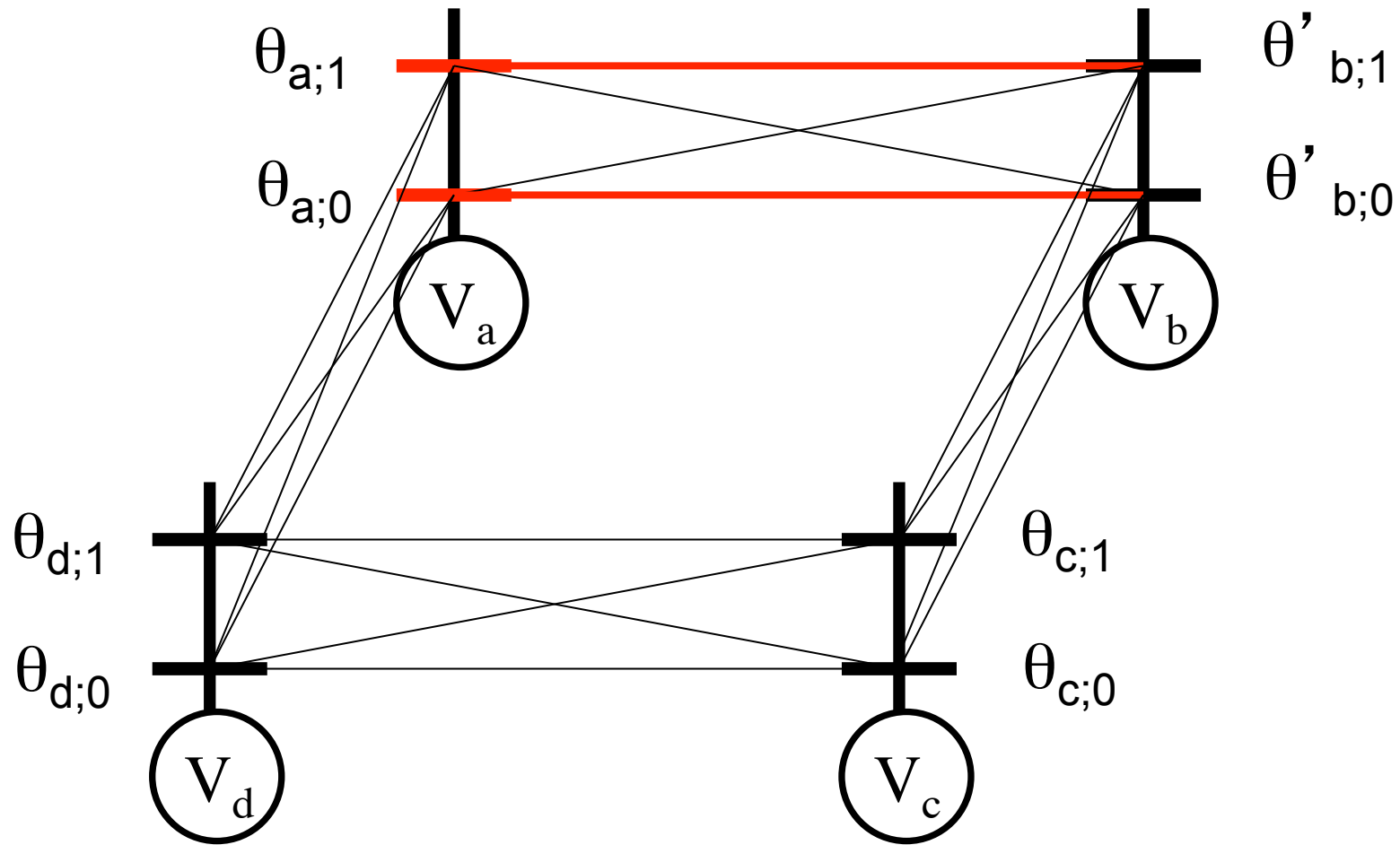


$$\theta'_{a;0} - \theta_{a;0} = q_{a;0}$$

$$\theta'_{a;1} - \theta_{a;1} = q_{a;1}$$

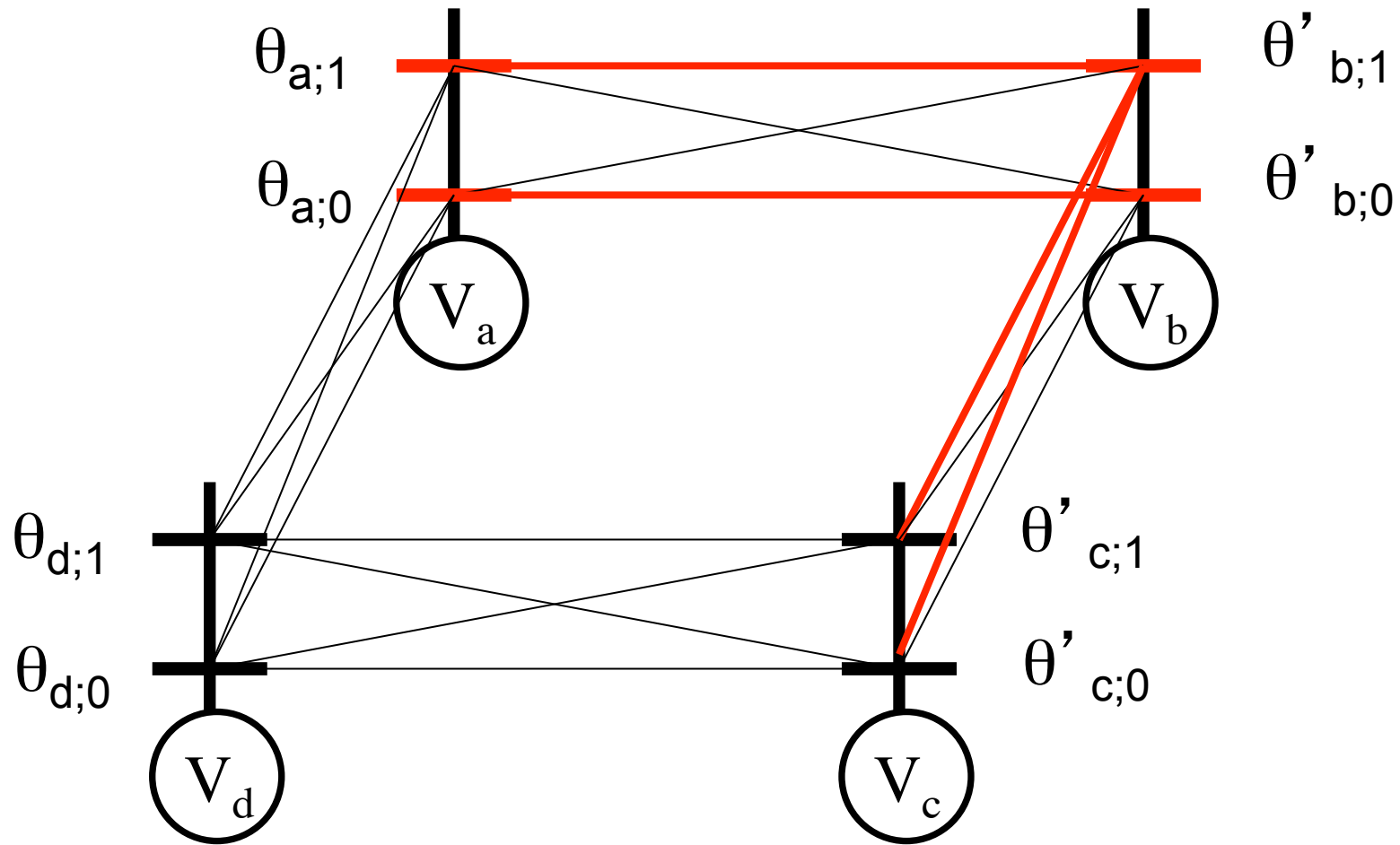
Pick minimum min-marginal. Follow red path.

# Belief Propagation on Cycles



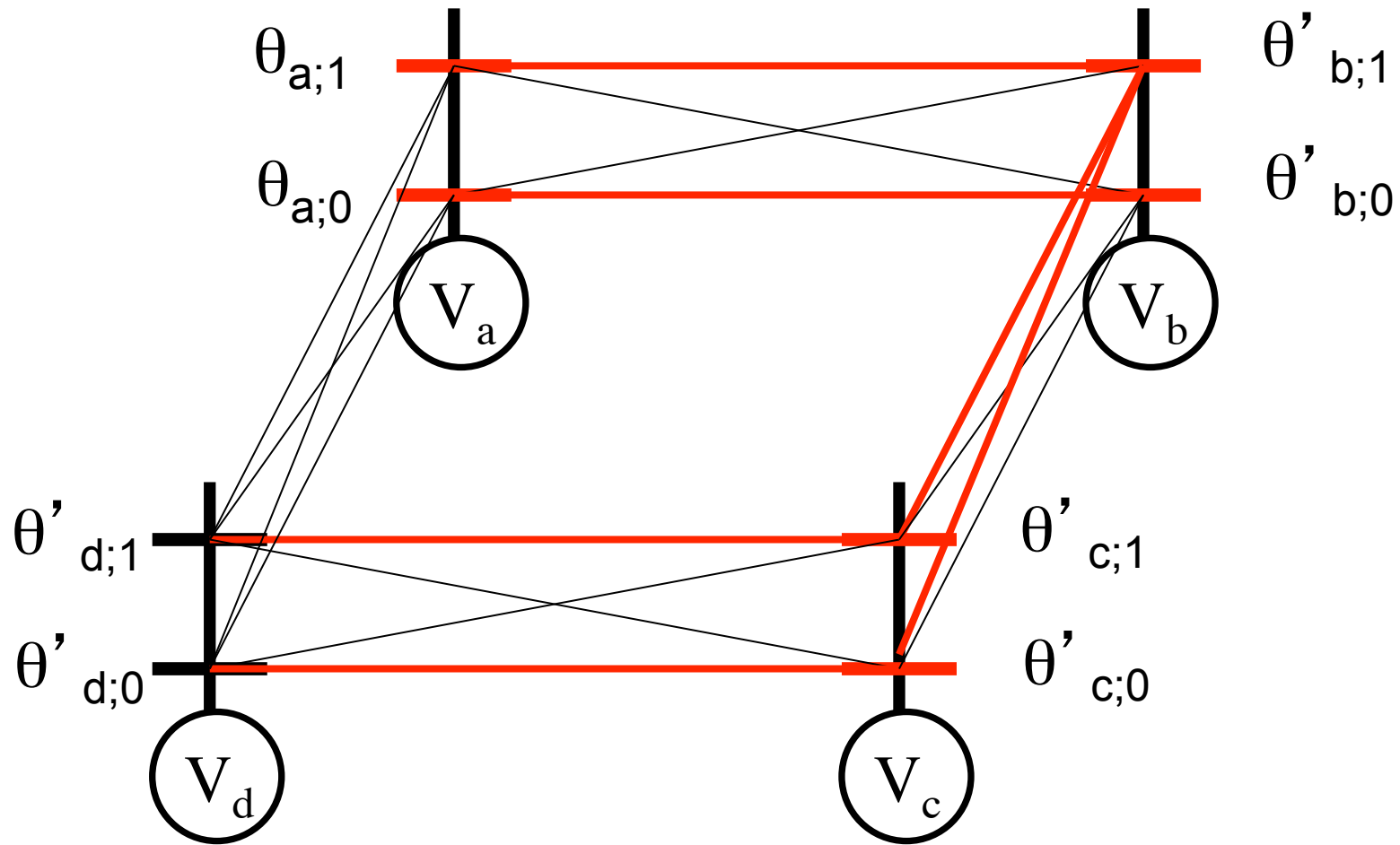
Potentials along the red path add up to 0

# Belief Propagation on Cycles



Potentials along the red path add up to 0

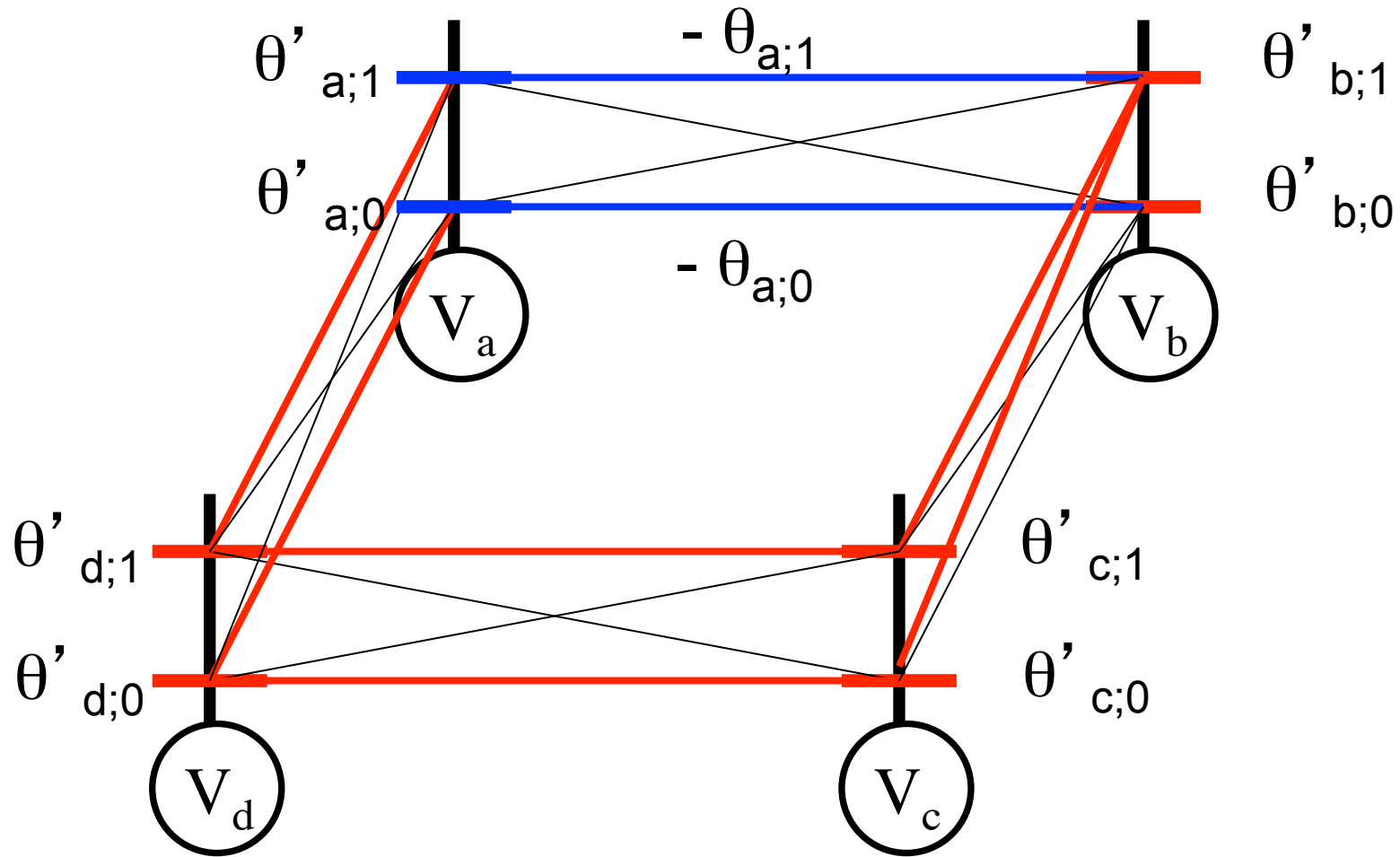
# Belief Propagation on Cycles



Potentials along the red path add up to 0



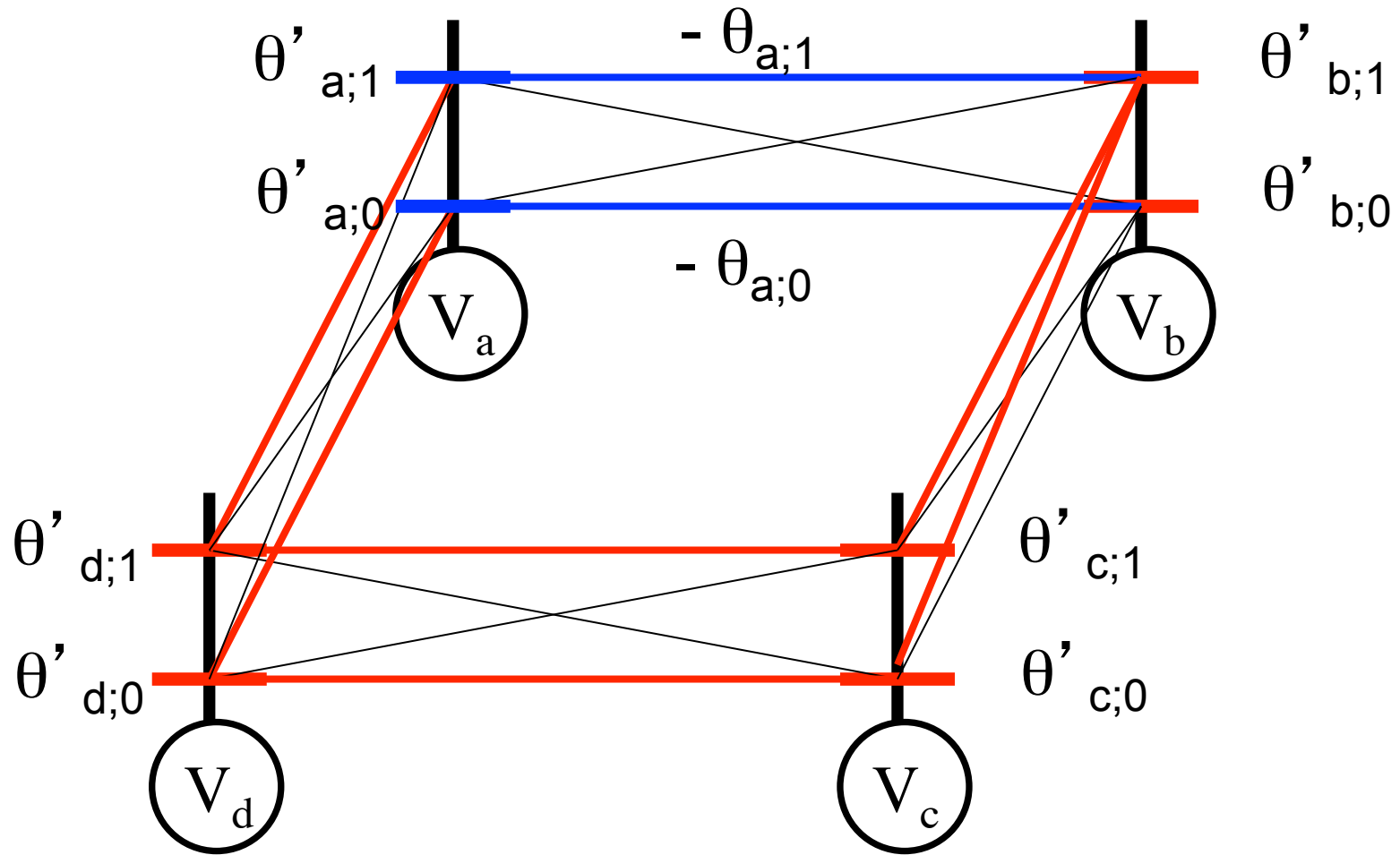
# Belief Propagation on Cycles



$$\theta'_{a;1} - \theta_{a;1} = q_{a;1} \leq \theta'_{a;0} - \theta_{a;0} = q_{a;0}$$

Potentials along the red path add up to 0

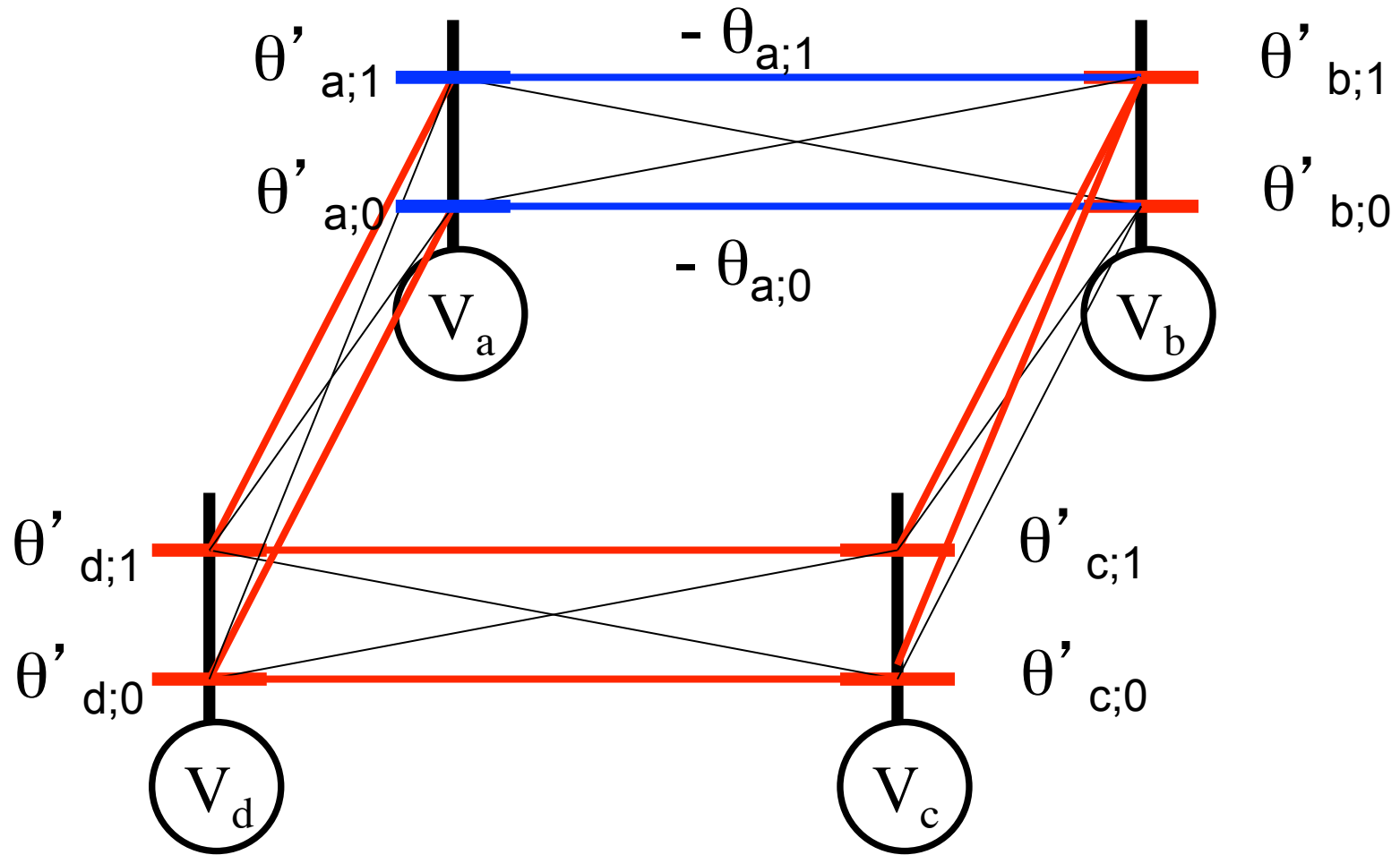
# Belief Propagation on Cycles



$$\theta'_{a;1} - \theta_{a;1} = q_{a;1} \leq \theta'_{a;0} - \theta_{a;0} = q_{a;0}$$

Problem Solved

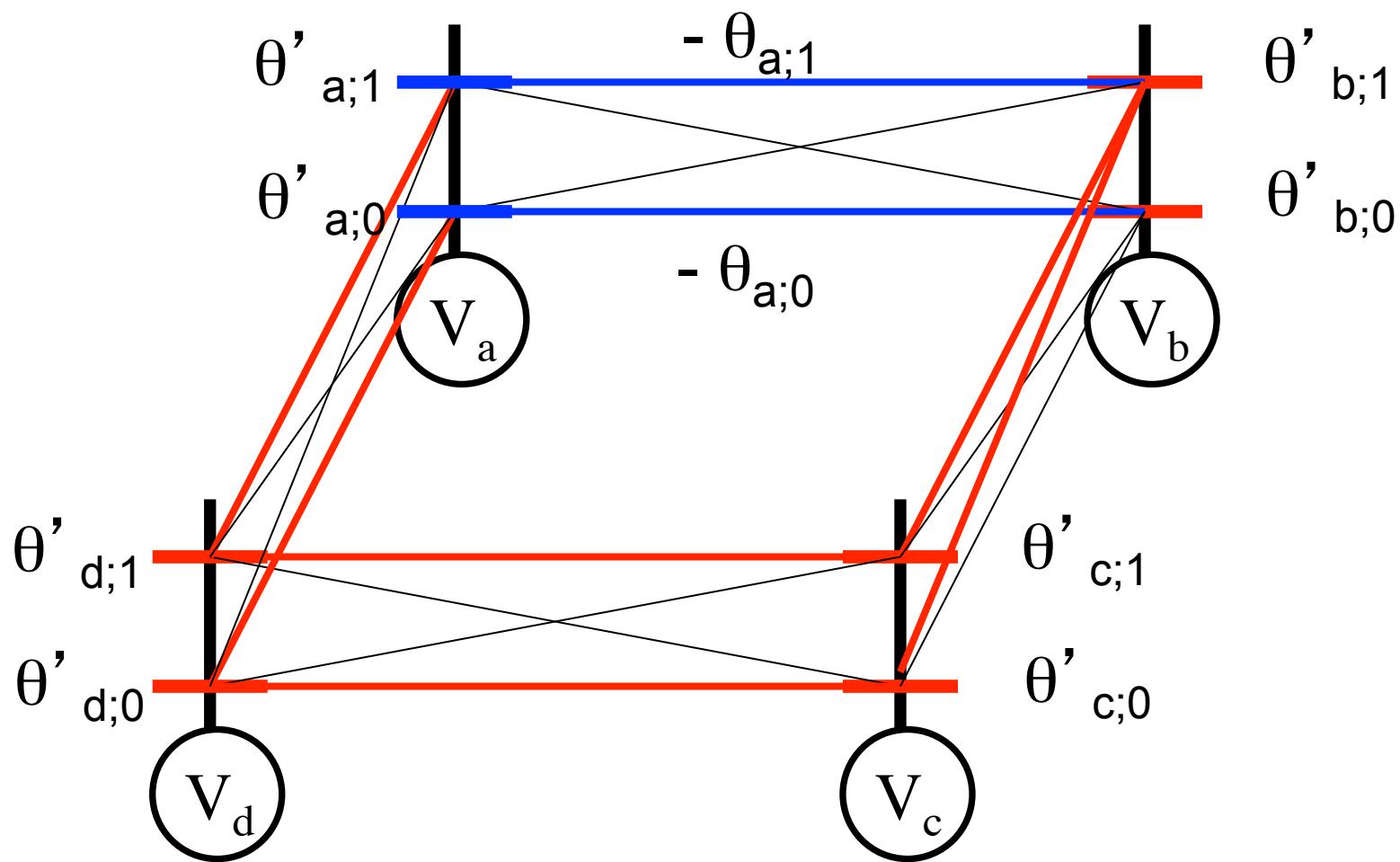
# Belief Propagation on Cycles



$$\theta'_{a;1} - \theta_{a;1} = q_{a;1} \geq \theta'_{a;0} - \theta_{a;0} = q_{a;0}$$

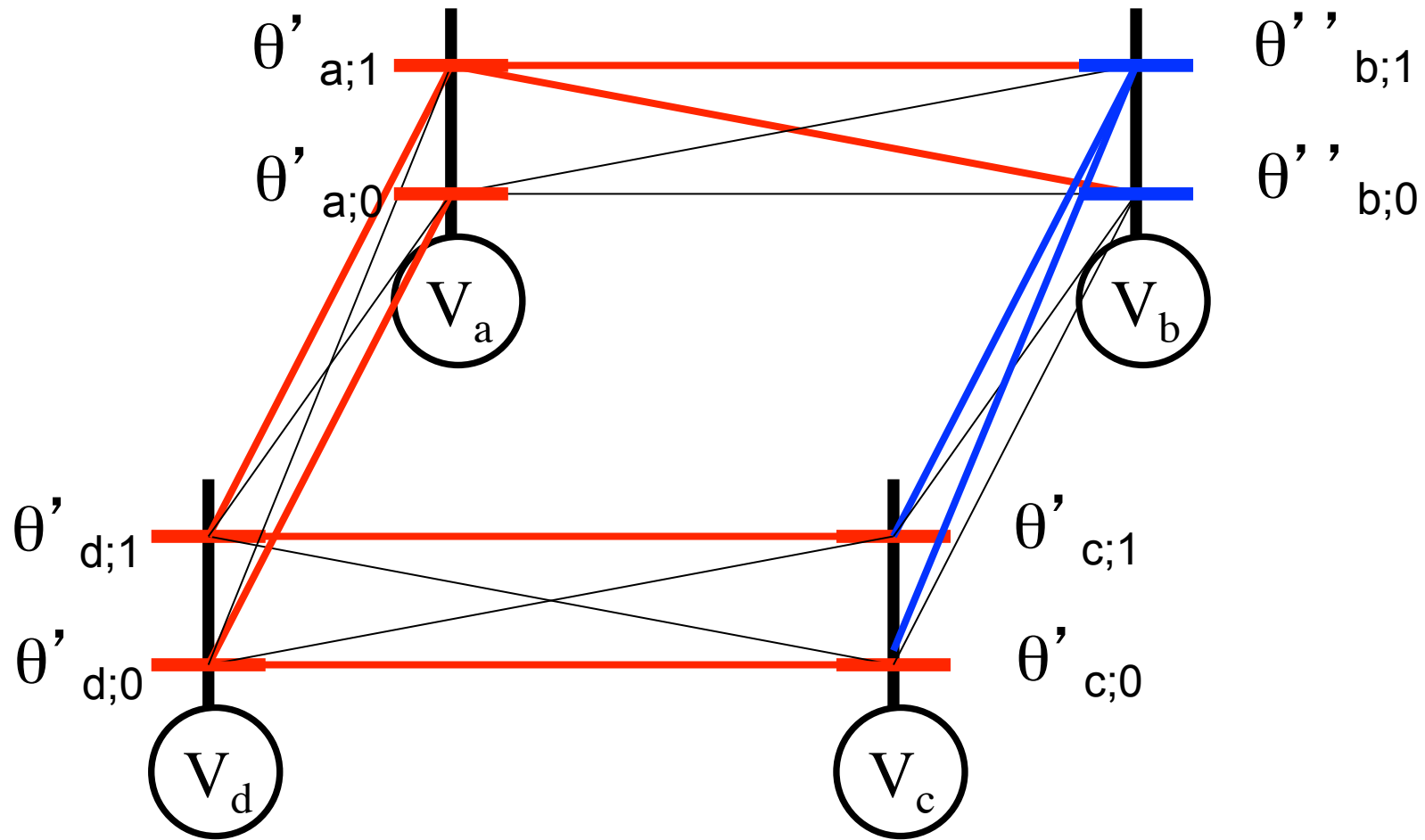
Problem Not Solved

# Belief Propagation on Cycles



Reparameterize (a,b) again

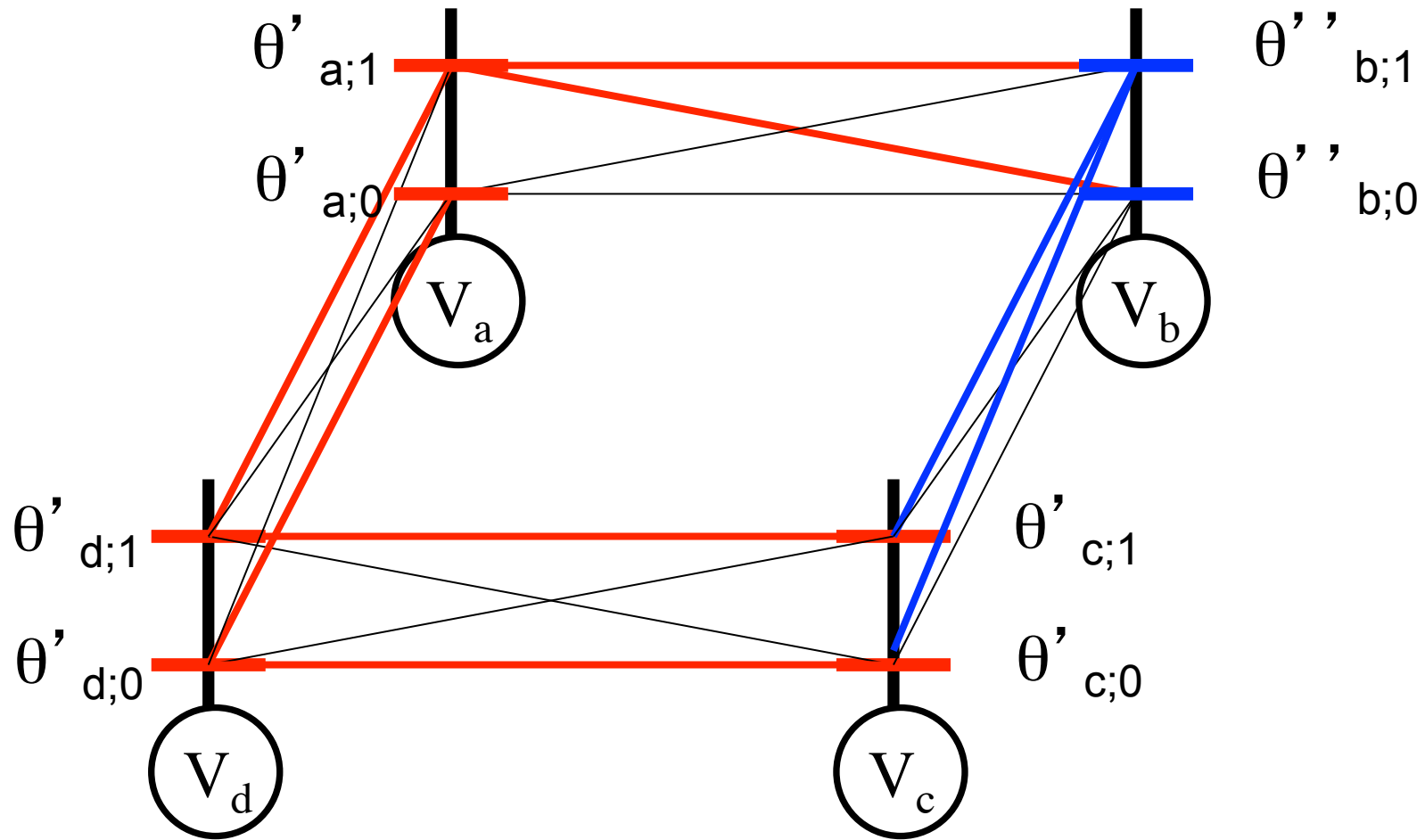
# Belief Propagation on Cycles



Reparameterize (a,b) again

But doesn't this overcount some potentials?

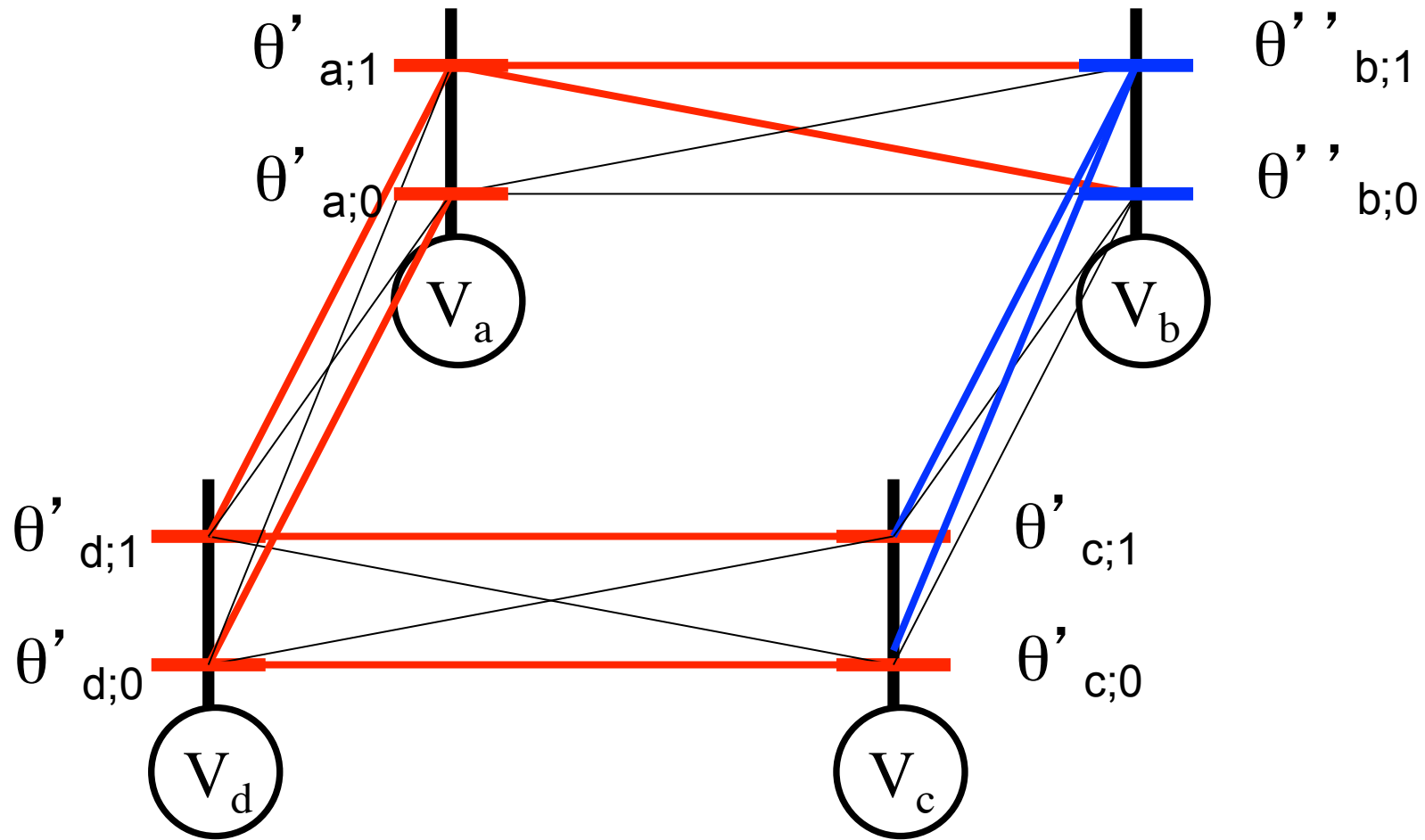
# Belief Propagation on Cycles



Reparameterize (a,b) again

Yes. But we will do it anyway

# Belief Propagation on Cycles



Keep reparameterizing edges in some order

Hope for convergence and a good solution

# Belief Propagation

- Generalizes to any arbitrary random field
- Complexity per iteration ?

$$O(|E||L|^2)$$

- Memory required ?

$$O(|E||L|)$$



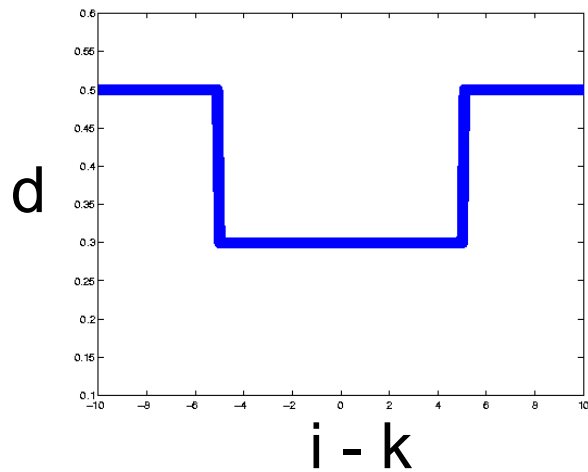
# Computational Issues of BP

Complexity per iteration

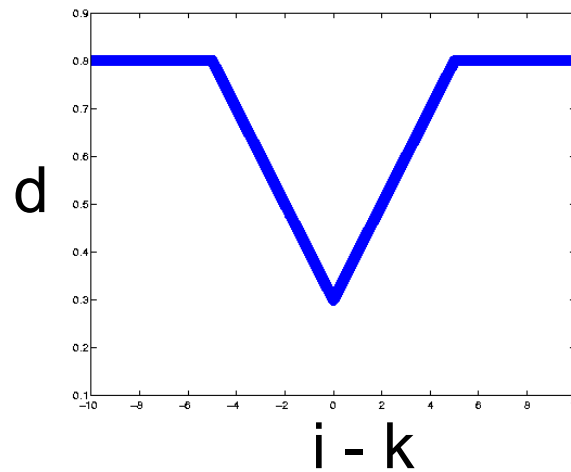
$$O(|E||L|^2)$$

Special Pairwise Potentials

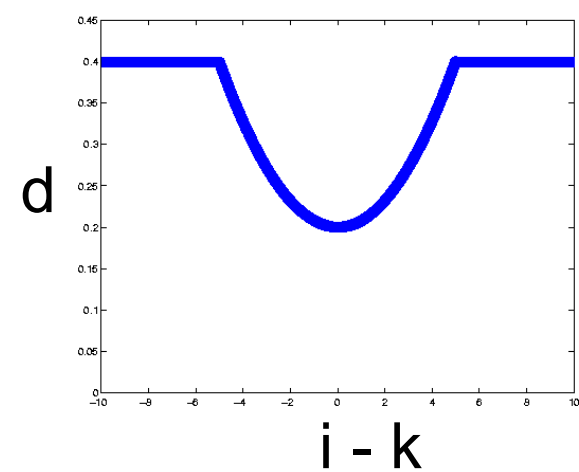
$$\theta_{ab;ik} = w_{ab}d(|i-k|)$$



Potts



Truncated Linear



Truncated Quadratic

$$O(|E||L|)$$

Felzenszwalb & Huttenlocher, 2004

# Computational Issues of BP

Memory requirements

$$O(|E||L|)$$

Half of original BP

Kolmogorov, 2006

Some approximations exist

Yu, Lin, Super and Tan, 2007

Lasserre, Kannan and Winn, 2007

**But memory still remains an issue**

# Computational Issues of BP

Order of reparameterization

Randomly

In some fixed order

The one that results in maximum change

Residual Belief Propagation

Elidan et al., 2006

# Summary of BP

Exact for chains

Exact for trees

Approximate MAP for general cases

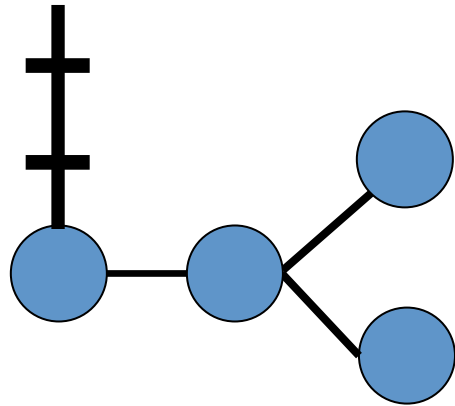
Not even convergence guaranteed

So can we do something better?

# Other alternatives

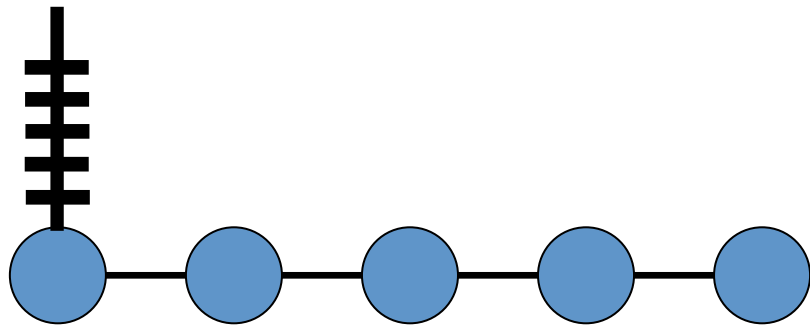
- Integer linear programming and relaxation
- TRW, Dual decomposition methods
- Extensively studied
  - Schlesinger, 1976
  - Koster et al., 1998, Chekuri et al., '01, Archer et al., '04
  - Wainwright et al., 2001, Kolmogorov, 2006
  - Globerson and Jaakkola, 2007, Komodakis et al., 2007
  - Kumar et al., 2007, Sontag et al., 2008, Werner, 2008
  - Batra et al., 2011, Werner, 2011, Zivny et al., 2014

# Where do we stand ?



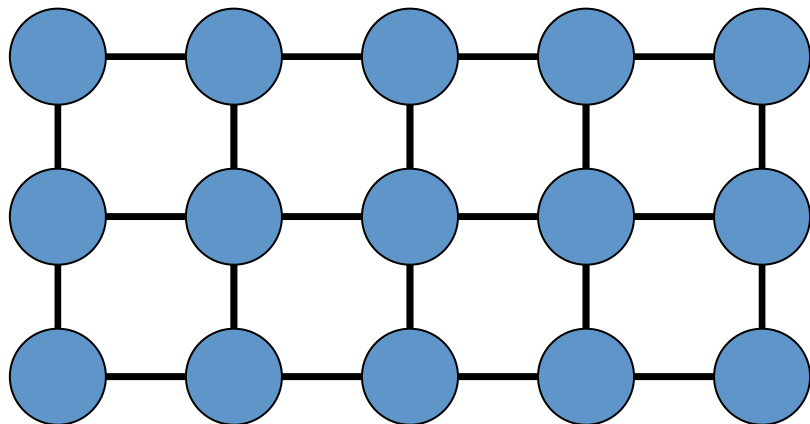
Chain/Tree, 2-label:

Use BP



Chain/Tree, multi-label:

Use BP



Grid graph:

Use TRW,  
dual decomposition,  
relaxation

# Note on Dynamic Programming

# Dynamic Programming (DP)

- DP  $\approx$  “careful brute force”
- DP  $\approx$  recursion + memoization + guessing
- Divide the problem into subproblems that are connected to the original problem
- Graph of subproblems has to be acyclic (DAG)
- Time = #subproblems  $\cdot$  time/subproblem



# 5 easy steps of DP

Analysis:

1. Define subproblems #subproblems
2. Guess part of solution #choices
3. Relate subproblems (recursion) time/subproblem
4. Recurse + memoize time  
OR build DP table bottom-up  
- check subprobs be acyclic / topological order
5. Solve original problem extra time

# 5 easy steps of DP

	Fibonacci	Shortest paths
1. Subproblems	$F_k, 1 \leq k \leq n$	$\delta_k(s, v), v \in V, 0 \leq k \leq V$
#subproblems	$n$	$V^2$
2. Guessing	$F_{n-1}, F_{n-2}$	edges coming into $v$
#choices	$1$	$\text{indegree}(v)$
3. Recurrence	$F_n = F_{n-1} + F_{n-2}$	$\delta_k(s, v) = \min\{ \delta_{k-1}(s, u) + w(u, v) \mid (u, v) \in E \}$
time/subproblem	$O(1)$	$O(\text{indegree}(v))$
4. Topological order	for $k = 1, \dots, n$	for $k = 0, 1, \dots, V - 1$ for $v \in V$
total time	$O(n)$	$O(V E)$
5. Original problem	$F_n$	$\delta_{V-1}(s, v)$
extra time	$O(1)$	$O(1)$